

A The Borel hierarchy

These exercises concern the details behind the proof of Aumann's (1961) theorem. Flicking through, you'll see there's quite a lot to cover, but the rest of the material doesn't depend on this technical development. It's only here to satisfy your curiosity about what happens deep inside the σ -algebra of Borel sets. If you enjoy these, take a closer look at *descriptive set theory*. Two classical textbooks are Moschovakis's (1987) selection of key, central results, and Kechris's (1995) comprehensive, detailed, and slightly more modern book.

Define by transfinite induction on $\omega_1 + 1$, the successor of the first uncountable ordinal:

$$\Sigma_\alpha^\mathcal{U}, \Pi_\alpha^\mathcal{U}, \Delta_\alpha^\mathcal{U} \subseteq \wp X \quad (\alpha \in \omega_1)$$

$$\Sigma_1^\mathcal{U} := \mathcal{U}$$

$$\Sigma_{\alpha+1}^\mathcal{U} := \left\{ \bigcup_{i \in I} A_i \mid I \subseteq \mathbb{N}, \vec{A} \in \mathcal{U} \cup \bigcup_{\beta \leq \alpha} \Pi_\beta^\mathcal{U} \right\} \quad (1 \leq \alpha \in \omega_1)$$

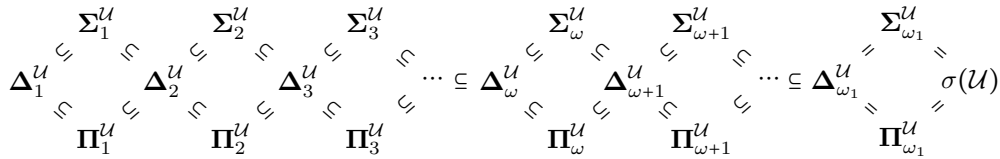
$$\Sigma_\gamma^\mathcal{U} := \bigcup_{\beta < \gamma} \Sigma_\beta^\mathcal{U} \quad (1 \leq \gamma \text{ a limit ordinal in } \omega_1)$$

$$\Pi_\alpha^\mathcal{U} := [\Sigma_\alpha^\mathcal{U}]^C := \{A^C \mid A \in \Sigma_\alpha^\mathcal{U}\} \quad \Delta_\alpha^\mathcal{U} := \Sigma_\alpha^\mathcal{U} \cap \Pi_\alpha^\mathcal{U}$$

▮ **A.1.** For every $\alpha \leq \omega_1$, we have $\Sigma_\alpha^\mathcal{U} \cup \Pi_\alpha^\mathcal{U} \subseteq \Delta_{\alpha+1}^\mathcal{U}$. ▮

▮ **A.2.** Prove that $\sigma(\mathcal{U}) = \Sigma_{\omega_1}^\mathcal{U} = \Pi_{\omega_1}^\mathcal{U} = \Delta_{\omega_1}^\mathcal{U}$. ▮

We therefore have the following relationships between the classes of the *Borel hierarchy*:



Given a set V whose elements represent variables, the σ -terms over V are the countably-infinitary terms generated by the following grammar:

$$t, s ::= x \mid x^C \mid \bigcup_{i \in I} t_i \mid \bigcap_{i \in I} t_i \quad (x \in V, I \subseteq \mathbb{N})$$

Given a *valuation* $e : V \rightarrow \sigma(\mathcal{U})$, we can interpret each σ -term t as a Borel subset $\llbracket t \rrbracket e \in \sigma(\mathcal{U})$. Note that every term t involves only countably many variables, we call these variables its *support*.

▮ **A.3.** Let $\mathcal{U} \subseteq \wp X$, $\mathcal{V} \subseteq \wp Y$. Show that for every measurable $f : \langle X, \sigma(\mathcal{U}) \rangle \rightarrow \langle Y, \sigma(\mathcal{V}) \rangle$, the inverse image f^{-1} is a homomorphism of σ -terms: ▮

$$f^{-1}[\llbracket t \rrbracket e] = \llbracket t \rrbracket (f^{-1} \circ e) \quad \Delta$$

▮ **A.4.** Show that if $e : V \rightarrow \mathcal{U}$ is surjective, then $\llbracket - \rrbracket e$ is surjective on $\sigma(\mathcal{U})$. ▮

We call a term *alternating* when, for every non-variable sub-term $f \langle t_i \rangle_i$, the root of each direct sub-tree is not the same operation symbol f .

▮ **A.5.** Show that every term is denotationally equivalent to an alternating term. You might enjoy presenting a denotation-preserving terminating rewriting system. ▮

The *Aumann rank* function assigns to each Borel set the first stage in the hierarchy in which it occurs in some $\Sigma^{\mathcal{U}}$ set:

$$\text{rank}^{\mathcal{U}} : \sigma(\mathcal{U}) \rightarrow \omega_1$$

$$\text{rank}^{\mathcal{U}} A := \min \{ \alpha \in \omega_1 \mid A \in \Sigma_{\alpha}^{\mathcal{U}} \}$$

Define the *alternating depth* of a σ -term as follows:

$$\text{alter} : \sigma\text{-Term} V \rightarrow \omega_1$$

$$\text{alter } x := \text{alter } x^{\mathbb{C}} := 0 \quad \text{alter } \bigcup_{i \in I} t_i := \bigvee_{i \in I} \text{alter } t_i \quad \text{alter } \bigcap_{i \in I} t_i := \bigvee_{i \in I} \text{alter } t_i + \bigvee_{i \in I} \left[t_i \neq \bigcap_{j \in J} s_j \right]$$

▮ **A.6.** Let t be a σ -term and e a valuation in some \mathcal{U} .

- Show that $\llbracket t \rrbracket e \in \Sigma_{\text{alter } t \vee \alpha}^{\mathcal{U}}$, where $\alpha := \bigvee_{x \in \text{supp } t} e(x) \in \omega_1$.
- Deduce that if $e : V \rightarrow \mathcal{U}$, then $\text{rank } \llbracket t \rrbracket e \leq \text{alter } t$. Generalise to any $e : V \rightarrow \sigma(\mathcal{U})$.
- Show that $\text{rank } A = \min \{ \text{alter } t \mid A = \llbracket t \rrbracket e \}$. ▮

▮ **A.7.** Prove that if $A \in \sigma(\mathcal{U})$ and $\rho := \text{rank}^{\mathcal{U}} A$, then:

$$\begin{aligned} A \cap [\Sigma_{\alpha}^{\mathcal{U}}] &\subseteq \Sigma_{\alpha}^{A \cap [\mathcal{U}]} \subseteq \Sigma_{(\rho+1) \vee \alpha}^{\mathcal{U}} & A \cap [\Pi_{\alpha}^{\mathcal{U}}] \\ &\subseteq \Pi_{\alpha}^{A \cap [\mathcal{U}]} \subseteq \Pi_{(\rho+1) \vee \alpha}^{\mathcal{U}} & A \cap [\Delta_{\alpha}^{\mathcal{U}}] \subseteq \Delta_{\alpha}^{A \cap [\mathcal{U}]} \subseteq \Delta_{(\rho+1) \vee \alpha}^{\mathcal{U}} \end{aligned} \quad \triangleleft$$

Let $\mathcal{U} \subseteq \wp X$, $\mathcal{V} \subseteq \wp Y$. When \mathcal{V} is countable, define:

$$\text{rank}^{\mathcal{U}, \mathcal{V}} : \mathbf{Meas}(\langle X, \sigma(\mathcal{U}) \rangle, \langle Y, \sigma(\mathcal{V}) \rangle) \rightarrow \omega_1$$

$$\text{rank } f := \bigvee_{A \in \mathcal{V}} f^{-1}[A]$$

Let $f : \langle X, \sigma(\mathcal{U}) \rangle \rightarrow \langle Y, \sigma(\mathcal{V}) \rangle$ be a measurable function.

▮ **A.8.** What's the rank of a continuous function between two topological spaces? ▮

▮ **A.9.** Bound the rank of $f^{-1}[A]$ for every $A \in \sigma(\mathcal{V})$, using $\text{rank } A$ and $\text{rank } f$.

Is your bound tight enough to deduce that $\text{rank } f^{-1}[A] \leq \text{rank } A$ when f is continuous for the topologies generated by \mathcal{U} and \mathcal{V} ? ▮

Let $\mathcal{U} \subseteq \wp(C \times X)$ and $\mathcal{V} \subseteq \wp X$ be two classes of subsets. We will regard subsets $\llbracket - \rrbracket \in \mathcal{U}$ as potential encodings for subsets in \mathcal{V} , where each element $c \in C$ encodes the section subset $\llbracket c \rrbracket := \{x \in X \mid x \in \llbracket c \rrbracket\}$.

We say that $\llbracket - \rrbracket \in \mathcal{U}$ is a \mathcal{U} - \mathcal{V} -encoder when $\mathcal{V} = \{\llbracket c \rrbracket \mid c \in C\}$. The intended meaning is that such an encoder lets us cover all the \mathcal{V} -subsets with a code in C . The literature uses the term *C-universal set for Ξ* for a \mathcal{U} - \mathcal{V} -encoder, when \mathcal{U} and \mathcal{V} belong to the same family of subset classes Ξ , such as $\mathcal{U} = \Sigma_{\alpha}(C \times X)$ and $\mathcal{V} = \Sigma_{\alpha}(X)$.

▮ **A.10.** Show that if $\llbracket - \rrbracket$ is a \mathcal{U} - \mathcal{V} -encoder, then $\llbracket - \rrbracket^{\mathbb{C}}$ is a $[\mathcal{U}]^{\mathbb{C}}$ - $[\mathcal{U}]^{\mathbb{C}}$ -encoder. ▮

▮ **A.11.** Let $\llbracket - \rrbracket$ be a \mathcal{U} - \mathcal{V} encoder, where $\mathcal{U} \subseteq \wp(C \times C)$ and $\mathcal{V} \subseteq \wp C$. Consider the diagonal function $\Delta := \lambda x. \langle x, x \rangle : C \rightarrow C \times C$.

Show that $\Delta^{-1}[\llbracket - \rrbracket^{\mathbb{C}}] \notin \mathcal{V}$. ▮

We'll use this diagonalisation technique to show that the Borel hierarchy doesn't collapse for the reals.

▮ **A.12.** Recall the Cantor space $\mathbb{G} \subseteq \mathbb{R}$, let \mathcal{V} be the open subsets of \mathbb{R} , let $\mathcal{V}' := \mathbb{G} \cap [\mathcal{U}]$ be the open subsets in \mathbb{G} , and \mathcal{U}' be the open subsets of $\mathbb{G} \times \mathbb{G}$.

- Show that if, for all $1 \leq \alpha < \omega_1$, we have $\Sigma_\alpha^{\mathcal{V}'} \neq \Pi_\alpha^{\mathcal{V}'}$, then $\Sigma_\alpha^{\mathcal{V}} \neq \Pi_\alpha^{\mathcal{V}}$ too, and so the Borel hierarchy for \mathbb{R} only stabilises at ω_1 .
- Show that if \mathbb{G} has a $\Sigma_\alpha^{\mathcal{U}'}\text{-}\Sigma_\alpha^{\mathcal{V}'}$ -encoder, then $\Sigma_\alpha^{\mathcal{V}'} \neq \Pi_\alpha^{\mathcal{V}'}$.
- Show that, for all $1 \leq \alpha \in \omega_1$, \mathbb{G} has both a $\Sigma_\alpha^{\mathcal{U}'}\text{-}\Sigma_\alpha^{\mathcal{V}'}$ encoder and a $\Pi_\alpha^{\mathcal{U}'}\text{-}\Pi_\alpha^{\mathcal{V}'}$ encoder. \triangleleft

The last exercise constructs a non-Borel set. This result doesn't fit the narrative, but we've already introduced most of the tools required for the job.

▮ **A.13.** A Borel set is *analytic* when it is empty, or a continuous image of the Baire space $\mathbb{Y} := \mathbb{N}^{\mathbb{N}}$. We denote by $\Sigma_1^1(S)$ the class of analytic subsets of S . One can show that every Borel set is analytic, but that would require a lot of additional machinery.

- Show that if $\mathcal{B}_{\mathbb{Y}} \subseteq \Sigma_1^1(\mathbb{Y})$ and we have a $\Sigma_1^1(\mathbb{Y} \times \mathbb{Y})\text{-}\Sigma_1^1(\mathbb{Y})$ -encoder, then $\mathcal{B}_{\mathbb{Y}} \subset \Sigma_1^1(\mathbb{Y})$.
- Show that we have a $\Pi_1^0(\mathbb{Y} \times \mathbb{Y})\text{-}\Pi_1^0(\mathbb{Y})$ -encoder.
- Construct a homeomorphism $\mathbb{Y} \cong \mathbb{Y} \times \mathbb{Y}$. Derive a $\Pi_1^0(\mathbb{Y} \times \mathbb{Y} \times \mathbb{Y})\text{-}\Pi_1^0(\mathbb{Y} \times \mathbb{Y})$ -encoder $\mathcal{F}[-]$.
- Show that setting $x \in [c]$ when $\exists z. \langle x, z \rangle \in \mathcal{F}[c]$ is an $\Sigma_1^1(\mathbb{Y} \times \mathbb{Y})\text{-}\Sigma_1^1(\mathbb{Y})$ -encoder.
Hint: the graph of a continuous function over \mathbb{Y} is a $\Pi_1^0(\mathbb{Y} \times \mathbb{Y})$ set. \triangleleft

References

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