

# Foundations for type-driven probabilistic modelling

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Computational golden era of:

logic & type rich  
computation

Statistical  
computation

# Computational golden era of:

logic & type rich  
computation

Expressive type systems:

Haskell, OCaml, Idris

Mechanised mathematics:

Agda, Coq, Isabelle/HOL, Lean

Verification:

SMT-powered, realistic  
systems

Statistical  
computation

generative modelling  
+

efficient inference:

Monte-Carlo simulation  
or gradient-based  
optimisation

"AI"

Computational golden era of:

logic type rich  
computation

Statistical  
computation

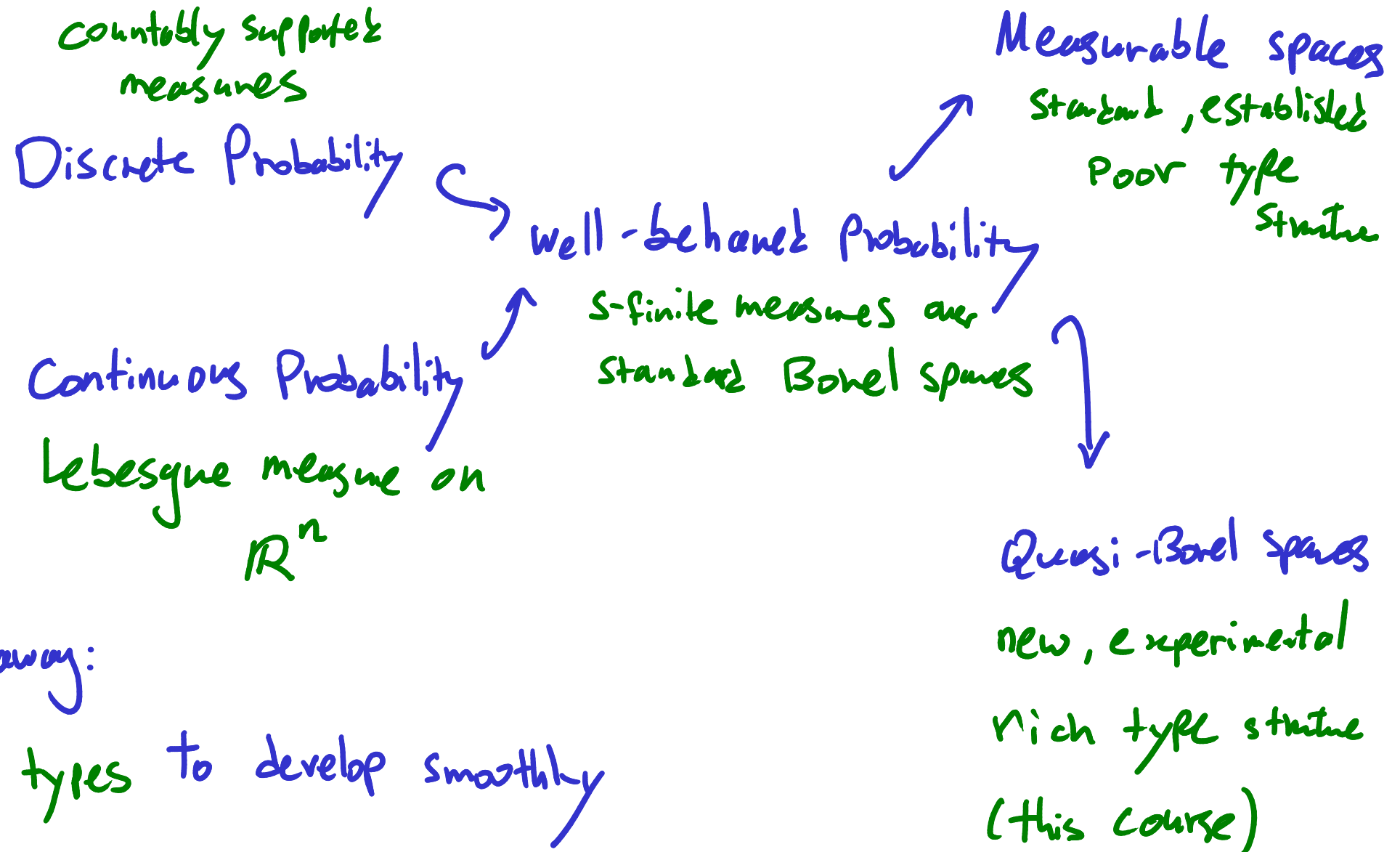
Clear connection to

Foundations:

- Reid's
- John's courses
- Michael's
- Dominik's

- this course

# Why foundations?



Takeaway:

use types to develop smoothly

Plan:

- 1) Type-driven probability: discrete case (Mon + Tue (?))
- 2) Borel sets & measurable spaces (Tue)
- 3) Quasi Borel spaces, simple type structure (Wed)
- 4) Dependent type structure & standard Borel spaces (Thu)
- 5) Integration & random variables (Fri)

please ask questions!



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# Language of distribution & Probability

$X$  type (=space) of values / outcomes

$DX$  type of distributions / measures over  $X$

$PX \subseteq DX$  Sub type of probability measures (total measure 1)

$BX$  type of measurable events - subsets of  $X$  we wish to measure

$W$  type of weights :  $[0, \infty]$

→ type judgment

$\mu : DX, E : BX \vdash c_{\mu}[E] : W$

↳ measure  $\mu$  assigns to  $E$



# Axioms for measures

---

Empty event:  $\emptyset : \mathcal{B}X$

Its measure is  $0 : \mathcal{W}$ :

$$\mu : \mathcal{D}X \vdash \underbrace{C_e[\emptyset]}_{\mu} = 0 : \mathcal{W}$$

# Axioms for measures

---

$\mathcal{B}X$  is a Boolean sub-algebra:

$$E : \mathcal{B}X \vdash E^c : \mathcal{B}X$$

$$E, F : \mathcal{B}X \vdash E \cup F, E \cap F : \mathcal{B}X$$

$E, C : \mathcal{B}X, \mu : \mathcal{D}X \vdash$  (disjoint additivity)

$$\mu[E] = \mu[E \cap C] + \mu[E \cap C^c] : \mathbb{W}$$

# Axioms for measures

---

$\omega := (\mathbb{N}, \leq)$      $(B, \subseteq)$      $(W, \leq)$  posets

$$(BX, \subseteq)^\omega := \left\{ (E_n)_{n \in \mathbb{N}} \in (BX)^\mathbb{N} \mid E_0 \subseteq E_1 \subseteq E_2 \subseteq \dots \right\}$$

$(BX, \subseteq)$  and  $(W, \leq)$  are  $\omega$ -chain-closed:

$$E_- : (BX, \subseteq)^\omega \vdash \bigcup_n E_n : BX \qquad a_- : (W, \leq)^\omega \vdash \sup_n a_n : W$$

$$E_- : (BX, \subseteq)^\omega, \mu : DX \vdash \qquad \text{(Scott continuity)}$$

$$C_e[\bigcup_n E_n] = \sup_n C_e[E_n] : W$$

# Axiom for Probability

$$\text{Cost} : PX \xrightarrow{\varepsilon} DX$$

$$1 : W$$

$$\mu : PX \vdash \underset{\text{Cost } \mu}{C_e[X]} = 1 : W$$

Avoid casting:

$$E : BX, \mu : PX \vdash \underset{\mu}{P_r[E]} := \underset{\text{Cost } \mu}{C_e[E]} : [0,1] \subseteq W$$

# Axioms for measures

Integration:

$$\mu: \mathcal{D}X, \varphi: \mathcal{W}^X \mapsto \int \mu \varphi : \mathcal{W} \quad (\text{Lebesgue integral})$$

Again, avoid casting:

$$\mu: \mathcal{P}X, \varphi: \mathcal{W}^X \mapsto \underbrace{E[\varphi]}_{\mu} := \int (\text{cast } \mu) \varphi : \mathcal{W} \quad (\text{Expectation})$$

More structure & notation later (...technical...)

Have: language + axioms

Want: model

today: discrete measures

rest of course: discrete + continuous

# Discrete model

type  $X$ : set

$$DX := \left\{ \mu: X \rightarrow \mathbb{W} \mid \mu \text{ is countably supported} \right\}$$

(next slide)

# Support

→ Powerset

$\mu: W^X$ ,  $S: \mathcal{P}X \vdash S$  supports  $\mu :=$

$\forall x: X. \mu x > 0 \Rightarrow x \in S$  : Prop

$\mu: W^X \vdash \text{supp } \mu := \{x \in X \mid \mu x > 0\} : \mathcal{P}X$

supp  $\mu$  is the smallest set supporting  $\mu$



# Discrete model

type  $X$ : set

$$DX := \{ \mu: X \rightarrow \mathbb{W} \mid \mu \text{ is countably supported} \}$$

$$:= \{ \mu: X \rightarrow \mathbb{W} \mid \text{supp } \mu \text{ is countable} \}$$

## Ex. measures

- $X$  ctbl, Counting measure  $\#_X : DX$

$$\#_X := \lambda x : X. 1 \quad (\text{NB: } \text{Supp} \#_X = X \quad \checkmark \text{ ctbl})$$

- Dirac measure:

$$x : X \mapsto \delta_x := \lambda x'. \begin{cases} x = x' : 1 \\ \text{o.w.} : 0 \end{cases} : DX$$

$$\text{NB: } \text{Supp} \delta_x = \{x\} \quad \checkmark \text{ ctbl}$$

- Zero measure  $\underline{0} := \lambda x. 0 : DX$

$$\text{NB: } \text{Supp} \underline{0} = \emptyset \quad \checkmark \text{ ctbl}$$

# Discrete model

type  $X$ : set

$DX := \{ \mu: X \rightarrow \mathbb{W} \mid \mu \text{ is countably supported} \}$

$$\mu: DX, E: BX \vdash \underbrace{C_e[E]}_{\mu} := \sum_{x \in E} \mu x$$

$$:= \sum_{x \in E \cap \text{Supp } \mu} \mu x$$

Lemma:  $\mu: DX, S \in \mathcal{P}_{\text{ctbl}} X, S \text{ supports } \mu, E: BX \vdash$

$$\underbrace{C_e[E]}_{\mu} = \sum_{x \in E \cap S} \mu x$$

Ex:

$$\bullet E: \mathcal{B}X \vdash \quad C_e[E] = |E| := \begin{cases} E \text{ has } n \text{ elements:} & n \\ E \text{ infinite:} & \infty \end{cases}$$

$\#_x$

$$\bullet E: \mathcal{B}X, n: X \vdash \quad C_e[E] = \begin{cases} x \in E: & 1 \\ x \notin E: & 0 \end{cases} =: [x \in E] : \mathbb{W}$$

$\delta_n$

$$\text{NB: } E: \mathcal{B}X \vdash [- \in E] : X \rightarrow \mathbb{W}$$

indicator  
function

$$\bullet E: \mathcal{B}X \vdash \quad C_e[E] = 0$$

0

# Validate axioms

$$\mu:DX \vdash C_{\mu}[E] = 0 \quad : W$$

$$E, C : BX, \mu:DX \vdash$$

$$C_{\mu}[E] = C_{\mu}[E \cap C] + C_{\mu}[E \cap C^c] \quad : W$$

$$E_{-} : (BX, \subseteq)^W, \mu:DX \vdash$$

$$C_{\mu}[\bigcup_n E_n] = \sup_n C_{\mu}[E_n] \quad : W$$

# Kernels

$\kappa$  from  $\Gamma$  to  $X$ :

$$\kappa : (\mathcal{D}X)^\Gamma$$

kernels are "open/parameterised" measures

Ex: Dirac kernel.  $\delta_- : (\mathcal{D}X)^X$

# Kock Integral

$$\mu : D\Gamma, \kappa : DX^\Gamma \vdash \oint \mu \kappa : DX$$

In discrete model:

$$\oint \mu \kappa := \lambda x : X. \sum_{\nu \in \Gamma} \mu \nu \cdot \overbrace{k(\nu; x)}{:= k \nu x}$$

# (Weak) disintegration problem:

Input:  $\mu: D\Gamma$   $V: DX$

Output: a kernel  $k: (DX)^\Gamma$  s.t.

$$\int \mu k = V$$

Call such  $k$  a (weak) disintegration of  $V$

w.r.t.  $\mu$ .

(non-standard terminology)



Ex disintegration:

$$\underline{n} := \{0, 1, 2, \dots, n-1\}$$

disintegrate  $\#_{\underline{z}^{n+1}}$  w.r.t.  $\#_{\underline{z}}$

$$k: (D(\underline{z}^{n+1}))^{\underline{z}} \quad k(x; f) := \begin{cases} f(n) = x: & 1 \\ \text{o.w.} & : 0 \end{cases}$$

$$\left( \int \#_{\underline{z}} k \right) f = \sum_{x \in \underline{z}} \overbrace{\#_{\underline{z}}^1 x}^1 \cdot k(x; f)$$

NB:  $\text{Supp}(kx)$   
 $\sqrt{c + b}$

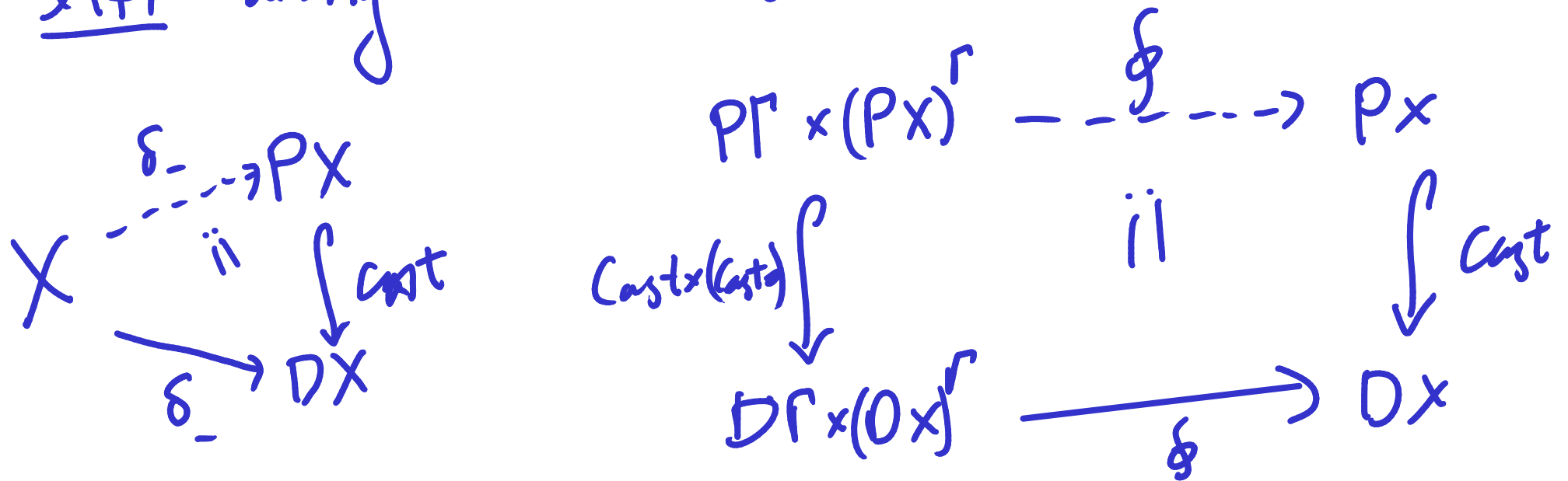
$$= k(0; f) + k(1; f) = k(fn; f) = 1 = \#_{\underline{z}^{n+1}}(f)$$

# Probability measures

$$PX := \{ \mu : DX \mid C_{\mu}[X] = 1 \} \xrightarrow{\subseteq} DX$$

Lemma:  $\delta_- : X \rightarrow DX$  and  $\wp : D\Gamma \times (DX)^{\Gamma} \rightarrow DX$

lift along the inclusion  $\text{cast} : P \xrightarrow{\subseteq} D :$



Prop (discrete Giry):

(Michèle Giry '82)

$(D, \delta_-, \oint)$  is a monad i.e.

$$\mu: \Gamma, \kappa: (DX)^\Gamma \vdash \oint \delta_\mu \kappa = \kappa r$$

$$\mu: DX \vdash \oint \mu(\delta x) \delta_x = \mu$$

: DX

$$\mu: D\Gamma, \kappa: (DX)^\Gamma, \tau: (DY)^X \vdash$$

$$\oint \mu(\delta \mu) (\oint (\kappa r) \tau) = \oint (\oint \mu \kappa) (\delta \kappa) \tau(x)$$

Corollary:  $(P, \delta_-, \oint)$  is a monad.

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→ type judgment

$\mu : DX, E : BX \vdash c_{\mu}[E] : W$

↳ measure  $\mu$  assigns to  $E$

Plan:

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# Language of distribution & Probability Recap

$X$  type (=space) of values / outcomes

$DX$  type of distributions / measures over  $X$

$PX \subseteq DX$  Sub type of probability measures (total measure = 1)

$BX$  type of measurable events - subsets of  $X$  we wish to measure

$W$  type of weights :  $[0, \infty]$

→ type judgment

$\mu : DX, E : BX \vdash c_{\mu}[E] : W$

↳ measure  $\mu$  assigns to  $E$

# Axioms for measures/distributions

Recap

$$\mu: \mathcal{D}X \vdash C_e[\emptyset] = 0 \quad : \mathbb{W}$$

$$E, C : \mathcal{B}X, \mu: \mathcal{D}X \vdash$$

$$C_e[E] = C_e[E \cap C] + C_e[E \cap C^c] \quad : \mathbb{W}$$

$$E_- : (\mathcal{B}X, \subseteq)^\omega, \mu: \mathcal{D}X \vdash$$

$$C_e[\bigcup_n E_n] = \sup_n C_e[E_n] \quad : \mathbb{W}$$



# Kernels & their Kock integral

Recap

Kernel from  $\Gamma$  to  $X$ :  $k: (DX)^\Gamma$  or  $k: \Gamma \rightarrow DX$

Dirac kernel:  $\delta_-: X \rightarrow DX$

Kock integral:  $\mu: D\Gamma, k: (DX)^\Gamma \vdash \int \mu k : DX$   
or  $\int \mu(dx) k(x)$  ( $dx$  binding occurs in  $k(x)$ )

Giry monads:  $(D, \delta_-, \int)$  &  $(P, \delta_-, \int)$ .

# Discrete model

Recap

$$\text{type : set} \quad W := [0, \infty] \quad \mathcal{B}_X := \mathcal{P}X$$

$$DX := \{ \mu : X \rightarrow W \mid \text{supp } \mu \text{ countable} \}$$

$$\mathcal{P}X := \{ \mu \in DX \mid \sum_{\mu} C_{\mu}[X] = 1 \}$$

$$C_{\mu}[E] := \sum_{x \in E} \mu x \quad \delta_x := \lambda x'. \begin{cases} x = x' : 0 \\ x \neq x' : 1 \end{cases}$$

$$\oint \mu k := \lambda x. \sum_{r \in \Gamma} \mu r \cdot k(r; x)$$

## Ex distributions

Recap

Counting measure ( $\chi_{\text{ctbl}}$ ):  $\#_X := \lambda x. 1$

Dirac measure  $\delta_x$  (prev slide)

Zero measure  $\underline{0} := \lambda x. 0$

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# Product measures

$$\mu: D_X, \nu: D_Y \vdash \mu \otimes \nu := \int \mu(x) \int \nu(y) \delta_{(x,y)} : D(X \times Y)$$

( $\otimes$ ) lifts along  $P \hookrightarrow D$

$$= \lambda(x,y). \mu x \cdot \nu y$$

↑ discrete model

$$\text{Ex: } \#_{X \times Y} = \#_X \otimes \#_Y$$

build measures compositionally

Indeed:

$$(\# \otimes \#)(x,y) = \#x \cdot \#y = 1 \cdot 1 = 1 = \#(x,y)$$

Notation:  $\lambda: D(X \times Y), \kappa: (DZ)^{X \times Y} \vdash \iint \lambda(dz, dy) \kappa(z, y)$   
 $= \oint \lambda \kappa$

Fubini - Tonelli Thm:

Integrate in any order:

$\mu: DX, \nu: DY, \kappa: (DZ)^{X \times Y} \vdash$

$$\oint \mu(dx) \oint \nu(dy) \kappa(x, y) = \iint (\mu \otimes \nu)(dz, dy)$$

$$= \oint \nu(dy) \oint \mu(dx) \kappa(x, y)$$

# Pushing a measure forward

$$\mu: D_\Omega, \alpha: X^\Omega, \mu_f := \int \mu(d\omega) \delta_{\alpha\omega} : DX$$

$$= \lambda x. \sum_{\substack{\omega \in \Omega \\ \alpha\omega = x}} \mu \omega$$

$\alpha: X^\Omega$ : random element

(w.r.t.  $\mu$ )

$\mu_\alpha: DX$ : the law of  $\alpha$

Ex: We can represent configurations of 2 dice  
using  $\underline{6} \times \underline{6}$

Letting  $(+)$ :  $\underline{6}^2 \rightarrow \mathbb{N}^2 \xrightarrow{(+)} \mathbb{N}$

We have that the law of  $(+)$ :

$(\#_{\underline{6}} \otimes \#_{\underline{6}})_{(+)} : \mathbb{D}\mathbb{N}$

is the number of  
rolls whose sum is given

build measures  
compositionally



# Scaling a measure

$$(\cdot) : W \times D_X \longrightarrow D_X$$

$$a \cdot \mu := \lambda x. a \cdot \mu x$$

$$\text{NB: } \text{supp}(a \cdot \mu) = \begin{cases} a=0: \emptyset \\ a \neq 0: \text{supp } \mu \\ \quad \checkmark c+|b| \end{cases}$$

$(\cdot) : W \times D_X \rightarrow D_X$  is an action of monoid  $(W, (\cdot), 1)$  on  $D_X$ :

$$\mu : D_X \vdash$$

$$1 \cdot \mu = \mu$$

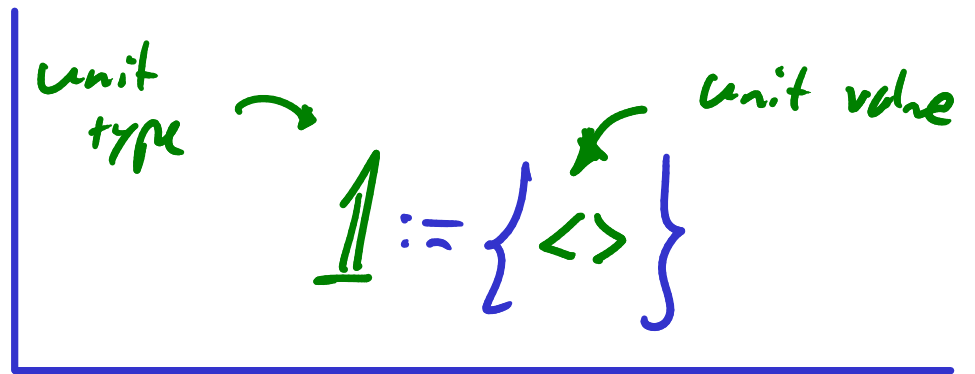
$$a, b : W, \mu : D_X \vdash$$

$$a \cdot (b \cdot \mu) = (a \cdot b) \cdot \mu$$

# Normalisation

$$\mu: D_X, \quad c_e[X] \neq 0, \infty \vdash$$

$$\|\mu\| := \left( \frac{1}{c_e[X]} \right) \cdot \mu \quad : \text{PX}$$



Ex:

$$\emptyset \neq A \subseteq_{\text{fin}} X \quad : \quad \bigcup_{A \subseteq X} \|\#_A\| \quad : \text{PX}$$

$$\mathbb{1} \xrightarrow{\#_A} DA \xrightarrow{(-)_{A \subseteq X}} DX \xrightarrow{\|\cdot\|} \text{PX}$$

$$\text{I.e. } \bigcup_{A \subseteq X} \|\cdot\| := \lambda x. \begin{cases} x \in A: & \frac{1}{|A|} \\ x \notin A: & 0 \end{cases} \quad \text{so } \bigcup_{\{x\} \subseteq X} \|\cdot\| = \delta_x$$

# Standard vocabulary

Joint distributions:  $\mu : D(X_1 \times X_2)$

Marginal distribution:  $X_1 \xleftarrow{\pi_1} X_1 \times X_2 \xrightarrow{\pi_2} X_2$   
law of Projection

$$\mu_{\pi_i} : D X_i$$

marginalisation:  $\mu_{\pi_i} = \int \mu(dx, dy) \delta_x$   
integrate out  $y$

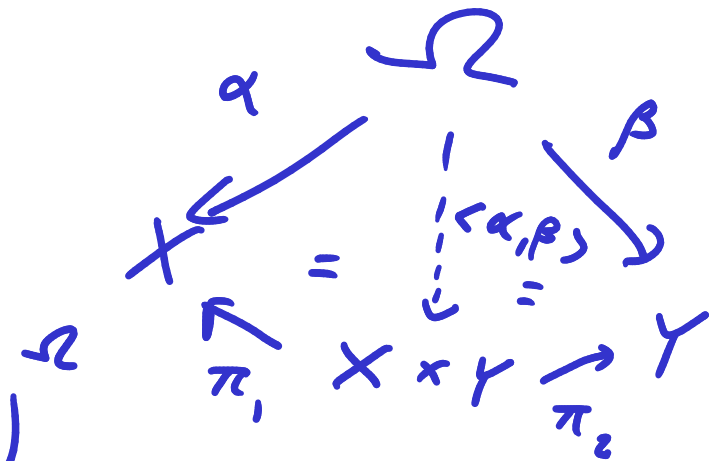
Exercise:  $\mu : P X, \nu : D X \vdash (\mu \otimes \nu)_{\pi_2} = \nu$

# independence

Pairing r.e.s:

$$\alpha : X^\Omega, \beta : Y^\Omega \vdash$$

$$\langle \alpha, \beta \rangle := \lambda \omega. \langle \alpha \omega, \beta \omega \rangle : (X * Y)^\Omega$$



$$\lambda : D^\Omega, \alpha : X^\Omega, \beta : Y^\Omega \vdash \alpha \perp_{\lambda} \beta := \lambda \langle \alpha, \beta \rangle = \lambda \alpha \oplus \lambda \beta$$

: Prop

$\alpha, \beta$  independent w.r.t.  $\lambda$

<sup>(Durmett)</sup>  
Ex represent outcomes of 3 coin tosses:

$$C := \{T, H\} \quad \Omega := C \times C \times C \quad \lambda: \bigcup_C \otimes \bigcup_C \otimes \bigcup_C : P_\Omega$$

$\pi_i: \Omega \rightarrow C$  outcome of  $i^{\text{th}}$  toss

$$\text{Same}_{ij}: \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{(\stackrel{?}{=})} B$$

where:

$$(\stackrel{?}{=}) : C^2 \rightarrow B := \{ \text{True}, \text{False} \}$$
$$x \stackrel{?}{=} y := \begin{cases} x = y : \text{True} \\ x \neq y : \text{False} \end{cases}$$

(Durmett)  
Ex represent outcomes of 3 coin tosses:

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$\pi_i: \Omega \rightarrow C$  outcome of  $i^{\text{th}}$  toss

$$\text{Same}_{ij}: \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{(\cdot)} \mathbb{B}$$

marginalisation

$$\lambda_{\text{Same}_{12}}^T = (\bigcup_C \otimes \bigcup_C)^T_{(\cdot)} = \frac{1}{4} + \frac{1}{4} = \frac{1}{2}$$

$\begin{matrix} \downarrow & & \downarrow \\ \bigcup_C(T) \cdot \bigcup_C(T) & & \bigcup_C(H) \cdot \bigcup_C(H) \end{matrix}$

So  $\lambda_{\text{Same}_{12}}^F = \frac{1}{2}$  too

(Durmett)  
Ex represent outcomes of 3 coin tosses:

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$\pi_i: \Omega \rightarrow C$  outcome of  $i^{\text{th}}$  toss

$i \neq j$ :  $\lambda_{\text{same}_{ij}} = \bigcup_{\mathcal{B}} \mathcal{B}$

$\text{Same}_{ij}: \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{(\cdot)} \mathcal{B}$

$$\lambda: \begin{aligned} (T, T) &\mapsto \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} + \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4} \\ \langle \text{same}_{12}, \text{same}_{23} \rangle &\quad \hookrightarrow \lambda(H, H, H) \quad \hookrightarrow \lambda(T, T, T) \end{aligned}$$

$$\begin{aligned} (T, F) &\mapsto \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} + \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4} \\ &\quad \hookrightarrow \lambda(H, H, T) \quad \hookrightarrow \lambda(T, T, H) \end{aligned}$$

(Durmett)  
Ex represent outcomes of 3 coin tosses:

$$C := \{T, H\} \quad \Omega := C \times C \times C \quad \lambda: \bigcup_C \otimes \bigcup_C \otimes \bigcup_C : P_\Omega$$

$\pi_i: \Omega \rightarrow C$  outcome of  $i^{\text{th}}$  toss

$i \neq j$ :  $\lambda_{\text{same}_{ij}} = \bigcup_{\mathbb{B}}$

$\text{same}_{ij}: \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{(\cdot)} \mathbb{B}$

$$\lambda_{\langle \text{same}_{12}, \text{same}_{23} \rangle} = \bigcup_{\mathbb{B} \times \mathbb{B}} = \bigcup_{\mathbb{B}} \otimes \bigcup_{\mathbb{B}} = \lambda_{\text{same}_{12}} \otimes \lambda_{\text{same}_{13}}$$

So  $\text{same}_{12} \perp_{\lambda} \text{same}_{13}$



# independence

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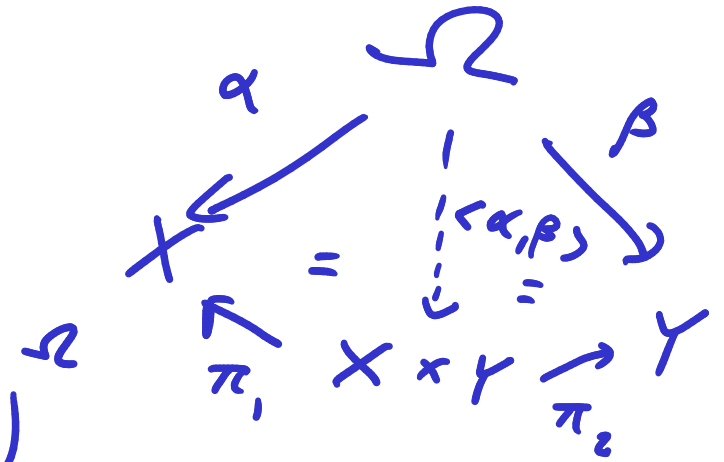
$$\lambda : D\Omega, \alpha : X^\Omega, \beta : Y^\Omega \vdash$$

$$\alpha \perp \beta :=$$

$$\lambda_{\langle \alpha, \beta \rangle} = \lambda_\alpha \otimes \lambda_\beta$$

$\alpha, \beta$  independent wr.t.  $\lambda$

: Prop



I-any version:

$$\lambda : D\Omega, \alpha_i : \prod_{i \in I} X_i^\Omega \vdash \prod_{i \in I} \lambda_{\alpha_i} :=$$

$\alpha_i$  independent  
wr.t.  $\lambda$

$$\forall J \subseteq_{\text{fin}} I.$$

$$\lambda_{\langle \alpha_j \rangle_{j \in J}} = \bigotimes_{j \in J} \lambda_{\alpha_j} : \text{Prop}$$

(Durmett)  
Ex represent outcomes of 3 coin tosses:

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$\pi_i: \Omega \rightarrow C$  outcome of  $i^{\text{th}}$  toss

$$\underline{i \neq j}: \lambda_{\text{Same}_{ij}} = \bigvee_{\mathbb{B}} \quad \text{Same}_{ij}: \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{(\cdot)} \mathbb{B}$$

$$\underline{i \neq j}: \text{Same}_{ij} \perp \text{Same}_{jk} \quad \perp_{\lambda} \{ \text{Same}_{12}, \text{Same}_{23}, \text{Same}_{13} \}$$

Intuition:  $\text{Same}_{13} = \text{IFF} (\text{Same}_{12}, \text{Same}_{23})$

Calc:

$$\lambda (T, T, T)_{\langle \text{Same}_{12}, \text{Same}_{23}, \text{Same}_{13} \rangle} = \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} + \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4} \neq \frac{1}{2^3} = \lambda_{\text{Same}_{12}} \otimes \lambda_{\text{Same}_{23}} \otimes \lambda_{\text{Same}_{13}}$$

$\hookrightarrow \lambda(H, H, H) \quad \hookrightarrow \lambda(T, T, T)$

# Vocabulary

(Discrete) Measure space  $(X, \mu: DX)$

measure preserving  $f: (X, \mu) \rightarrow (Y, \nu)$

function  $f: X \rightarrow Y$  s.t.  $\mu_f = \nu$

$\mu: DX, f: X \rightarrow Y$   $\mu$  invariant under  $f :=$

$f: (X, \mu) \rightarrow (X, \mu)$

Ex:

$\mu: DX, \nu: DY$

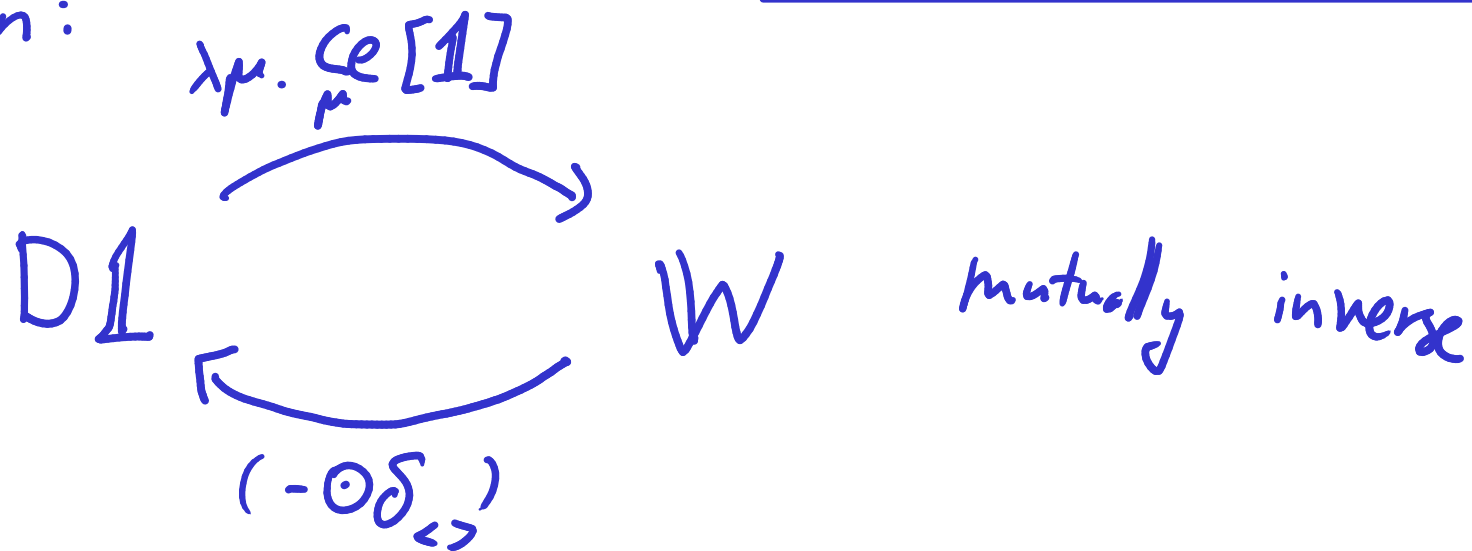
Swap:  $(X \times Y, \mu \otimes \nu) \rightarrow (Y \times X, \nu \otimes \mu)$  so

$\mu: DX$   $\mu \otimes \mu$  invariant under swap

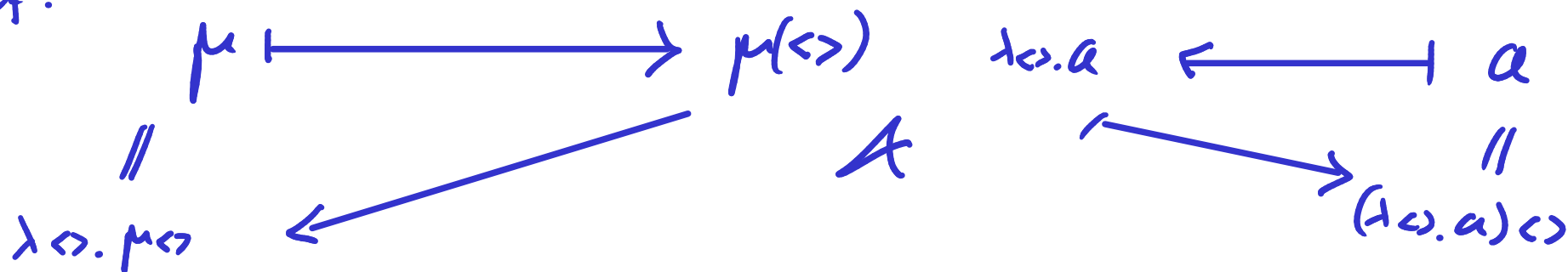
# Weights as measures

NB: unit type  $\rightarrow$   $\mathbb{1} := \{\langle \rangle\}$  unit value

Observation:



Proof:



# Integration

$$\mu: DX, \varphi: W^X \mapsto \int \mu \varphi : W \quad (\text{Lebesgue integral})$$
$$:= \sum_{x \in X} \mu x \cdot \varphi x$$

Can derive it:

$$\begin{array}{ccc} DX \times W^X & \xrightarrow{DX \times (\cong \circ -)} & DX \times (D\mathbb{1})^X \\ \int \downarrow & \cong & \downarrow \phi \\ W & \xleftarrow{\cong} & D\mathbb{1} \end{array}$$

## Additivity:

$$\begin{aligned} I \text{ ctbl}, \mu_- : (DX)^I \vdash \sum_{i \in I} \mu_i & : DX \\ & := \lambda x. \sum_{i \in I} \mu_i x \end{aligned}$$

NB:

$$\begin{aligned} \text{supp} \sum_i \mu_i & \subseteq \\ & \cup_i \text{supp} \mu_i \\ & \checkmark \text{ctbl} \end{aligned}$$

Ex: Bernoulli distribution

$$p : [0,1] \vdash B(p) := p \cdot \underset{\text{True}}{\delta} + (1-p) \cdot \underset{\text{False}}{\delta} : P/B$$

$$\text{i.e. } B_p : \begin{aligned} \text{True} & \mapsto p \\ \text{False} & \mapsto 1-p \end{aligned}$$

Thm (affine-linearity):

$\phi$  is affine-linear in each argument:

I ctbl

$$\mu_- : (D\Gamma)^I, \kappa_- : (Dx)^I, \vdash \int (\sum_{i \in I} a_i \cdot \kappa_i) \mu_- = \sum_{i \in I} a_i \cdot \int \mu_- \kappa_i$$

$a_- : W^I$

I ctbl,  $\mu : D\Gamma$ ,  $a_- : W^I$ ,  $\kappa_- : Dx^I, \vdash$

$$\int \mu(dx) \left( \sum_{i \in I} a_i \cdot \kappa_i(x) \right) = \sum_{i \in I} a_i \cdot \int \mu \kappa_i$$

Prop:  $\mathbb{W} \cong D\mathbb{1}$  is a  $\sigma$ -semi-ring isomorphism:

$$(\mathbb{W}, \Sigma, (\cdot), \mathbb{1}) \cong (D\mathbb{1}, \Sigma', (\cdot), \delta_{\langle \rangle})$$

and  $(\cdot): \mathbb{W} \times D\mathbb{X} \rightarrow D\mathbb{X}$  makes  $D\mathbb{X}$  into a module:

$$\left( \sum_{i \in I} a_i \right) \cdot \mu = \sum_{i \in I} (a_i \cdot \mu) \quad a \cdot \sum_{i \in I} \mu_i = \sum_{i \in I} a \cdot \mu_i$$

Corollary:  $\int$  is affine-linear in each argument.



## Random variable :

NB:  $\bar{\mathbb{R}} := [-\infty, \infty]$

A random element  $\alpha: \bar{\mathbb{R}}^\Omega$  (wrt some  $\mu: D \rightarrow \Omega$ )

Can add, multiply r.v.'s.

To integrate r.v.'s:

$$(-)^{\pm}: \bar{\mathbb{R}}^\Omega \rightarrow \mathbb{W}^\Omega$$

$$\alpha^+ := \lambda \omega. \begin{cases} \alpha \cdot \omega \geq 0: \alpha \omega \\ 0.w: 0 \end{cases} = [\alpha \geq 0] \cdot |\alpha|$$

$$\alpha^- := \lambda \omega. \begin{cases} \alpha \cdot \omega \leq 0: |\alpha \omega| \\ 0.w: 0 \end{cases} = [\alpha \leq 0] \cdot |\alpha|$$

So  $\alpha = \alpha^+ - \alpha^-$

$\mu: D\Omega, \alpha: \bar{\mathbb{R}}^{\Omega}, \int \mu \alpha^+ < \infty$  or  $\int \mu \alpha^- < \infty$  +

$$\int \mu \alpha := \int \mu \alpha^+ - \int \mu \alpha^- : \bar{\mathbb{R}}$$

Ex. The (discrete) Lebesgue  $p$ -space:

$$p: [1, \infty), \mu: P\Omega \vdash \mathcal{L}_p(\Omega, \mu) :=$$

$$\left\{ \alpha: \bar{\mathbb{R}}^{\Omega} \mid \int_{\mu} |\alpha|^p < \infty \right\}$$

$\mathcal{L}_p(\Omega, \mu)$  has a norm  $\|\alpha\| := \sqrt[p]{\int_{\mu} |\alpha|^p}$  almost Banach

$\mathcal{L}_2(\Omega, \mu)$  has an inner product  $\langle \alpha, \beta \rangle := \int_{\mu} \alpha \cdot \beta$  almost Hilbert

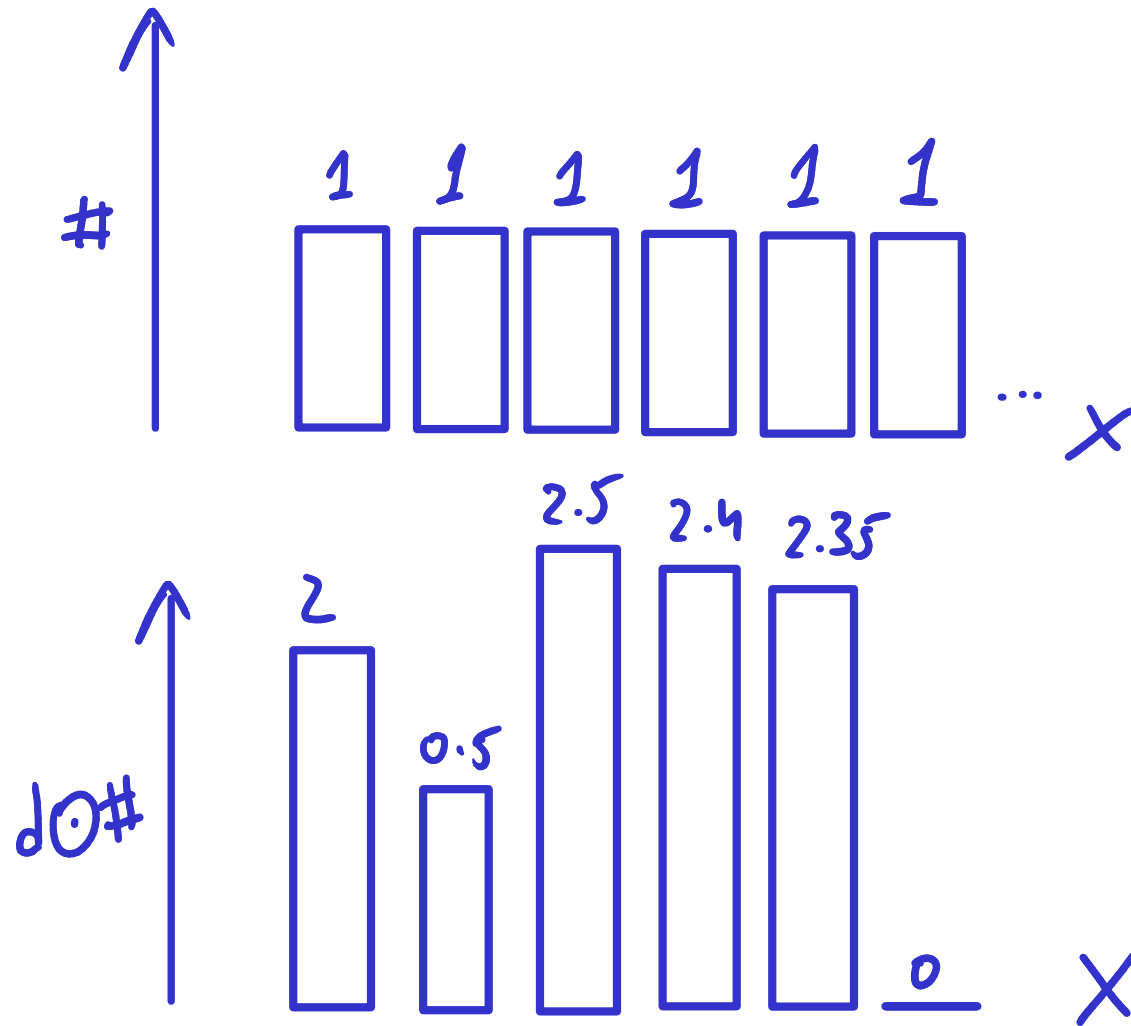
# Density

a density over  $X$  :  $d : X \rightarrow W$

$$d : W^X, \mu : DX \vdash d \odot \mu : DX \\ := \int \mu(dx) (dx \cdot \delta_x)$$

Warning The types of measures & densities in the discrete model are close, but still different. They coincide on countable sets, so people often confuse them. Types help us keep them separate.

Intuition:



# Almost certain properties

$E: \mathcal{B}X, \mu: DX \vdash \mu(dx)$ -almost certainly  $x \in E$  : Prop

$$:= [- \in E] \odot \mu = \mu$$

$$\uparrow \text{NB: } [- \in E] = \begin{cases} x \in E : 1 \\ x \notin E : 0 \end{cases} : W$$

When  $\mu: P X$  we say instead

$\mu(dx)$ -almost surely  $x \in E$

Exercise Look up the def. of a normed space

and modify the definition so that  $L_p(\Omega, \mu)$  is a normed space up to almost sure equality.

## Absolute continuity

$d$  is a density of  $\mu$  w.r.t.  $\nu$  or

$d$  is a Radon-Nikodym derivative w.r.t.  $\nu$

$$\mu, \nu: \mathcal{D}X, d: \mathcal{W}^X \vdash d = \frac{d\mu}{d\nu} \quad : \text{Prop}$$

$$:= \mu = d \circ \nu$$

$\mu, \nu: \mathcal{D}X \vdash \mu \ll \nu := \mu$  is absolutely continuous w.r.t.  $\nu$  : Prop

$$:= \exists d: \mathcal{W}^X. d = \frac{d\mu}{d\nu}.$$

$:= \mu$  has a density w.r.t.  $\nu$

Lemma:  $\mu, \nu: \mathcal{D}X,$   
 $\mu \ll \nu,$   
 $h: (\mathcal{D}X)^X$

$$\int \nu(dx) \frac{d\mu}{d\nu}(x) \cdot kx = \int \mu(dx) kx$$

$$\underline{\text{Ex:}} \quad U_{A \subseteq X} \ll (\#_A)_{\text{Cost: } A \subseteq X}$$

$$\frac{dU_{A \subseteq X}}{d(\#_A)_{\text{Cost}}} = \lambda x. \left\{ \begin{array}{l} x \in A: \frac{1}{|A|} \\ \text{O.W.}: 0 \end{array} \right.$$

but also:

$$\frac{dU_{A \subseteq X}}{d(\#_A)_{\text{Cost}}} = \lambda x. \frac{1}{|A|}$$

Radon-Nikodym Thm: (discrete version)

$\mu, \nu: \mathcal{P}X \vdash \mu \ll \nu$  iff  $\forall x. \nu x = 0 \Rightarrow \mu x = 0$

i.e.  $\text{Supp } \mu \subseteq \text{Supp } \nu$

In that case, if  $d_1, d_2 = \frac{d\mu}{d\nu}$  then

$\nu(dx)$ -a.s.  $d_1 x = d_2 x$

Ex: for ctbl  $X$ ,  $\forall \mu: \mathcal{D}X. \mu \ll \#_X$ . Proof: vacuously, as  $\#_X x \neq 0$ .

Then  $\lambda x. \mu x = \frac{d\mu}{d\#}$ .



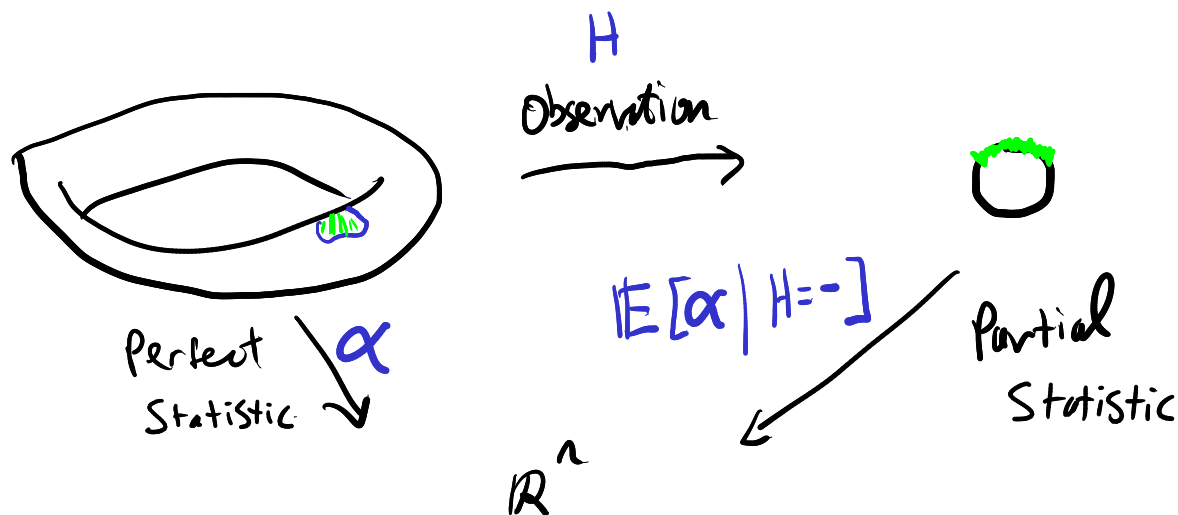
# Conditional expectation

$\beta$  is a conditional expectation of  $\alpha$  w.r.t.  $\mu$  along  $H$

$$\mu: \mathcal{D}\Omega, H: X^\Omega, \alpha: \mathcal{L}_1(\Omega, \mu), \beta: \mathcal{L}_1(X, \mu_H)$$

$$\vdash \beta = \mathbb{E}[\alpha | H = -] \quad : \text{Prop}$$

$$:= \forall \varphi: \mathcal{L}_1(X, \mu_H^M). \int \mu_H(d\alpha) \beta(\alpha) \cdot \varphi(\alpha) = \int \mu(d\omega) \alpha(\omega) \cdot \varphi(H\omega)$$



Thm (Kolmogorov): (discrete version)

There is a function

$$\underline{\mathbb{E}}[-|-] \in \prod_{\mu: P_{\Omega}} \prod_{H: X^{\Omega}} L_1(\Omega, \mu) \rightarrow L_1(X, \mu_H)$$

s.t.  $\mathbb{E}_{\mu}[\alpha | H = -]$  is a conditional expectation of  $\alpha$  w.r.t.  $\mu$  along  $H$ .

# Conditional Probability (discrete version):

$$H: X^\Omega, \mu: P_X \vdash P_\Gamma[- | H = -] : (P_\Omega)^X$$

$$:= \lambda x_0: X. \lambda \omega_0: \Omega. \mathbb{E}_{\omega \sim \mu} [ [ \omega_0 = \omega ] | H\omega = x_0 ]$$

# Bayes's Theorem (discrete version, adapted from Williams):

Let  $\lambda: P(X \times \mathcal{H})$  joint probability distribution.

Assume  $\mu: D_X, \nu: D_{\mathcal{H}}$  s.t.  $\lambda \ll \mu \otimes \nu$ .

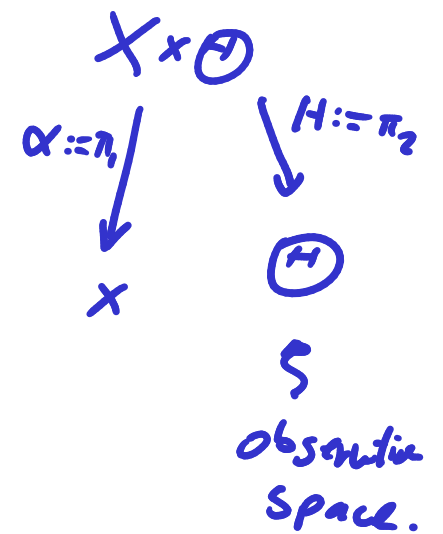
with  $d_{X, \mathcal{H}} = \frac{d\lambda}{d(\mu \otimes \nu)}$ .

obs 1:  $d_X: \mathcal{W}^X$

$$d_X := \lambda_{\mathcal{H}} \int \nu(d\theta) d_{X, \mathcal{H}}(x, \theta)$$

then  $d_X = \frac{d\lambda_{\mathcal{H}}}{d\mu}$

A similar  $(d_{\mathcal{H}}: \mathcal{W}^{\mathcal{H}}) := \lambda_X \int \mu(dx) d_{X, \mathcal{H}}(x, \theta) = \frac{d\lambda_X}{d\nu}$

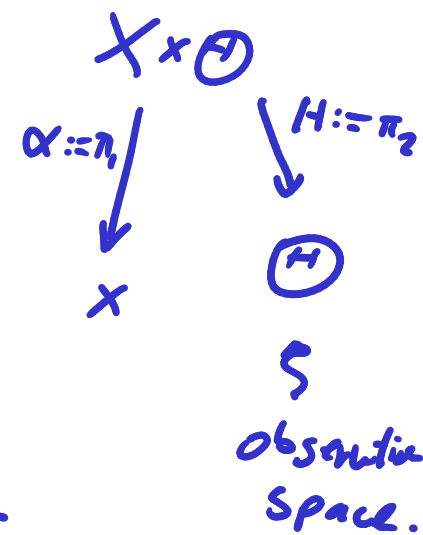


# Bayes's Thm (discrete version, adapted from Williams):

Let  $\lambda: P(X \times \Theta)$  joint probability distribution.

Assume  $\mu: D_X, \nu: D_\Theta$  s.t.  $\lambda \ll \mu \otimes \nu$ .

with  $d_{X,H} = \frac{d\lambda}{d(\mu \otimes \nu)}$  .  $d_X = \frac{d\lambda_\alpha}{d\mu}$       $d_\Theta = \frac{d\lambda_H}{d\nu}$



Let  $d_{X|H}(-|-): X \times \Theta \rightarrow W$

$$d_{X|H}(\alpha|\theta) := \begin{cases} d_\theta \neq 0: & \frac{d_{X,H}(\alpha,\theta)}{d_\theta} \\ \text{o.w.:} & 0 \end{cases}$$

$$\lambda_{X|H=-} : \Theta \rightarrow P_X$$

$$\lambda_{X|H=\theta} := d_{X|H}(-|\theta) \otimes \mu$$

Bayes's formula:

$$P_\lambda[-|H=-] = \lambda_{X|H=-}$$

# Summary

$\mu \otimes \nu$  Product measures & Fubini-Tonelli

$\mu_H$  Push-forward / law

$(\mathcal{D}_X, \Sigma, (\cdot))$  module structure over affine linearity of  $\mathcal{F}$

} Lebesgue integration

Standard vocabulary: joint dist., marginalisation, independence, invariance

density & Radon-Nikodym derivatives (heed the **warning**)

almost certain properties

Conditional expectation & Probability  
with Bayes's Thm.

Plan:

- 1) Type-driven probability: discrete case (Mon + Tue) ✓
- 2) Borel sets & measurable spaces (Wed) ✓
- 3) Quasi Borel spaces, simple type structure (Wed)
- 4) Dependent type structure & standard Borel spaces (Thu)
- 5) Integration & random variables (Fri)

please ask questions!



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# Foundations for type-driven probabilistic modelling

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Plan:

- 1) Type-driven probability: discrete case (Mon + Tue)
- 2) Borel sets & measurable spaces (Wed)
- 3) Quasi Borel spaces (Web) Simple type structure (Thu)
- 4) Dependent type structure & standard Borel spaces (Thu)
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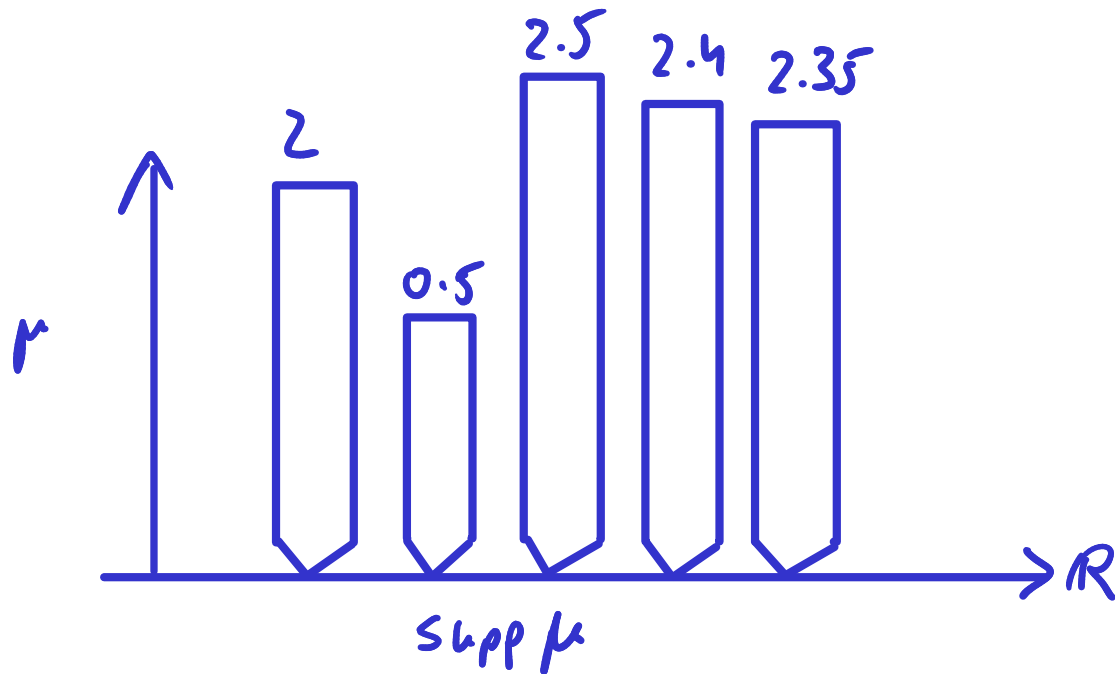
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discrete model measure only histograms:



Want:

- lengths
- areas
- volumes.

Continuous **Caveat:**

Thm: No  $\lambda: \mathcal{P}\mathbb{R} \rightarrow [0, \infty]$ :

$$\lambda(a, b) = b - a$$

(generalises length)

$$\lambda(r + A) = \lambda A$$

(translation invariant)

$$\lambda\left(\bigcup_{n=0}^{\infty} A_n\right) = \sum_{n=0}^{\infty} \lambda A_n$$

$\sigma$ -additive

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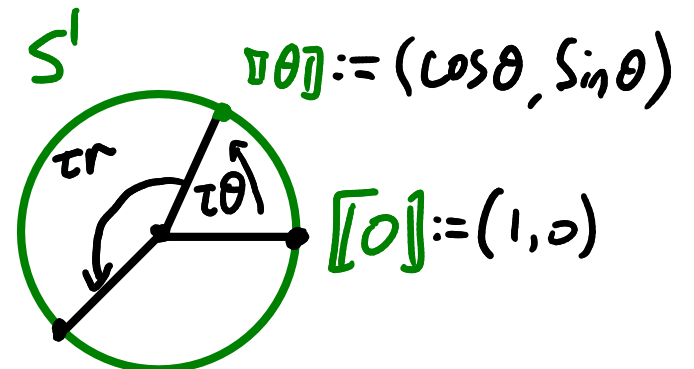
(translation invariant)

$$\lambda\left(\bigcup_{n=0}^{\infty} A_n\right) = \sum_{n=0}^{\infty} \lambda A_n$$

$\sigma$ -additive

Direct proof in standard analysis courses. Idea behind typical proof is:

Thm: no  $\lambda: \mathcal{P}S^1 \rightarrow [0, \infty]$  st.



$$r: \mathbb{R} \mapsto \text{rotate}_r, [0] := [\theta + \tau r]$$

a) satisfy measure axioms for  $\mathcal{B}S^1 := \mathcal{P}S^1$

b) invariant under rotations:  $E \in \mathcal{B}S^1$

$$\lambda \text{ rotate}[E] = \lambda E$$

c)  $\lambda S^1 = \tau (= 2\pi)$

Reduce  $(S', \lambda^{S'})$  to  $(\mathbb{R}, \lambda^{\mathbb{R}})$  via restriction & push forward

$$\lambda^{\mathbb{R}}|_{\mathcal{P}[0,1]} := \lambda_{E \subseteq [0,1]}. \quad \lambda E : \mathcal{P}[0,1] \rightarrow \mathbb{W}$$

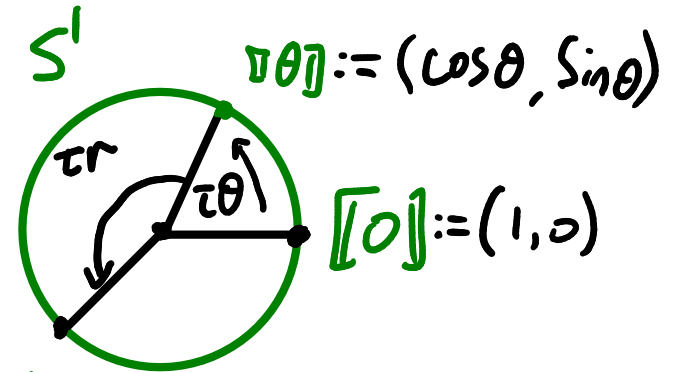
$$\lambda^{S'} := \lambda_{E \subseteq S'}. \quad \lambda^{\mathbb{R}}(\mathbb{I}^{-1}[E]) : \mathcal{P}S' \xrightarrow{\mathbb{I}^{-1}} \mathcal{P}[0,1] \xrightarrow{\lambda^{\mathbb{R}}|_{\mathcal{P}[0,1]}} \mathbb{W}$$

noting

rotations in  $S'$   $\iff$  translations in  $\mathbb{R}$

Since  $\nexists \lambda^{S'}$ , we have  $\nexists \lambda^{\mathbb{R}}$  either.

Thm: no  $\lambda: \mathcal{P}S^1 \rightarrow [0, \infty]$   
 s.t.



a) Satisfy measure axioms for  $BS' := \mathcal{P}S^1$

b) invariant under rotations:  $E: BS' \rightarrow \mathbb{R}$

c)  $\lambda S^1 = \tau$  ( $:= 2\pi$ )

$$\lambda \text{rotate}_r[E] = \lambda E$$

Proof:  $a + b \Rightarrow \neg c$ :

1) Using axiom of choice (AOC):

$$S^1 = \bigoplus_{i=0}^{\infty} E_i$$

$$E_i = \text{rotate}_{r_i}[E_0]$$

$$2) \lambda S^1 = \sum_{i=0}^{\infty} \lambda E_i = \sum_i \lambda \text{rotate}_{r_i} E_0 = \sum_{i=0}^{\infty} \lambda E_0 = \begin{cases} \lambda E_0 = 0: 0 \\ \lambda E_0 > 0: \infty \end{cases} \neq \tau$$

Constructing  $E_i$ :

$$x, y: S' \vdash x \sim y := \exists q \in \mathcal{Q}. \text{rotate}_q x = y \quad : \text{Prop}$$

$$\equiv \exists q \in [0, 1) \cap \mathcal{Q}. \text{rotate}_q x = y$$

$\sim$ -equivalence classes:

$$x: S' \vdash [x]_{\sim} := \{ y \in S' \mid x \sim y \} \quad : \mathcal{P}S'$$

$$C := \{ [x]_{\sim} \in \mathcal{P}S' \mid x \in S' \}$$

$$\forall e \in C, e \neq \emptyset, \text{ so by } AOC: \exists \xi: C \rightarrow S'. \xi_e \in e.$$

---

NB:  $\xi$  injective

Take  $C_0 := \{ \sum_e \in S' \mid e \in C \} \in \mathcal{P} S'$

Note:  $x \sim y, x, y \in C_0 \vdash x = y$ .

$q: \mathbb{Q} \vdash C_q := \text{rotate}_q[C_0] \in \mathcal{P} S'$

Let  $(r_i)_{i=0}^{\infty}$  enumerate  $\mathbb{Q} \cap [0, 1)$  st.  $r_0 = 0$

Take  $E_i := C_{r_i}$

By fiat:  $E_i = C_{r_i} = \text{rotate}_{r_i}[C_0] = \text{rotate}_{r_i}[E_0]$

RTP:  $S' = \bigoplus_{i=0}^{\infty} E_i$

NB:  $x, y: S' \vdash$   
 $x \sim y: \text{Prop}$

$C = \sim$ -equiv.

$\sum: C \rightarrow S'$

$e: C \vdash \sum_e \in e$

$$E_i \cap E_j = \emptyset, \quad i \neq j:$$

---

$$x \in E_1 \cap E_2 \Rightarrow \exists y_i \in \mathcal{C}. \quad x = \text{rotate}_{r_i} y_i$$

$$\Rightarrow y_1 \sim x \sim y_2 \Rightarrow y_1 = y_2 =: y$$

$$\Rightarrow \text{rotate}_{r_2 - r_1} y = y, \quad |r_2 - r_1| < 1$$

$$\Rightarrow r_1 = r_2$$



$S' = \bigcup_{i=0}^{\infty} E_i$  :  $x \in S'$ . letting  $e := \xi_{[x]_n} : \mathcal{P}S'$

$$\xi_e, x \in E \Rightarrow \xi_e \sim x$$

$$\Rightarrow \exists q \in (\mathbb{Q} \cap [0, 1)). \text{ rotate }_q \xi_e = x.$$

As  $\xi_e \in C_0$  :  $x \in C_q$ . Find  $i$  s.t.  $r_i = q$

and  $x \in C_{r_i} = E_i$ .



Takeaway: Taking  $BIR := DIR$

Excludes measures such as:

length, area, volume

Workaround: only measure well-behaved subsets

Def: The Borel subsets  $\mathcal{B}_{\mathbb{R}} \subseteq \mathcal{P}\mathbb{R}$ :

- open intervals  $(a, b) \in \mathcal{B}_{\mathbb{R}}$

closure under  $\sigma$ -algebra operations:

---

$$\emptyset \in \mathcal{B}_{\mathbb{R}}$$

↙  
empty set

---

$$A \in \mathcal{B}_{\mathbb{R}} \quad \overline{A^c} := \mathbb{R} \setminus A \in \mathcal{B}$$

↙  
complements

---

$$\vec{A} \in \mathcal{B}_{\mathbb{R}}^{\mathbb{N}} \quad \overline{\bigcup_{n=0}^{\infty} A_n} \in \mathcal{B}_{\mathbb{R}}$$

↙  
countable unions

# Examples

discrete Countable:  $\{r\} = \bigcap_{\epsilon \in \mathbb{Q}^+} (r-\epsilon, r+\epsilon) \in \mathcal{B}_{\mathbb{R}}$

$I$  countable  $\Rightarrow I = \bigcup_{r \in I} \{r\} \in \mathcal{B}_{\mathbb{R}}$

closed intervals:  $[a, b] = (a, b) \cup \{a, b\}$

Non-examples?

More complicated: analytic, Lebesgue

Def: Measurable space  $V = (V, B_V)$

Set (carrier)  $\checkmark$   
 Family of Subsets  
 $B_V \subseteq P(V)$

closed under  $\sigma$ -algebra operations:

$\emptyset \in B_V$   
 $\uparrow$   
 empty set

$A \in B_V$   


---

 $A^c := V \setminus A \in B_V$   
 $\uparrow$   
 complements

$\vec{A} \in B_V^{\mathbb{N}}$   


---

 $\bigcup_{n=0}^{\infty} A_n \in B_V$   
 $\uparrow$   
countable unions

Idea: structure all spaces after the worst-case scenario

# Examples

- Discrete spaces  $X^{\text{meas}} = (X, P_X)$
- Euclidean spaces  $\mathbb{R}^n$  — replace intervals with chests  $\prod_{i=1}^n (a_i, b_i)$   
 $\mathbb{R}^{\mathbb{N}}$  similarly  $\{C \cap A \mid C \in \mathcal{B}_V\}$
- Sub spaces:  $A \in P_{\mathcal{L}V}$   $A := (A, [B_V] \cap A)$
- Products:  $A \times B := (\mathcal{L}A_1 \times \mathcal{L}B_1, \sigma([B_A] \times [B_B]))$

Def: Borel measurable functions  $f: V_1 \rightarrow V_2$

- functions  $f: V_1 \rightarrow V_2$
- inverse image preserves measurability:

$$f^{-1}[A] \in B_{V_1} \iff A \in B_{V_2}$$

### Examples

- $(+), (\cdot) : \mathbb{R}^2 \rightarrow \mathbb{R}$
- $| \cdot |, \sin : \mathbb{R} \rightarrow \mathbb{R}$
- any continuous function  $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$
- any function  $f: X \rightarrow V$

# Category Meas

Objects: Measurable spaces

Morphisms: Measurable functions

Identities:

$$\text{id} : V \rightarrow V$$

Composition:

$$f : V_2 \rightarrow V_3 \quad g : V_1 \rightarrow V_2$$

---

$$f \circ g : V_1 \rightarrow V_3$$



# Meas Category

Products, Coproducts / disjoint union, Subspaces  
Categorical limits, colimits, but:

Thm [Aumann '61] No  $\sigma$ -algebras  $B_{B_{\mathbb{R}}}, B_{\mathbb{R}^{\mathbb{R}}}$  for measurable

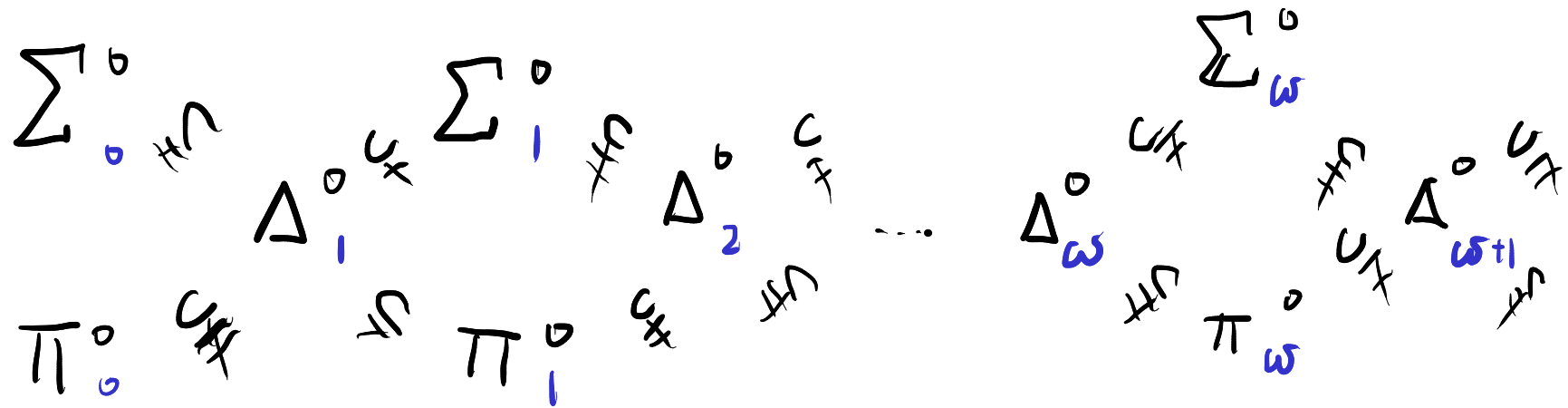
membership predicate  $\leftarrow$   $(\ni) : (B_{\mathbb{R}}, B_{B_{\mathbb{R}}}) \times \mathbb{R} \longrightarrow \text{Bool}$   
 $(U, r) \mapsto [r \in U]$

eval :  $(\text{Meas}(\mathbb{R}, \mathbb{R}), B_{\mathbb{R}^{\mathbb{R}}}) \times \mathbb{R} \rightarrow \mathbb{R}$   
 $(f, r) \mapsto f(r)$

Questions! skip proof?

Proof (sketch):

Borel hierarchy:



Stabilises at  $\Delta_{\omega_1}^0 = \mathcal{B}(\Sigma_0^0) = \Delta_{\omega_1+1}^0$

$$\text{rank } A := \min \{ \alpha < \omega_1 \mid A \in \Delta_\alpha^0 \}$$

then  
for  $B_{B_{\mathbb{R}}} = P(B_{\mathbb{R}})$

$$(\exists) : (B_{\mathbb{R}}, B_{B_{\mathbb{R}}}) \times \mathbb{R} \rightarrow \mathbb{R}$$

$$(U, r) \mapsto [r \in U]$$

If measurable:

$$B_{V \times U} = B([B_V] \times [B_U])$$

$$\alpha := \text{rank}((\exists)^{-1}[\text{true}]) < \omega,$$

Take  $A \in B_{\mathbb{R}}$ ,  $\text{rank} A > \alpha$

But:

$$\alpha < \text{rank} A = \text{rank}(A, \rightarrow)^{-1}[(\exists)^{-1}[\text{true}]] \leq \text{rank}((\exists)^{-1}[\text{true}]) \leq \alpha$$

#

More details in Ex. B

Sequential Higher-order structure:

$$I \text{ Countable} : V^I = \prod_{i \in I} V$$

$\Rightarrow$  Some higher-order structure in Meas:

$$\text{Cauchy} \in B_{[-\infty, \infty]}^{\mathbb{N}}$$

$$\text{Cauchy} = \bigcap_{\epsilon \in \mathbb{Q}^+} \bigcup_{k \in \mathbb{N}} \bigcap_{\substack{m, n \in \mathbb{N} \\ m, n \geq k}} \{ \vec{y} \in [-\infty, \infty]^{\mathbb{N}} \mid |y_m - y_n| < \epsilon \}$$

$$\text{lim sup} : [-\infty, \infty]^{\mathbb{N}} \rightarrow [-\infty, \infty] \quad \text{lim} : \text{Cauchy} \rightarrow \mathbb{R}$$

Compose higher-order building blocks:

lim is measurable!  
↗

$$\text{VanishingSeq}(\mathbb{R}) := \left\{ \vec{r} \in \mathbb{R}^{\mathbb{N}} \mid \lim_{n \rightarrow \infty} r_n = 0 \right\} \in \mathcal{B}_{\mathbb{R}^{\mathbb{N}}}$$

$$\text{approx}_\Delta : \text{VanishingSeq}(\mathbb{R}^+) \times \mathbb{R} \longrightarrow \mathbb{Q}^{\mathbb{N}}$$

$$\text{s.t.} : \left| \left( \text{approx}_{\Delta} \vec{r} \right)_n - r \right| < \Delta_n$$

Slogan: Measurable by Type! ▽

Not all operations of interest fit:

$$\text{lim sup} : ([-\infty, \infty]^{\mathbb{R}})^{\mathbb{N}} \longrightarrow [-\infty, \infty]^{\mathbb{R}}$$

$$\text{lim sup} := \lambda \vec{f}. \lambda x. \limsup_{n \rightarrow \infty} f_n x$$

Intrinsically higher-order! ▽

Want

Slogan: measurability by type!

But

For higher-order building blocks

defer measurability proofs until

we resume 1<sup>st</sup> order fragment  $\Rightarrow$  non compositional

Plan:

- 1) Type-driven probability: discrete case (Mon + Tue)
- 2) Borel sets & measurable spaces (Wed)
- 3) Quasi Borel spaces (Web) Simple type structure (Thu)
- 4) Dependent type structure & standard Borel spaces (Thu)
- 5) Integration & random variables (Fri)

please ask questions!

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# Plan

Def:  $V \in \text{Meas}$  is Standard Borel when

$$V \cong A \quad \text{for some } A \in \mathcal{B}_{\mathbb{R}}$$

the "good part" of  $\text{Meas}$  — the subcategory

$$\text{Sbs} \hookrightarrow \text{Meas}$$



Sbs includes

- Discrete  $\mathbb{I}$ ,  $\mathbb{I}$  countable
- Countable products of Sbs:

$$\mathbb{R}^n, \mathbb{R}^{\mathbb{N}}, \mathbb{Z}^{\mathbb{N}}, \mathbb{N}^{\mathbb{N}}$$

- ~ Borel subspaces of Sbs:

$$\mathbb{I} := [0, 1]$$

$$\mathbb{R}^+ := (0, \infty) \quad \mathbb{R}_{\geq 0} := [0, \infty]$$

- Countable coproducts of Sbs:

$$\mathbb{W} := [0, \infty]$$

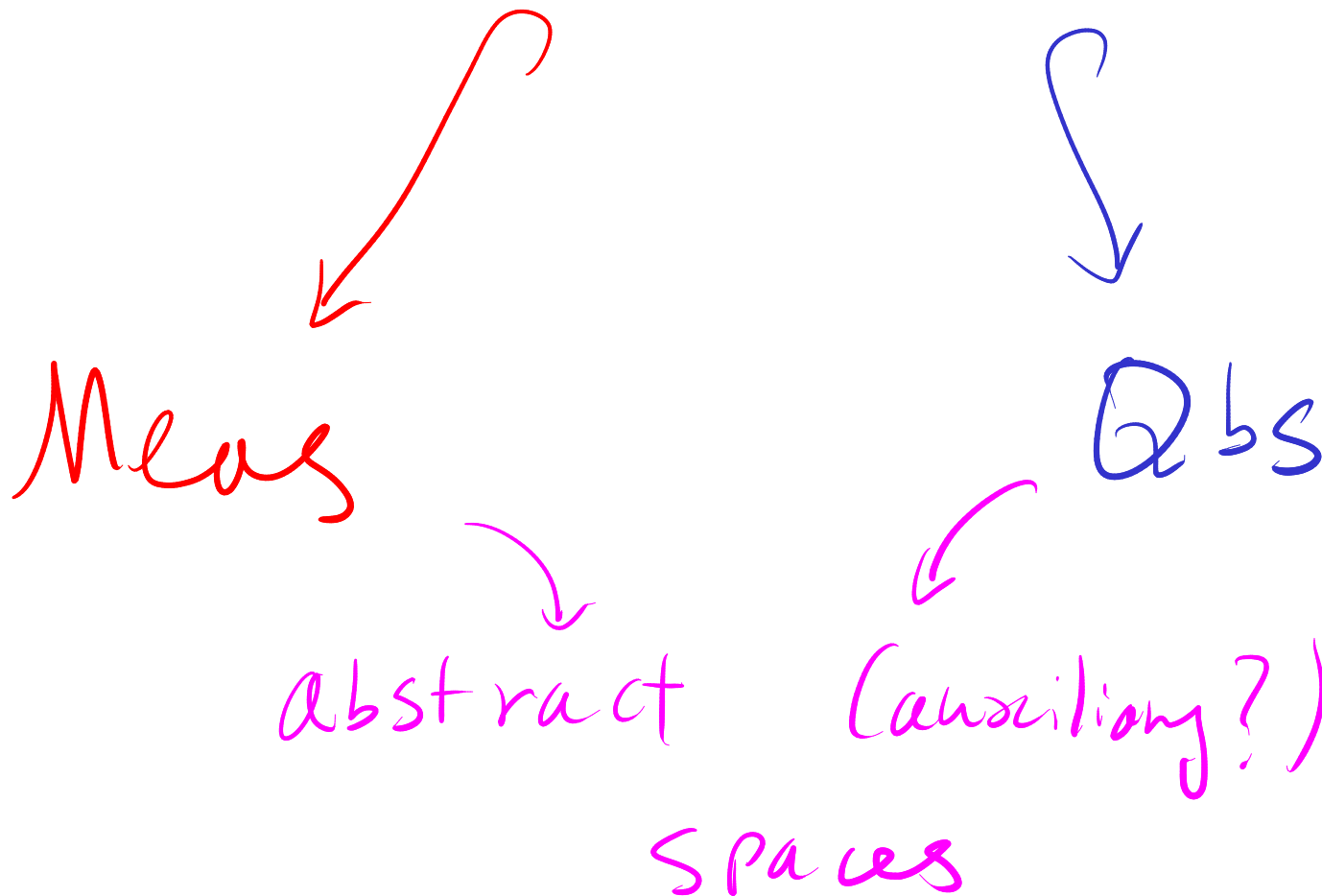
$$\overline{\mathbb{R}} := [-\infty, \infty]$$

Conservative extensions:

Concrete spaces

we "observe"

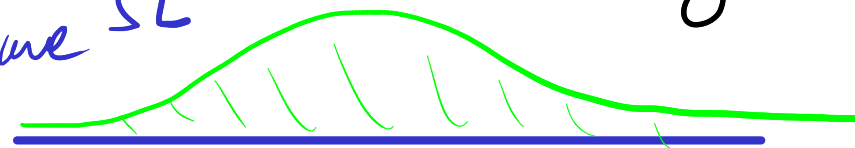
Standard Borel spaces



# Cone idea

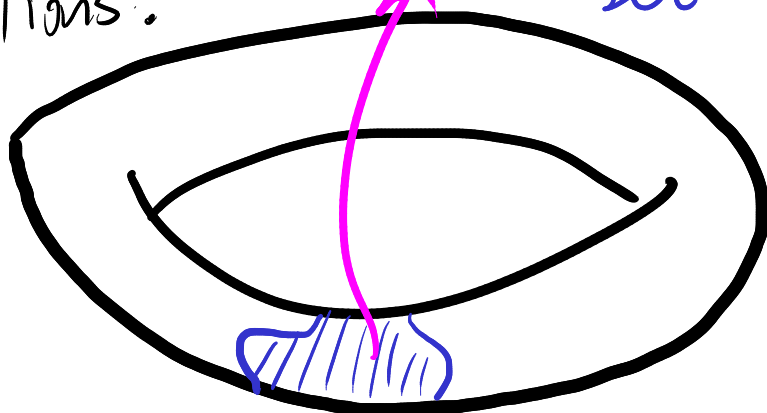
Measure Theory

sample space  $\Omega$  Obs Theory



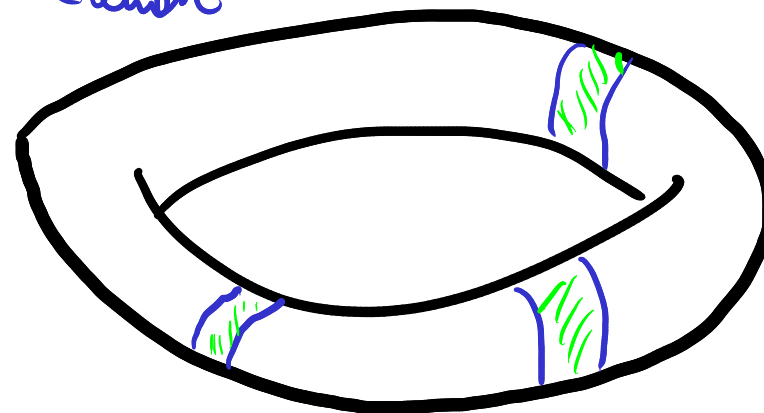
Primitive notions:

measurable subset



random element

$\downarrow \alpha$



Derived

measure

notions:

random

elements

$\alpha: \Omega \rightarrow \text{Space}$

measurable subsets

Def: Quasi-Borel space  $X = (\mathcal{L}X, \mathcal{R}_X)$

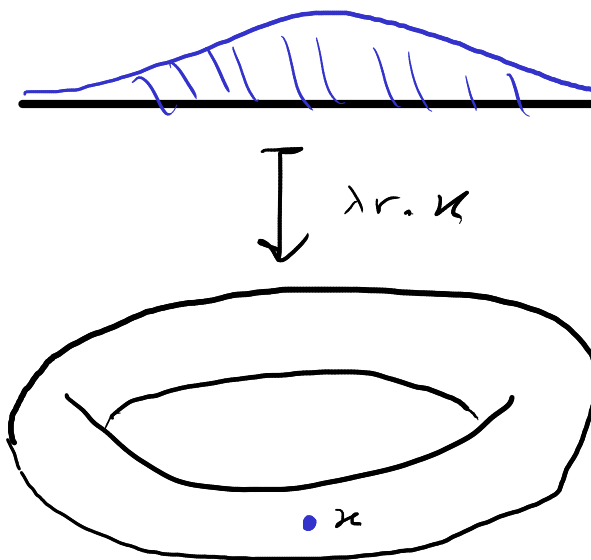
$\mathcal{R}_X \subseteq \mathcal{L}X^{\mathcal{L}\mathbb{R}}$  closed under:

Set  
"carrier"

Set of  
functions  $\alpha: \mathbb{R} \rightarrow \mathcal{L}X$   
"random elements"

- Constant  $S$ :

$$\frac{x \in \mathcal{L}X}{(\lambda r. x) \in \mathcal{R}_X}$$



- Precomposition:

- recombination

Def: Quasi-Borel space

$$X = (\mathcal{L}X, \mathcal{R}_X)$$

Set  
"carrier"

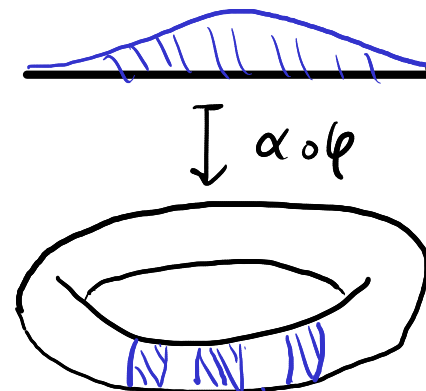
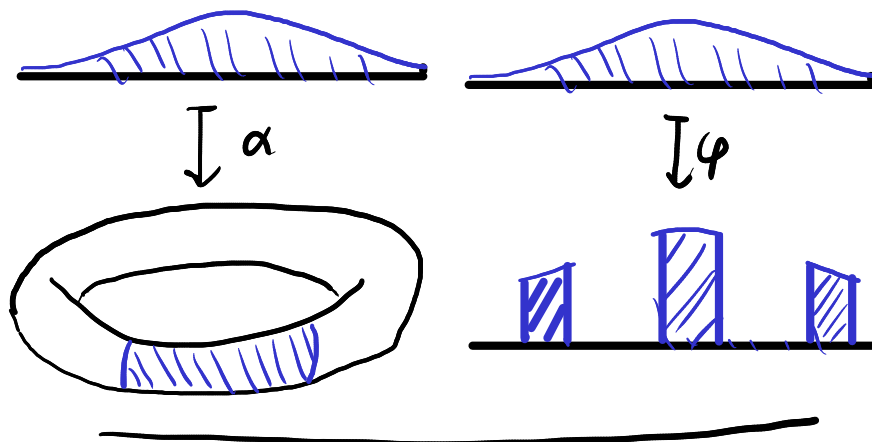
Set of  
functions  $\alpha: \mathbb{R} \rightarrow \mathcal{L}X$   
"random elements"

$\mathcal{R}_X \subseteq \mathcal{L}X^{\mathbb{R}}$  Closed under:

- Precomposition:

$\alpha \in \mathcal{R}_X \quad \varphi: \mathbb{R} \rightarrow \mathbb{R} \text{ in Sbs}$

$$\varphi \circ \alpha: \mathbb{R} \xrightarrow{\varphi} \mathbb{R} \xrightarrow{\alpha} \mathcal{L}X \in \mathcal{R}_X$$



Def: Quasi-Borel space

$$X = (\mathcal{L}X, \mathcal{R}_X)$$

Set  
"carrier"

Set of  
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$$\mathcal{R}_X \subseteq \mathcal{L}X^{\mathbb{L}\mathbb{R}}$$

Closed under:

- re combination

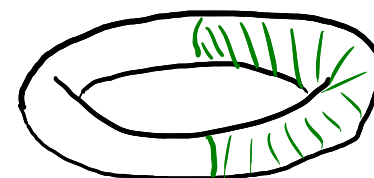
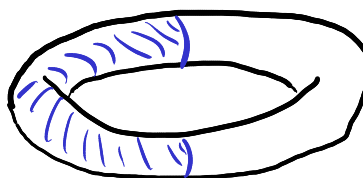
$$\vec{\alpha} \in \mathcal{R}_X^{\mathbb{N}} \quad \mathbb{R} = \bigcup_{n=0}^{\infty} A_n \quad \in \mathcal{B}_{\mathbb{R}}$$

$$\lambda r. \begin{cases} r \in A_n: \alpha_n^r \\ \vdots \end{cases}$$

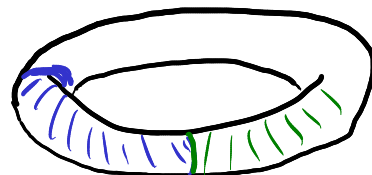


$\downarrow \alpha$

$\downarrow \beta$



$\downarrow \lambda r. \begin{cases} r \in A: \alpha r \\ r \in B: \beta r \end{cases}$



Def: Quasi-Borel space

$$X = (\mathcal{L}X, \mathcal{R}X)$$

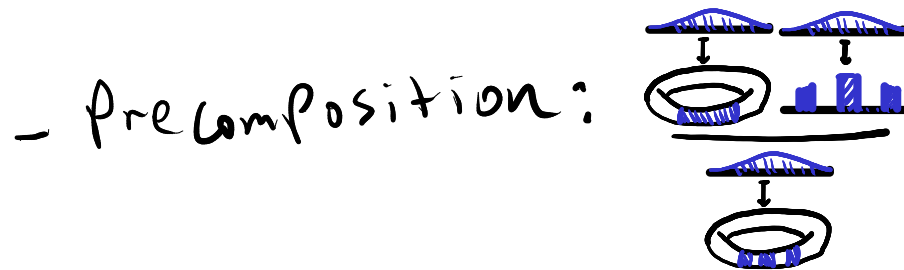
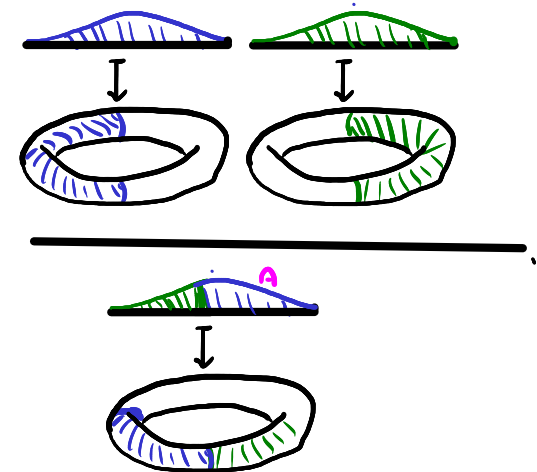
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$\mathcal{R}X \subseteq \mathcal{L}X^{\mathbb{R}}$  Closed under:



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# Examples

recombination of constants

$$- \mathbb{R} = (\mathbb{R}, \text{Meas}(\mathbb{R}, \mathbb{R}))$$

qbs underlying  $\mathbb{R}$

$$- X \in \text{Set}, \quad \overset{\text{qbs}}{X} := (X, \sigma\text{-simple}(\mathbb{R}, X))$$

$\lambda_r.$   $\left\{ \begin{array}{l} \vdots \\ r \in A_n: x_n \\ \vdots \end{array} \right.$

discrete qbs on  $X$

$$- \quad \underset{\text{Qbs}}{\mathbb{R}} X := (X, X^{\mathbb{R}})$$

all functions

Indiscrete qbs on  $X$



Obs morphism  $f: X \rightarrow Y$

- function  $f: X_1 \rightarrow Y_1$

-  $\alpha \downarrow \in R_X$

---

$\alpha \downarrow \in R_Y$   
 $f \downarrow$

Example

- Constant functions

are obs  
morphisms

-  $\sigma$ -simple functions

are obs morphisms

Category Obs  $\Leftarrow$

- identity, composition

# Full model

type : Obs     $w := [0, \infty]$      $B_X := (\text{Thur})$

$D_X := (\text{Fri})$

$P_X := \{ \mu \in D_X \mid C_{\mu}[X] = 1 \}$      $(\text{Thu})$

$C_{\mu}[E] := (\text{Fri})$      $\delta_x := (\text{Fri})$

$\phi_{\mu k} := (\text{Fri})$

# Plan:

- 1) Type-driven probability: discrete case (Mon + Tue)
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Please ask questions!

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# Foundations for type-driven probabilistic modelling

Ohad Kammar  
University of Edinburgh

Logic Summer School  
Australian National University  
4–16 December, 2023  
Canberra, ACT, Australia



THE UNIVERSITY of EDINBURGH  
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## Plan:

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## Full model

type : Obs     $w := [0, \infty]$      $\mathcal{B}_X := (\text{Thur})$

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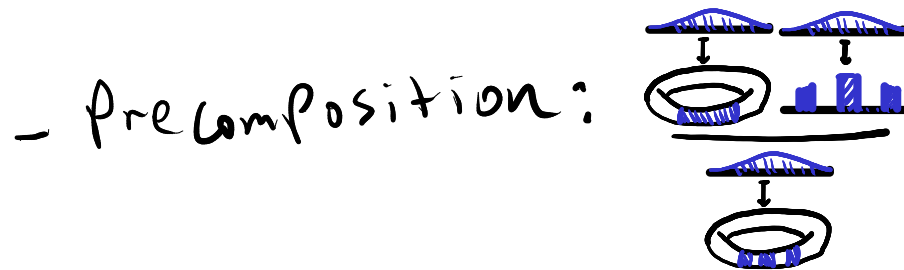
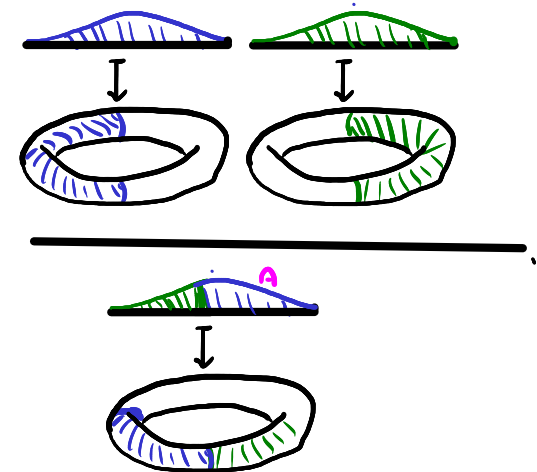
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discrete qbs on  $X$

$$- \quad \underset{\text{Qbs}}{\mathbb{R}} X := (X, X^{\mathbb{R}})$$

$\hookrightarrow$  all functions

Indiscrete qbs on  $X$



Validate gbs axioms for:  $W := ([0, \infty], \text{Meas}(\mathbb{R}, W))$

• Constants:

$E : \mathcal{B}_W, \alpha : W \vdash$

$$(\lambda r : \mathbb{R}. \alpha)^{-1}[E] = \begin{cases} \alpha \in E : & \mathbb{R} \\ \alpha \notin E : & \emptyset \end{cases} \in \mathcal{B}_{\mathbb{R}} \quad \checkmark$$

Validate qbs axioms for:  $\mathbb{W} := ([0, \infty], \text{Meas}(\mathbb{R}, \mathbb{W}))$

• Precomposition:

$\alpha: \text{Meas}(\mathbb{R}, \mathbb{W}), \varphi: \text{Meas}(\mathbb{R}, \mathbb{R}) \vdash$

$$\mathbb{R} \xrightarrow{\varphi} \mathbb{R} \xrightarrow{\alpha} \mathbb{W} \in \text{Meas}(\mathbb{R}, \mathbb{W})$$

$\downarrow$   
Meas is a cat.

Explicitly:

$$(\alpha \circ \varphi)^{-1}[E] \in \beta_{\mathbb{R}} \xleftarrow{\varphi^{-1}} \alpha^{-1}[E] \in \beta_{\mathbb{R}} \xleftarrow{\alpha^{-1}} E \in \beta_{\mathbb{W}} \quad \checkmark$$

Validate gbs axioms for:  $\mathcal{W} := ([0, \infty], \text{Meas}(\mathbb{R}, \mathcal{W}))$

•  $\mathbb{R}$  combination

$I$  ctbl,  $\alpha_i: \text{Meas}(\mathbb{R}, \mathcal{W})$ ,  $E_i: \mathcal{B}_{\mathbb{R}}$ ,  $\mathbb{R} = \bigcup_{i \in I} E_i$ ,  $F: \mathcal{B}_{\mathcal{W}}$

$$\left( \lambda r. \begin{cases} \vdots \\ r \in E_i : \alpha_i r \\ \vdots \end{cases} \right)^{-1} [F]$$

$\beta :=$

$$= \bigcup_{i \in I} \alpha_i^{-1} [F] \cap E_i \in \mathcal{B}_{\mathbb{R}}$$

In fact:

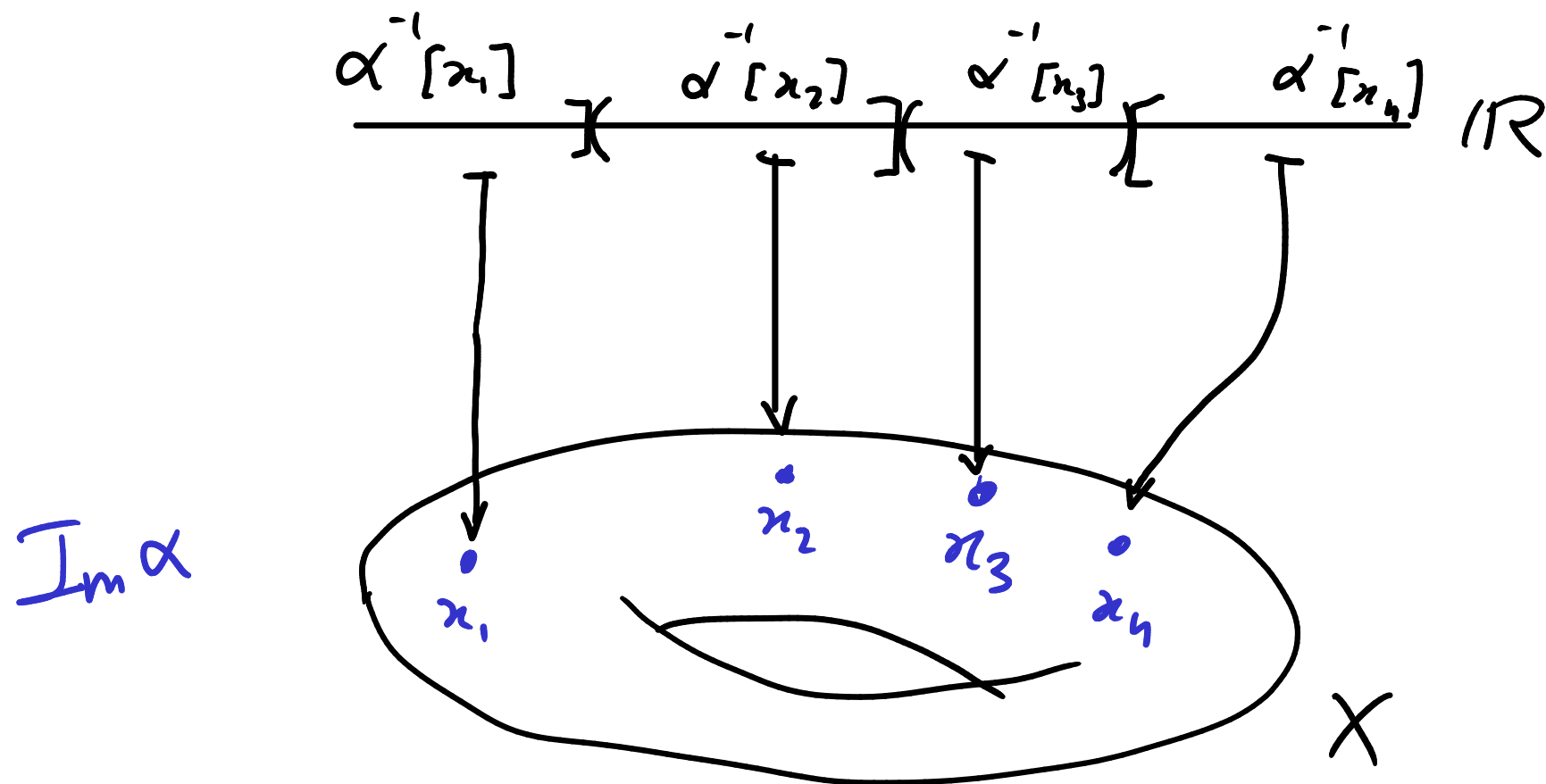
$$r \in \text{LHS} \Leftrightarrow \beta r \in F \Leftrightarrow \exists i \in I. r \in E_i \wedge \alpha_i r \in F \Leftrightarrow r \in \text{RHS}$$

✓

# $\sigma$ -simple function

$\alpha: \mathbb{R} \rightarrow X$  s.t.  $\text{Im } \alpha := \alpha[\mathbb{R}]$  is ctbl  $\wedge$

$\forall x \in \text{Im } \alpha. \alpha^{-1}[x] \in \mathcal{B}_{\mathbb{R}}$



Validate qbs axioms for:  $\Gamma^{\text{Obs}}$ ,  $X := (X, \sigma\text{-simple}(X))$

- Constants

$$\text{Im}(\lambda r. r) = \{r\} \text{ ctbl } \checkmark$$

NB:  $f$   $\sigma$ -simple:  
 $\text{Im } f$  ctbl  $\wedge$   
 $f^{-1}[x] \in \mathcal{B}_{\mathbb{R}}$

$$g: X \vdash (\lambda r. r)^{-1}[y] = \begin{cases} x=y: \mathbb{R} \\ x \neq y: \emptyset \end{cases} \in \mathcal{B}_{\mathbb{R}} \checkmark$$

Validate qbs axioms for:  $\Gamma X^{\text{obs}}$  :=  $(X, \sigma\text{-simple}(X))$

• Precomposition:

$\alpha: \sigma\text{-simple}(X), \varphi: \text{Meas}(\mathbb{R}, \mathbb{R}) \vdash$

$\text{Im}(\alpha \circ \varphi) \subseteq \text{Im} \alpha$  ctbl ✓

NB:  $f$   $\sigma$ -simple:  
 $\text{Im} f$  ctbl  $\wedge$   
 $f^{-1}[x] \in \mathcal{B}_{\mathbb{R}}$

$x: X \vdash$

$(\alpha \circ \varphi)^{-1}[x] = \varphi^{-1}[\alpha^{-1}(x)] \in \mathcal{B}_{\mathbb{R}}$  ✓

$\alpha^{-1}(x) \in \mathcal{B}_{\mathbb{R}}$

$\varphi: \mathbb{R} \rightarrow \mathbb{R}$  measurable

Validate qbs axioms for:  $\Gamma^{\text{obs}}$ ,  $X := (X, \sigma\text{-simple}(X))$

• recombination:

$$\alpha_i : (\sigma\text{-simple}(X))^I, E_i \in \mathcal{B}_{\mathbb{R}}, R = \bigoplus_{i \in I} E_i \vdash$$

NB:  $f$   $\sigma$ -simple:  
 $\text{Im } f$  ctbl  $\wedge$   
 $f^{-1}[x] \in \mathcal{B}_{\mathbb{R}}$

$$\text{Im} [E_i \cdot \alpha_i]_{i \in I} \subseteq \bigcup_{i \in I} \text{Im } \alpha_i \quad \text{ctbl} \quad \checkmark$$

$x : X \vdash$

$$[E_i \cdot \alpha_i]_{i \in I}^{-1}(x) = \bigcup_{i \in I} \alpha_i^{-1}[x] \cap E_i \in \mathcal{B}_{\mathbb{R}} \quad \checkmark$$

Prop:  $X: \text{Set}, A: \text{Obs} \vdash$

$$\bullet \forall f: X \rightarrow \perp A, \tilde{f}: \overset{\text{Obs}}{X} \rightarrow A$$

$$\bullet \forall f: \perp A \rightarrow X, \tilde{f}: A \rightarrow \overset{X}{\perp \text{Obs}}$$



Prop:  $X: \text{Set}, A: \text{Obs} \vdash$

$$\bullet \forall f: X \rightarrow \mathcal{L}A, \tilde{f}: \overset{\text{Obs}}{X} \rightarrow A$$

Pf:  $\alpha: \mathcal{R}_{\overset{\text{Obs}}{X}} \vdash \alpha \text{ } \sigma\text{-simple} \Rightarrow$

$$\alpha = [\alpha^{-1}[x].\lambda r. x]_{x \in \text{Im } \alpha} \Rightarrow$$

$$(f \circ \alpha) = [\alpha^{-1}[x].\lambda r. fx]_{x \in \text{Im } \alpha} \in \mathcal{R}_A \quad \checkmark$$

recombination

constat  $\in \mathcal{B}_A$  ctbl

Borel

Prop:  $X : \text{Set}, A : \text{Qbs} \vdash$

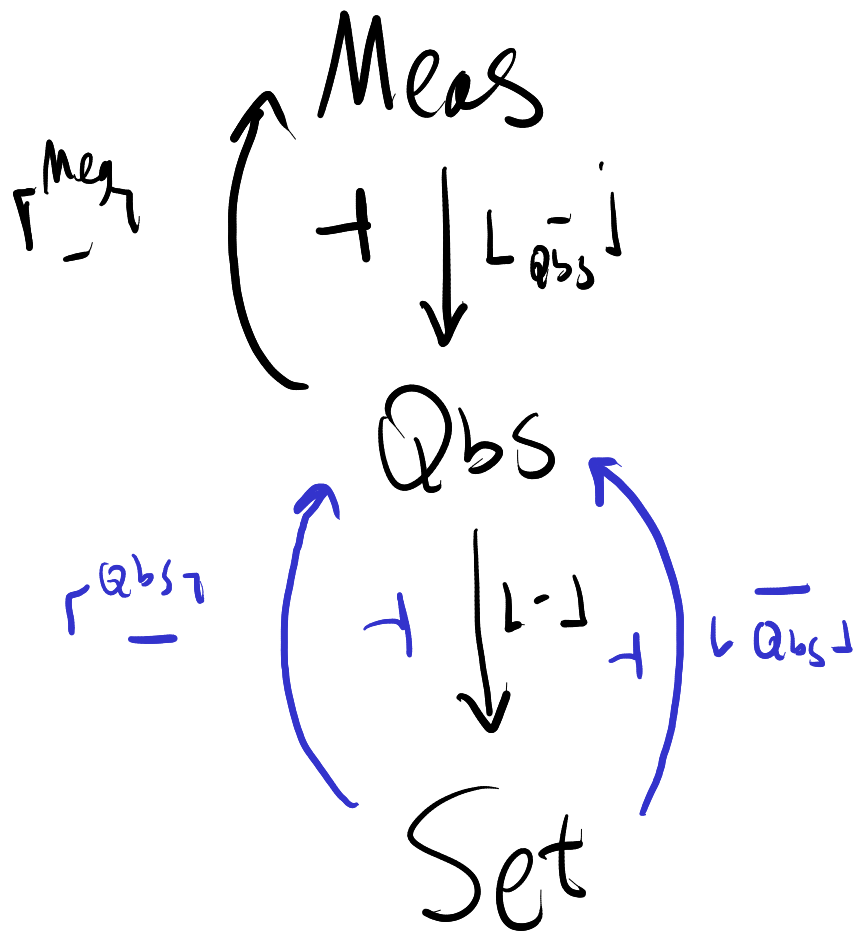
•  $\forall f : X \rightarrow \perp A, \quad f : \ulcorner X^{\text{Qbs}} \urcorner \rightarrow A$

•  $\forall f : \perp A, \quad \rightarrow X, \quad f : A \rightarrow \ulcorner X^{\text{Qbs}} \urcorner$

Prf:  $\alpha : R_A \vdash (f \circ \alpha : R \rightarrow X) \in R_{\ulcorner X^{\text{Qbs}} \urcorner}$  always. ✓



# Useful adjunctions:



$$\mathcal{L}_{\text{Obs}}^{\text{V}} := (\mathcal{L}_{\text{V}}, \text{Meas}(\mathbb{R}, V))$$

$$(V \in \text{Meas})$$

$$\mathcal{L}_{\text{X}}^{\text{meas}} := \left\{ A \subseteq \mathcal{L}_{\text{X}} \mid \forall \alpha \in \mathbb{R}_{\text{X}}. \alpha^{-1}[A] \in \mathcal{B}_{\mathbb{R}} \right\}$$

- limits (products, subspaces)  
and colimits (coproducts, quotients)  
as in Set

- Slogan: every measurable space is carried by a qbs

# Example

Product  $(X \times Y, \pi_1, \pi_2)$ :

-  $L_{X \times Y} = L_{X_1 \times Y_1}$  *necessarity!*

correlated  
random  
elements

-  $R_{X \times Y} = \{ \lambda r_0(\alpha r, \beta r) \mid \alpha \in R_X, \beta \in R_Y \}$

rest of structure as in Set.

# Function Spaces

Straight forward!

$$- \mathcal{Y}^X := \text{Obs}(X, \mathcal{Y})$$

$$- \mathcal{R}_{\mathcal{Y}^X} := \text{uncurry}[\text{Obs}(\mathbb{R} \times X, \mathcal{Y})]$$

$$= \left\{ \alpha: \mathbb{R} \rightarrow \mathcal{Y}^X \mid \lambda(r, x). \alpha r x: \mathbb{R} \times X \rightarrow \mathcal{Y} \right\}$$

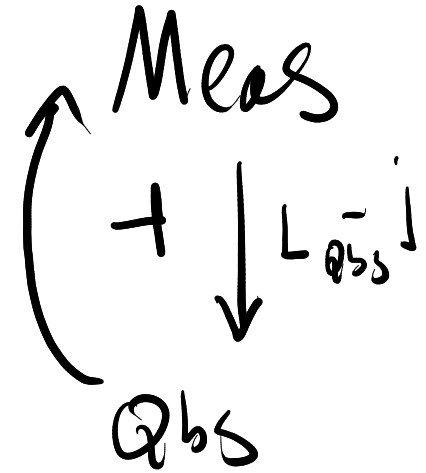
$$- \text{eval}: \mathcal{Y}^X \times X \rightarrow \mathcal{Y}$$
$$\text{eval}(f, x) := fx$$

# Meas vs Obs

By generalities:

$\sigma$ -algebra on  $\text{Meas}(\mathbb{R}, \mathbb{R})$

$\Gamma_{\text{meas}}$



$\Gamma_{\text{meas}} \mathbb{R}$

$\mathbb{R} \times \mathbb{R}$

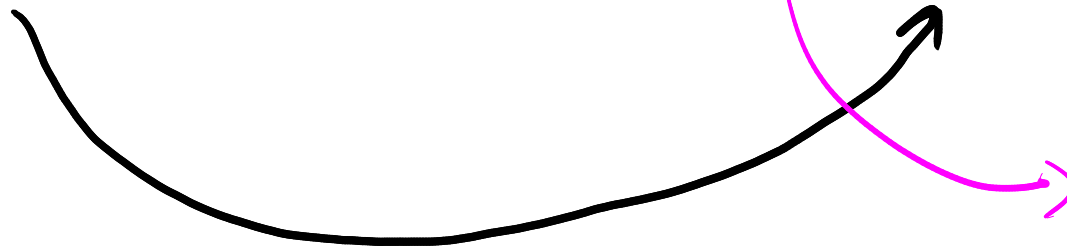
$\longrightarrow$

$\Gamma_{\text{meas}} \mathbb{R}$

$\mathbb{R} \times \mathbb{R}$

~~$\longrightarrow$~~

$\Gamma \mathbb{R} = \mathbb{R}$



$\Gamma_{\text{eval}}$

$\left( \Gamma_{\text{eval}} \mathbb{R}^{\mathbb{R}} \times \mathbb{R} \neq \Gamma \mathbb{R}^{\mathbb{R}} \times \Gamma \mathbb{R} \right)$

No factorisation by Aumann's Theorem.

# Simple Type Structure

"Simple" because:

- Simply-typed  $\lambda$ -calculus
- types are simple:  $A, B : \text{Type} \vdash B^A : \text{Type}$ 
  - no polymorphism
  - no term dependency
- Contexts for terms:  $\Gamma \vdash t : A$   
are simple:  $\Gamma = x_1 : A_1, \dots, x_n : A_n$   
i.e.  $\text{List}(\text{Type})$

# Simple Type Structure

"Simple" because:

- interpretation is simple:

$$\llbracket x_1:A_1, \dots, x_n:A_n \rrbracket := \prod_{i=1}^n A_i$$

$$\llbracket \Gamma \vdash t:A \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow A$$

in Qbs



# Simple Type Structure

Curry-Howard-Lambek

$$\frac{\Gamma \vdash t : A \quad \Gamma \vdash s : B}{\Gamma \vdash \langle t, s \rangle : A \times B} \rightsquigarrow$$

$$[\Gamma] \xrightarrow{\lambda r. \langle tr, sr \rangle} A \times B$$

is measurable

$$\frac{\Gamma \vdash t : A \times B \quad \Gamma, x:A, y:B \vdash s : C}{\Gamma \vdash \text{let } (x, y) = t \text{ in } s : C} \rightsquigarrow$$

$$\Gamma \vdash \text{let } (x, y) = t \text{ in } s : C \rightsquigarrow$$

$$\lambda r. \text{let } (a, b) = tr \text{ in } s \vdash [x \mapsto a, y \mapsto b]$$

$$[\Gamma] \xrightarrow{\quad} C$$

is measurable. etc.

measurability  
by  
type!

# Random element Space

$$\mathcal{R}_X := X^{\mathbb{R}} \quad \text{since} \quad \llbracket X^{\mathbb{R}} \rrbracket = \mathcal{R}_X \quad \text{as sets.}$$

Why?

$$(1) \quad \alpha \in \llbracket X \rrbracket^{\mathbb{R}} \Rightarrow \alpha: \mathbb{R} \rightarrow X \text{ in Obs.}$$

$$\text{id}_{\mathbb{R}}: \mathbb{R} \rightarrow \mathbb{R} \text{ measurable} \Rightarrow \text{id} \in \mathcal{R}_{\mathbb{R}}$$

$$\Rightarrow \alpha = \alpha \circ \text{id} \in \mathcal{R}_X$$

$$(2) \quad \alpha \in \mathcal{R}_X \Rightarrow \forall \psi \in \mathcal{R}_{\mathbb{R}} = \text{Meas}(\mathbb{R}, \mathbb{R}). \quad \alpha \circ \psi \in \mathcal{R}_X \Rightarrow \alpha: \mathbb{R} \rightarrow X \Rightarrow \alpha \in \llbracket X \rrbracket^{\mathbb{R}}$$

Pre composition  
↙

# Subspaces

For  $X \in \text{Obs}$ ,  $A \subseteq X$  Set:

$$R_A := \{ \alpha: \mathbb{R} \rightarrow A \mid \alpha \in R_X \}$$

Then  $A = (A, R_A)$  is the *subspace* qbs

We write  $A \hookrightarrow X$

# Borel subspaces ensemble

The  $\sigma$ -algebra  $B_X := \left\{ A \subseteq X \mid \forall \alpha \in R_X. \alpha^{-1}[A] \in B_{\mathbb{R}} \right\}$

internalises as  $B_X = \mathcal{Z}^X$ , the qbs of

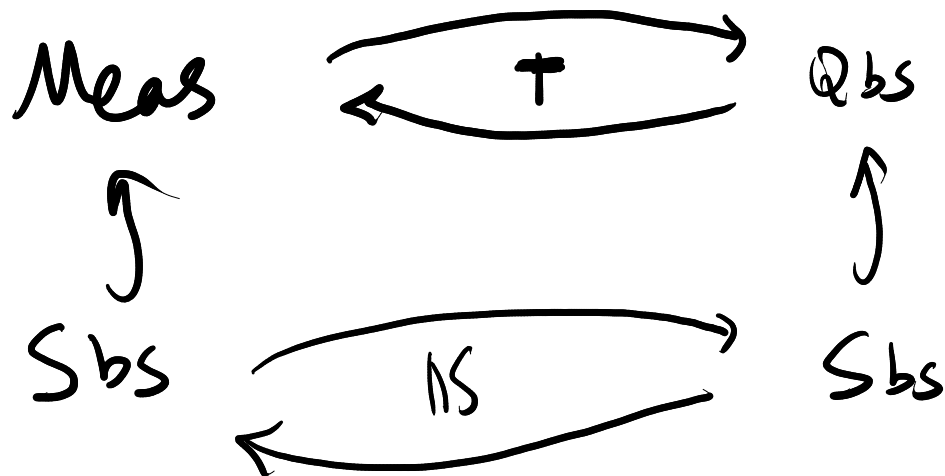
Borel subsets.

$\left( \begin{array}{l} B \\ \downarrow \\ \mathcal{L}(B_{\mathbb{R}}) \end{array} \right)$  are the Borel-on-Borel sets from  
descriptive set theory.  
cf. [Sabou et al. '21]

# Standard Borel Spaces

Def: A qbs  $S$  is **standard Borel** when

$S \cong A$  for some  $A \in \mathcal{B}_{\mathbb{R}}$



**Slogan:** Qbs conservative extension of Sbs

Example  $C_0 := \{f: \mathbb{R} \rightarrow \mathbb{R} \mid f \text{ continuous}\} \hookrightarrow \mathbb{R}^{\mathbb{R}}$

$C_0$  is sbs. (Well-known!)

Proof:

$C'_0 \in B_{\mathbb{R}^{\mathbb{Q}}}$  <sup>sbs!</sup>

$C'_0 := \left\{ g \in \mathbb{R}^{\mathbb{Q}} \mid \begin{array}{l} \forall a, b \in \mathbb{Q}, \varepsilon \in \mathbb{Q}^+ \\ \exists \delta \in \mathbb{Q}^+ \forall p, q \in \mathbb{Q}^+ \cap [a, b] \\ |p - q| < \delta \Rightarrow |g(p) - g(q)| < \varepsilon \end{array} \right\}$

on closed intervals  
(= compact intervals)  
Continuity  $\iff$  uniform continuity

Borel measurable } by type checks

then  $C_0 \cong C'_0 \in B_{\mathbb{R}^{\mathbb{Q}}}$ :

$C_0 \rightarrow C'_0$

$C'_0 \rightarrow C_0$

$\varphi \mapsto \varphi|_{\mathbb{Q}}$

$\psi \mapsto \lambda r. \lim_{n \rightarrow \infty} g(\text{approx } r \text{ by } (\frac{i}{n})_{i \in \mathbb{N}})_n$

## Example (ctd)

$C_0$  is sbs, and  $\text{eval}: C_0 \times \mathbb{R} \rightarrow \mathbb{R}$

is measurable.

Avoids:

- constructing complete separable metrics
- proving that evaluation is measurable w.r.t. metric  $\sigma$ -algebra.

# Non-examples

~ [Sabok et al. '21]

$$- \{ A \in \mathcal{B}_{\mathbb{R}} \mid A \neq \emptyset \} \hookrightarrow \mathcal{B}_{\mathbb{R}}$$

$$- \{ (A_1, A_2) \in \mathcal{B}_{\mathbb{R}}^2 \mid A_1 \subseteq A_2 \} \hookrightarrow \mathcal{B}_{\mathbb{R}}^2$$

$$- \{ A \in \mathcal{B}_{\mathbb{R}} \mid A \text{ open} \} \hookrightarrow \mathcal{B}_{\mathbb{R}}$$



# Plan:

- 1) Type-driven probability: discrete case (Mon + Tue)
- 2) Borel sets & measurable spaces (Wed)
- 3) Quasi Borel spaces (Wed) Simple type structure (Thu) ✓
- 4) Dependent type structure & standard Borel spaces (Thu) ✓
- 5) Integration & random variables (Fri)

please ask questions!



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# Dependent Type Structure

Types can contain terms: a type referring to a term

$$X: \text{Type}, E: B_X \vdash \{x \in X \mid x \in E\}: \text{Type}$$

a type, just like  
STLC

a term!

# Dependent Type Structure

Types can contain terms:

$$X: \text{Type}, E: B_X \vdash \{x \in X \mid x \in E\}: \text{Type}$$

a type referring to a term

a type, just like  
STLC

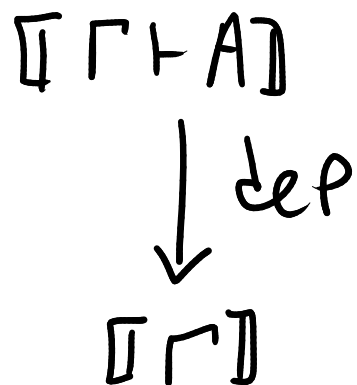
a term!

Content formation:

$$\frac{\Gamma \vdash A: \text{Type}}{\Gamma, x: A \vdash}$$

# Dependent Type Structure

types denote spaces-in-Content



Dependent types denote spaces-in-Content

$\Gamma \vdash$  ← Contract

$\llbracket \Gamma \vdash A \rrbracket$  ← Space in Content

$\Gamma \vdash A$   
 ↗  
 Type in Content

$\llbracket \Gamma \rrbracket$  ← Context Space

assigns environment

E.g.:

$A$

↓

↑

Simple types

$\llbracket E : B_A + \{x \in A \mid x \in E\} \rrbracket$

$\{ (E, a) \in B_A^{x_A} \mid a \in E \}$

↓  $\pi_1$   
 $B_A$

decoder

# Content extension

$$\frac{\Gamma \vdash A}{\Gamma, a:A \vdash}$$

$$\begin{array}{c} \llbracket \Gamma \vdash A \rrbracket \\ \text{dep} \downarrow \\ \llbracket \Gamma \rrbracket \quad \llbracket \Gamma, a:A \rrbracket := \llbracket \Gamma \vdash A \rrbracket \end{array}$$

# Substitution

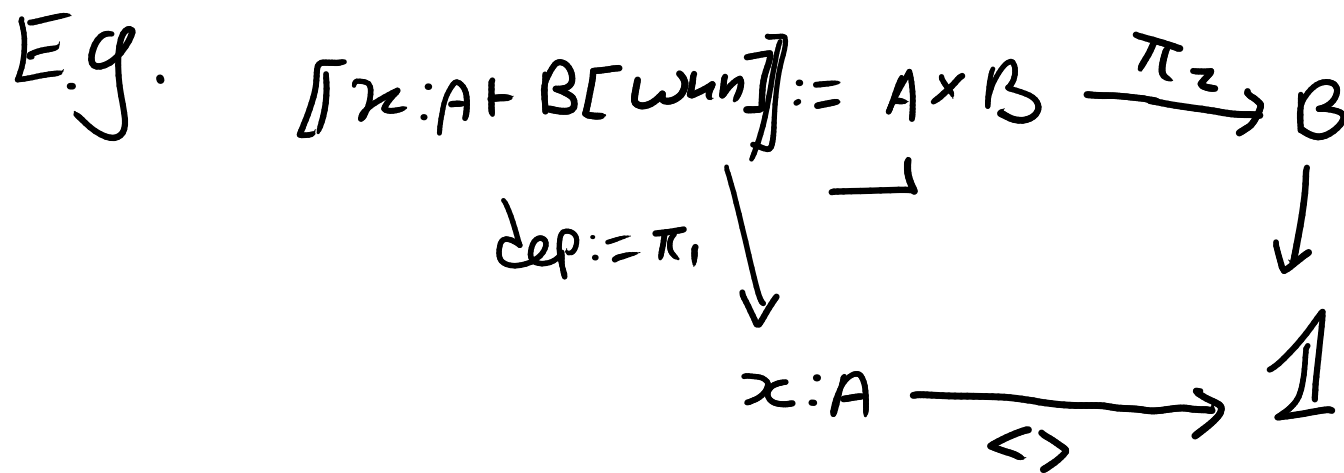
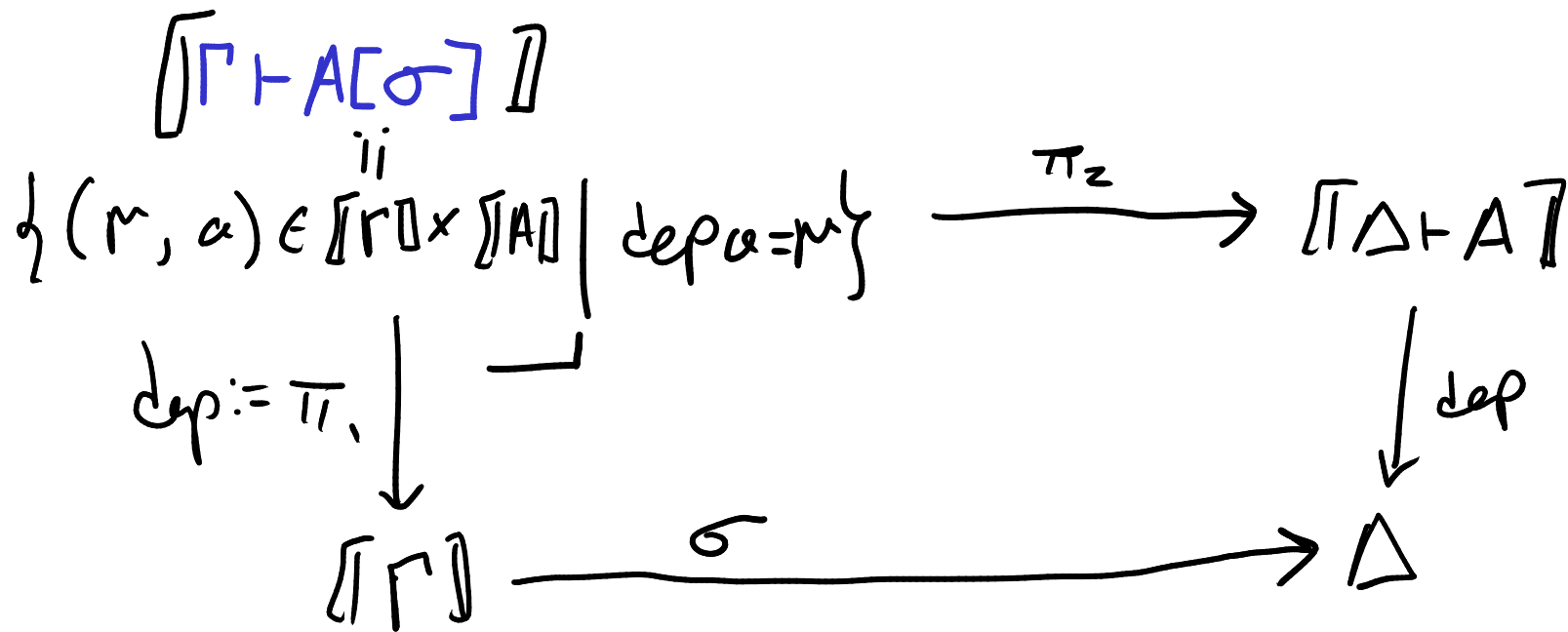
$$\frac{\Gamma \vdash \sigma:\Delta}{\text{E.g. weakening}}$$

$$\llbracket \sigma \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow \llbracket \Delta \rrbracket$$

$$\Gamma, a:A \vdash \text{Wkn} : \Gamma$$

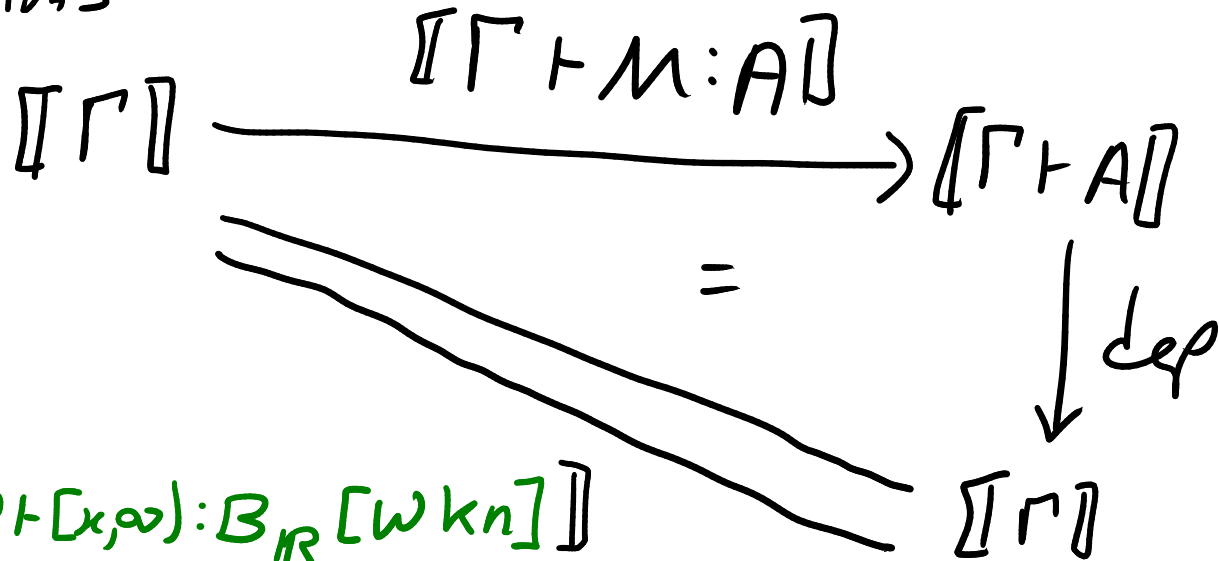
$$\llbracket \Gamma, a:A \rrbracket := \llbracket \Gamma \vdash A \rrbracket \xrightarrow[\text{dep}]{\text{Wkn}} \llbracket \Gamma \rrbracket$$

# Action of substitution on types

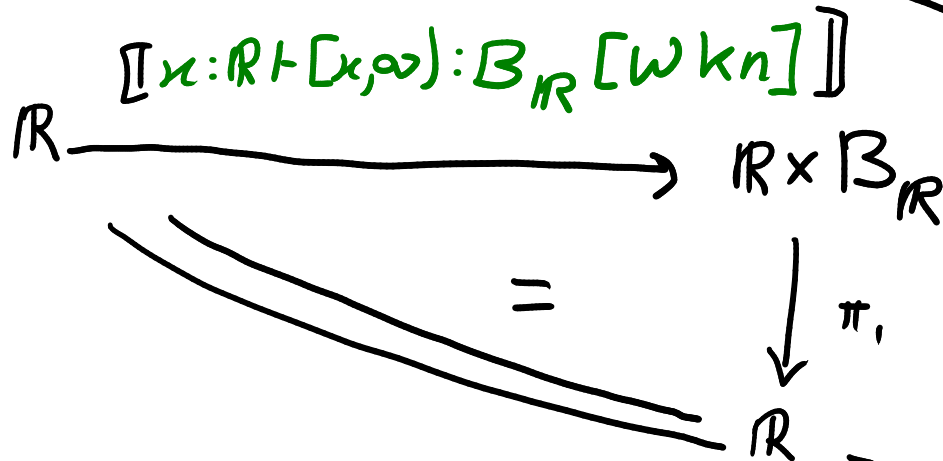


simple type

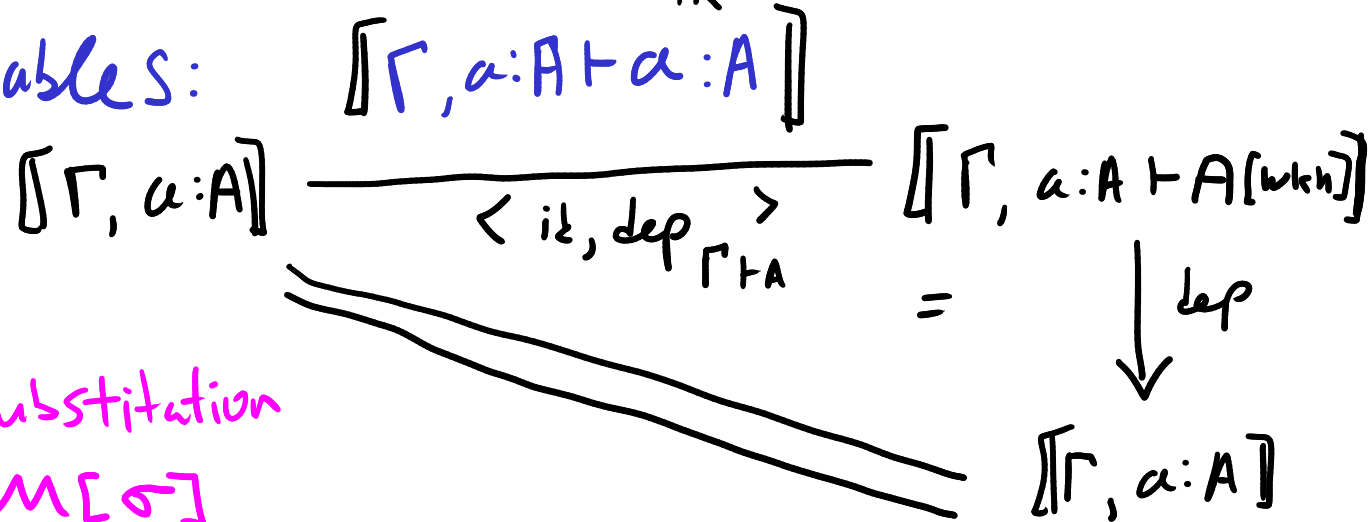
Terms : sections



e.g.



E.g. Variables:



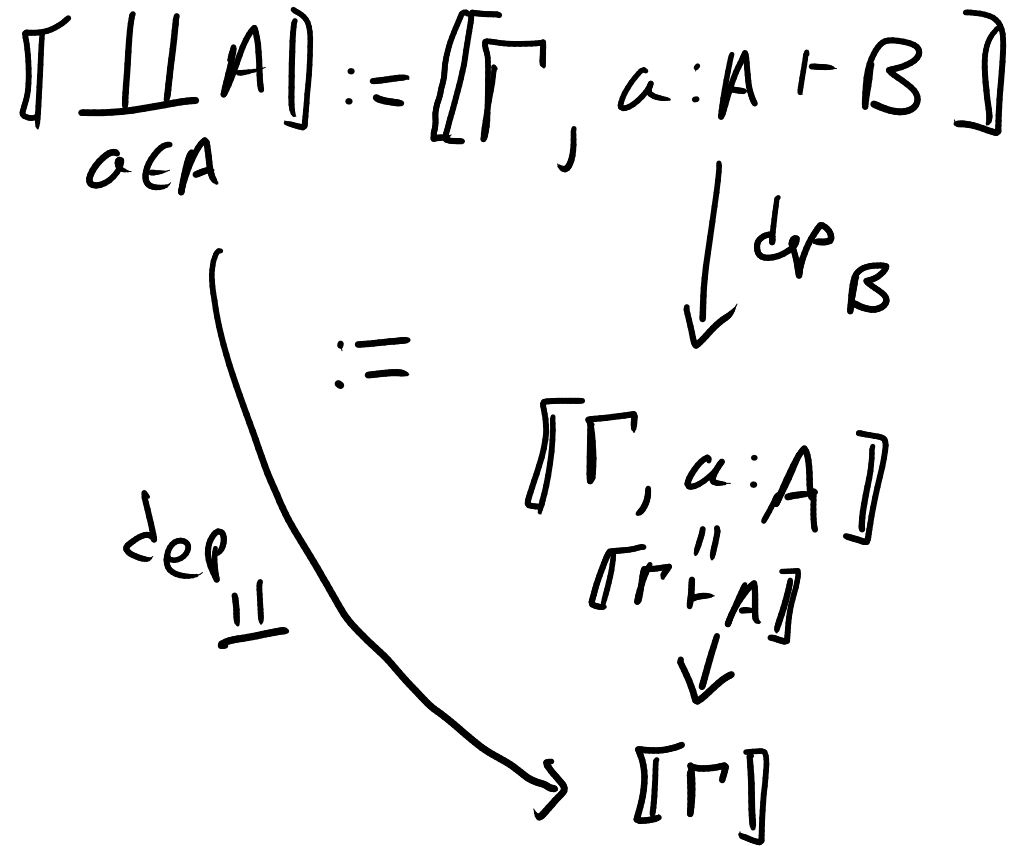
Exercise:

action of substitution  
 $M[\sigma]$



# Dependent Pairs

$$\frac{\Gamma, a:A \vdash B}{\Gamma \vdash \frac{\llcorner}{a \in A} B}$$



# Dependent products

$$\frac{\Gamma, a:A \vdash B}{\Gamma \vdash \prod_{a \in A} B}$$

$$\Gamma \vdash \prod_{a \in A} B$$

$$\llbracket \prod_{a \in A} B \rrbracket :=$$

$$\left\{ (M_0, f : \{ a \in \llbracket A \rrbracket \mid \text{dep } a = M_0 \} \rightarrow \llbracket \prod_{a:A} B \rrbracket) \mid \forall a \in \llbracket \prod_{a:A} B \rrbracket. \text{dep } a = M_0 \Rightarrow \text{dep } (f a) = a \right\}$$

aha:  $(a:A) \rightarrow B$

Exercise: find the random elements.

# Full model

type: Obs     $w := [0, \infty]$      $\mathcal{B}_X \cong \mathcal{B}^X$

$\mathcal{D}X := (\text{Fri})$

$\mathcal{P}_X := \left\{ \mu \in \mathcal{D}X \mid \mathcal{C}_\mu[X] = 1 \right\}$

$\mathcal{C}_\mu[E] := (\text{Fri})$      $\delta_x := (\text{Fri})$

$\oint \mu_k := (\text{Fri})$

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Please ask questions!



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# Foundations for type-driven probabilistic modelling

Ohad Kammar  
University of Edinburgh

Logic Summer School  
Australian National University  
4–16 December, 2023  
Canberra, ACT, Australia



THE UNIVERSITY of EDINBURGH  
**informatics**

**lfcs**

Laboratory for Foundations  
of Computer Science



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Alan Turing  
Institute

Facebook Research NCSC

## Partiality cf. [Vakar et al. '19]

A Borel embedding  $e: X \hookrightarrow Y$

- injective function  $e: \llbracket X \rrbracket \rightarrow \llbracket Y \rrbracket$

- its image is Borel:  $e[\llbracket X \rrbracket] \in \mathcal{B}_Y$

-  $e$  is Strong:  $\alpha \in R_X \iff e \circ \alpha \in R_Y$

## Examples

•  $\mathbb{1} \hookrightarrow \mathbb{2}$

•  $S$  is sbs  $\iff \exists S \hookrightarrow \mathbb{R}$

Def: A Partial map  $f: X \rightarrow Y$  is a morphism

$$f: X \rightarrow Y \perp \{\perp\}$$

Its domain of definition

$$f: (Y \perp \{\perp\})^X \quad \text{Dom } f := \{x \in X \mid f x \neq \perp\} : \text{Type}$$

Depend-type  
interpretation

$$\begin{array}{ccc}
 \llbracket \text{Dom } f \rrbracket & \xrightarrow{\quad} & \{y \in Y \mid y \in E\} \\
 \downarrow \text{dep} & \lrcorner & \downarrow \text{dep} \\
 \llbracket f : (Y \perp \{\perp\})^X \rrbracket & \xrightarrow{[E \mapsto \lambda x. f x \neq \perp]} & \llbracket E : B_Y \rrbracket
 \end{array}$$

## Plan:

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Please ask questions!

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# Full model

type: Obs     $w := [0, \infty]$      $\mathcal{B}_X \cong \mathcal{B}^X$

$\mathcal{D}X := (\text{Fri})$

$\mathcal{P}_X := \{ \mu \in \mathcal{D}X \mid C_{\mu}[X] = 1 \}$

$C_{\mu}[E] := (\text{Fri})$      $\delta_x := (\text{Fri})$

$\oint \mu_k := (\text{Fri})$

Def: A measure  $\mu$  over  $\mathbb{R}$  is a function

$$\mu: \mathcal{B}_{\mathbb{R}} \rightarrow \mathbb{W} := [0, \infty]$$

Satisfying the measure axioms:

$$E: \mathcal{B}^{\omega}$$

$$\mu \emptyset = 0, \quad \mu E = \mu(E \cap F) + \mu(E \cap F^c), \quad \mu\left(\bigcup_n E_n\right) = \sup_n \mu E_n$$

For measurable spaces, replace  $\mathbb{R}$  with  $V$

We write  $\mathcal{G}V$  for the set of measures on  $V$

For qbs  $X$ , take  $\mathcal{G}^{\text{meas}} X$

Thm (Lebesgue measure):

There is a unique measure  $\lambda \in \mathcal{LGR}$  s.t.:

$$\lambda(a, b) = b - a$$

Thm (Lebesgue measure):

There is a unique measure  $\lambda \in \mathcal{L}(\mathbb{R})$  s.t.:

$$\lambda(a, b) = b - a$$

Proof sketch (standard analysis textbook):

1) restrict attention to  $(0, 1]$  & extend via  $\sigma$ -additivity

2) Take  $\Sigma_0 \subseteq \mathcal{B}_{(0, 1]}$   $E \in \Sigma_0 \Leftrightarrow E = \bigcup_{i=1}^n (a_i, b_i]$  ←

3) Defining  $\lambda: \Sigma_0 \rightarrow \mathbb{W}$ ,  $\lambda \left( \bigcup_{i=1}^n (a_i, b_i] \right) := \sum_{i=1}^n (b_i - a_i)$  independent of

4)  $\lambda \emptyset = 0$ ,  $\lambda E = \lambda(E \cap F) + \lambda(E \cap F^c)$  (straightforward)

↳

5) **Technical gadget**:  $\forall (E_n \supseteq E_{n+1})$  in  $\Sigma_0$ ,

$$\inf \lambda E_n > 0 \Rightarrow \bigcap E_n \neq \emptyset.$$

6)  $\lambda$  is continuous on  $\Sigma_0$ : If  $(E_n \subseteq E_{n+1})_n$  in  $\Sigma_0$

$$\text{and } \bigcup_n E_n \in \Sigma_0 \text{ then } \lambda \bigcup E_n = \sup_n \lambda E_n$$

7) Noting that:  $\Sigma_0$  is a Boolean algebra

$$\& \sigma(\Sigma_0) = \mathcal{B}_{(0,1]}$$

we use Carathéodory's extension theorem:

$$\lambda \text{ extends uniquely to } \lambda: \mathcal{B}_{(0,1]} \rightarrow \mathbb{W}.$$

# The unrestricted Giry spaces

Equip  $\mathcal{G}V$  with two qbs structures:

$$X \quad \mathcal{R}_{\mathcal{G}V} := \left\{ \alpha: \mathbb{R} \rightarrow \mathcal{G}V \mid \forall A \in \mathcal{B}_V, \lambda r. \alpha(r, A): \mathbb{R} \rightarrow \mathcal{W} \right\}$$

$$\checkmark \quad \mathcal{G}V \longleftrightarrow \mathcal{W}^{\mathcal{B}_X}$$

$\hookrightarrow \alpha$  is a kernel.

- Fewer random elements  
 $\mathcal{R}_{\mathcal{G}V} \subseteq \mathcal{R}_{\mathcal{G}'V}$
- Lebesgue integral  
measurable in  
both arguments.

## Farewell Meas

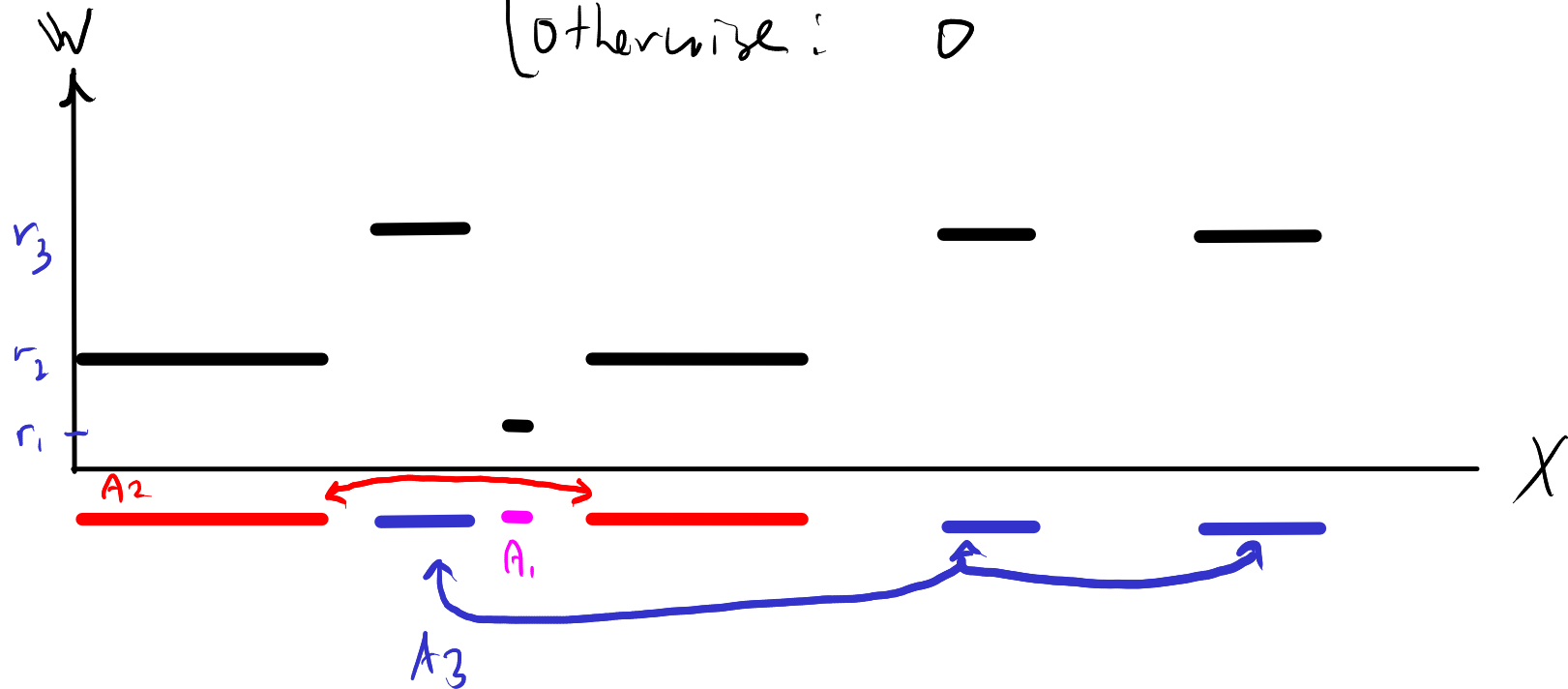
Now on:

1. All spaces are quasi-Borel (upcoming)
2. "measurable function" means qbs morphism!

Def: Simple function  $\varphi: X \rightarrow W$  when

$\exists n \in \mathbb{N}, \vec{A} \in \mathcal{B}_X^n, A_i \cap A_j = \emptyset, \vec{r} \in W$  s.t.  
 $(i \neq j)$

$$\varphi(x) = \begin{cases} r_i & x \in A_i \\ 0 & \text{otherwise} \end{cases}$$



Encode into a space:

$$\text{Simple Code} := \prod_{n \in \mathbb{N}} B_X^n \times W^n$$

$$\text{Simple} := \{ f \in W^X \mid f \text{ simple} \} \hookrightarrow W^X$$

and define an interpretation:

$$\llbracket - \rrbracket : \text{Simple Code} \longrightarrow \text{Simple}$$

$$\llbracket (n, \vec{A}, \vec{r}) \rrbracket := \sum_{i=1}^n r_i \cdot [- \in A_i]$$

↳ characteristic function  
for  $A_i$



Lemma:  $f: X \rightarrow W$  is measurable → remember!  
965  
morphisms!

iff  $f = \lim_{n \rightarrow \infty} f_n$  for some monotone sequence

$f_n \in \text{Simple}$ .

Moreover, we have measurable such choice:

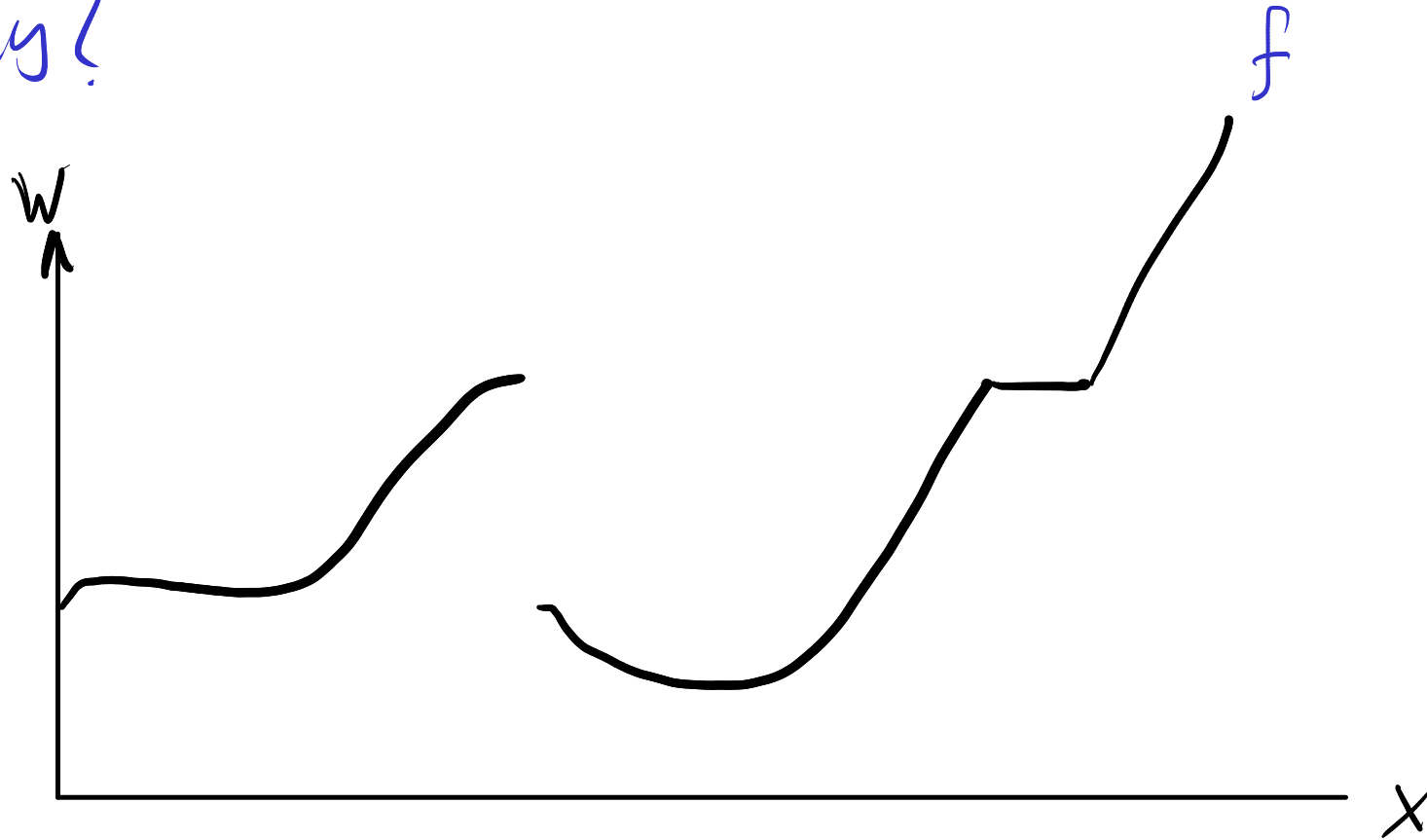
Simple Approx:

$\left\{ \vec{\Delta} \in \mathbb{R}^+ \mid \Delta_n \rightarrow 0 \right\} \times \left\{ \vec{a} \in W^{\mathbb{N}} \mid \begin{array}{l} \vec{a} \text{ monotone} \\ a_n \rightarrow \infty \end{array} \right\} \times W^X \rightarrow \text{Simple Code}$

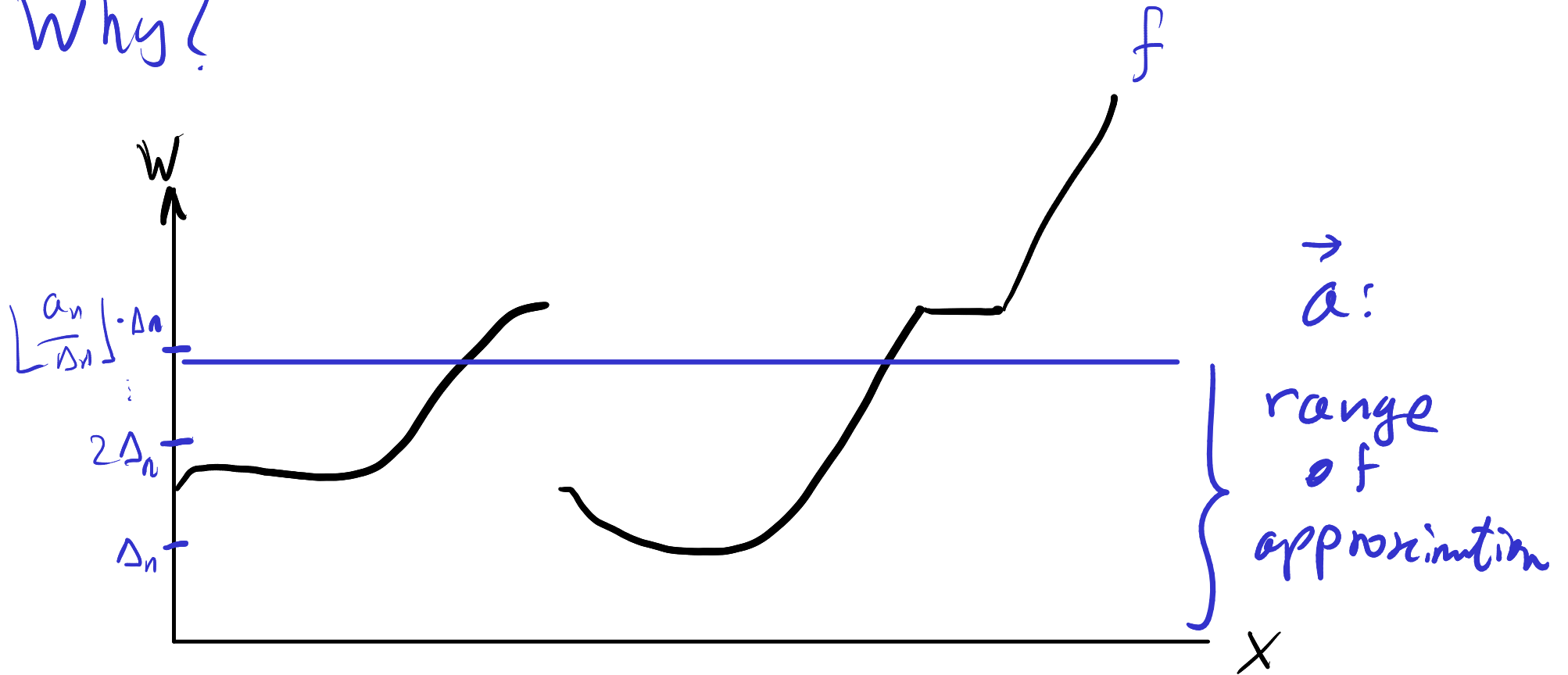
↑  
rate of  
convergence

↑  
range of  
approximation

Why?

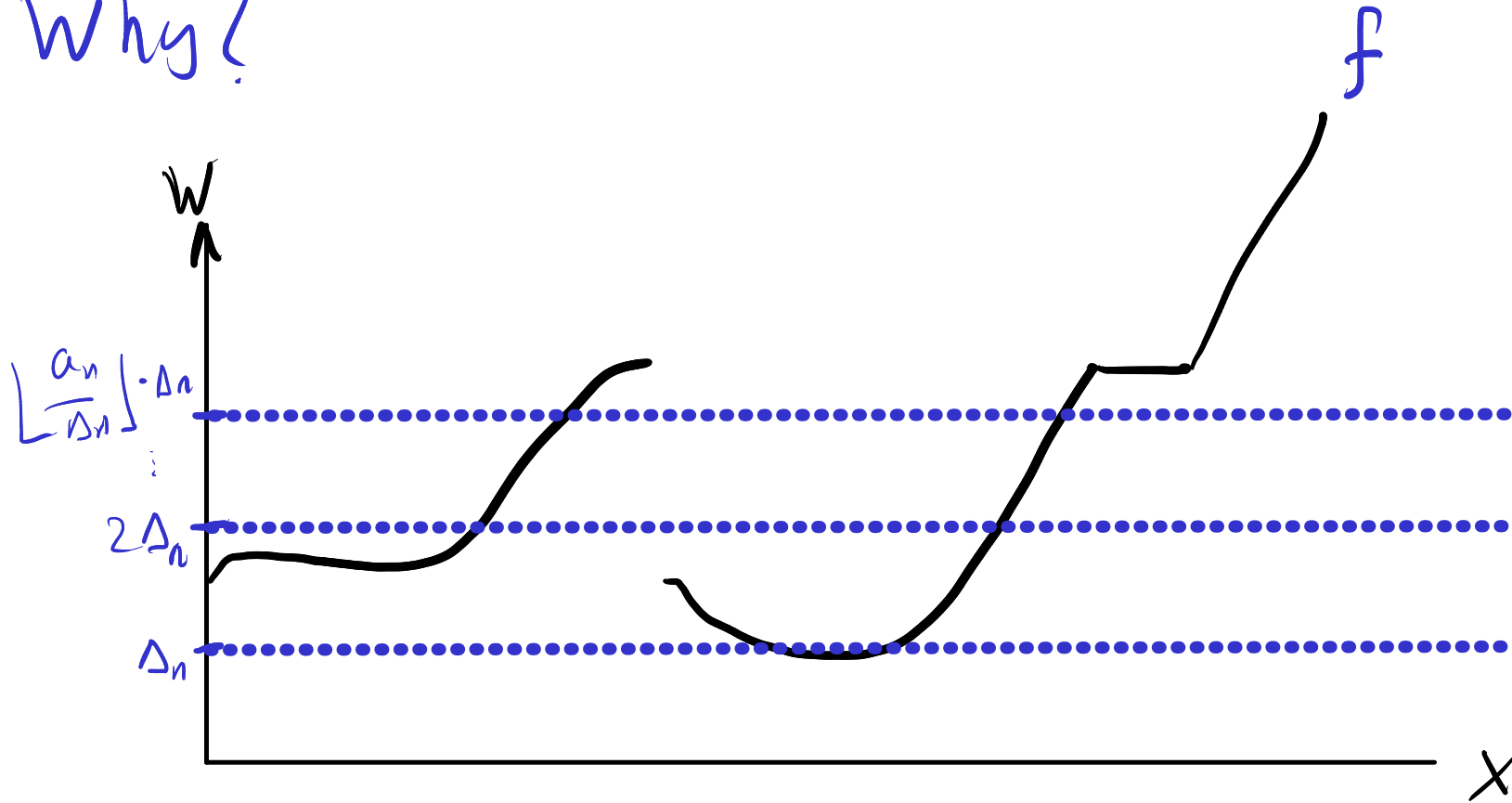


Why?

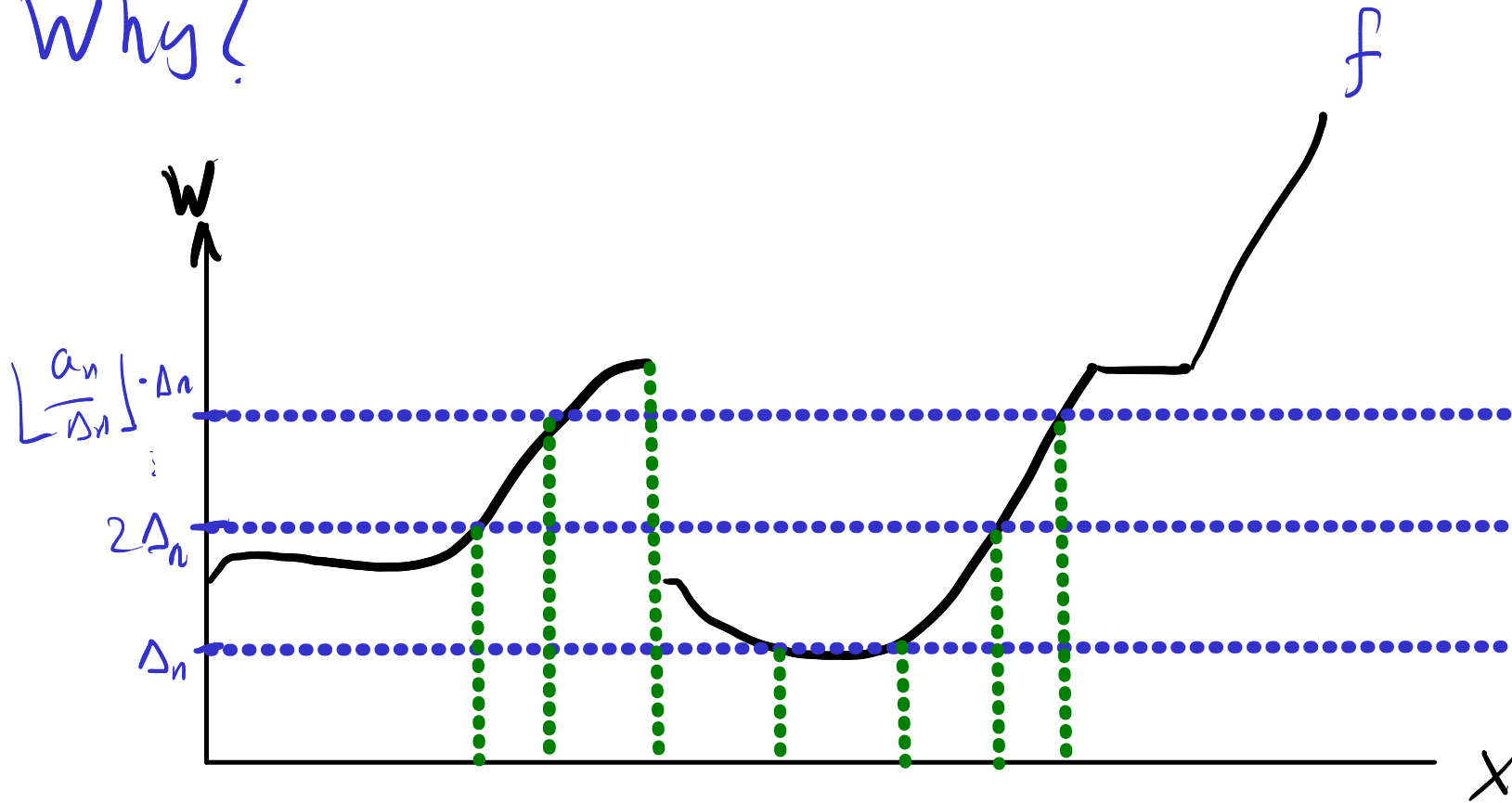


↑ resolution of approximation

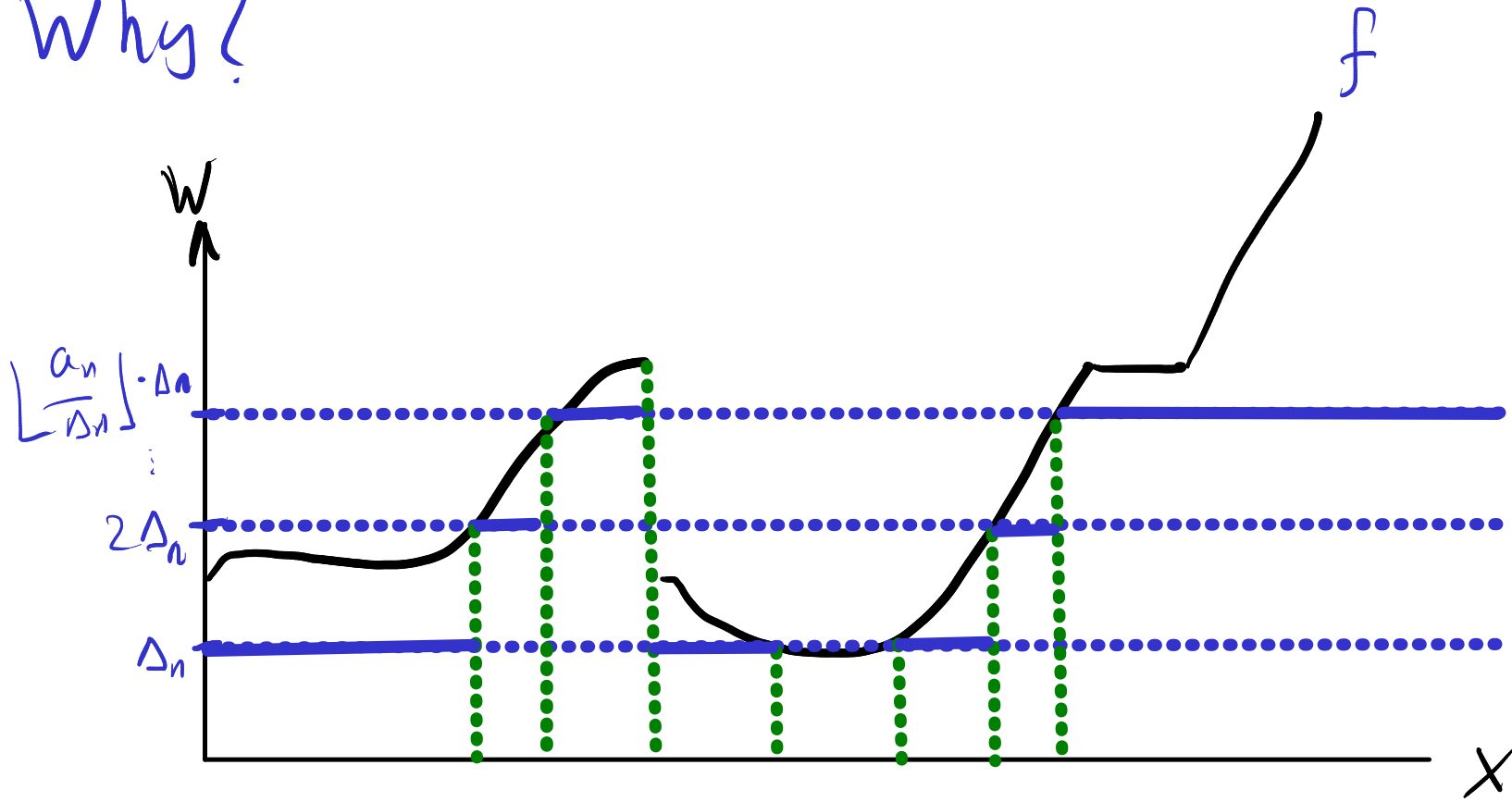
Why?



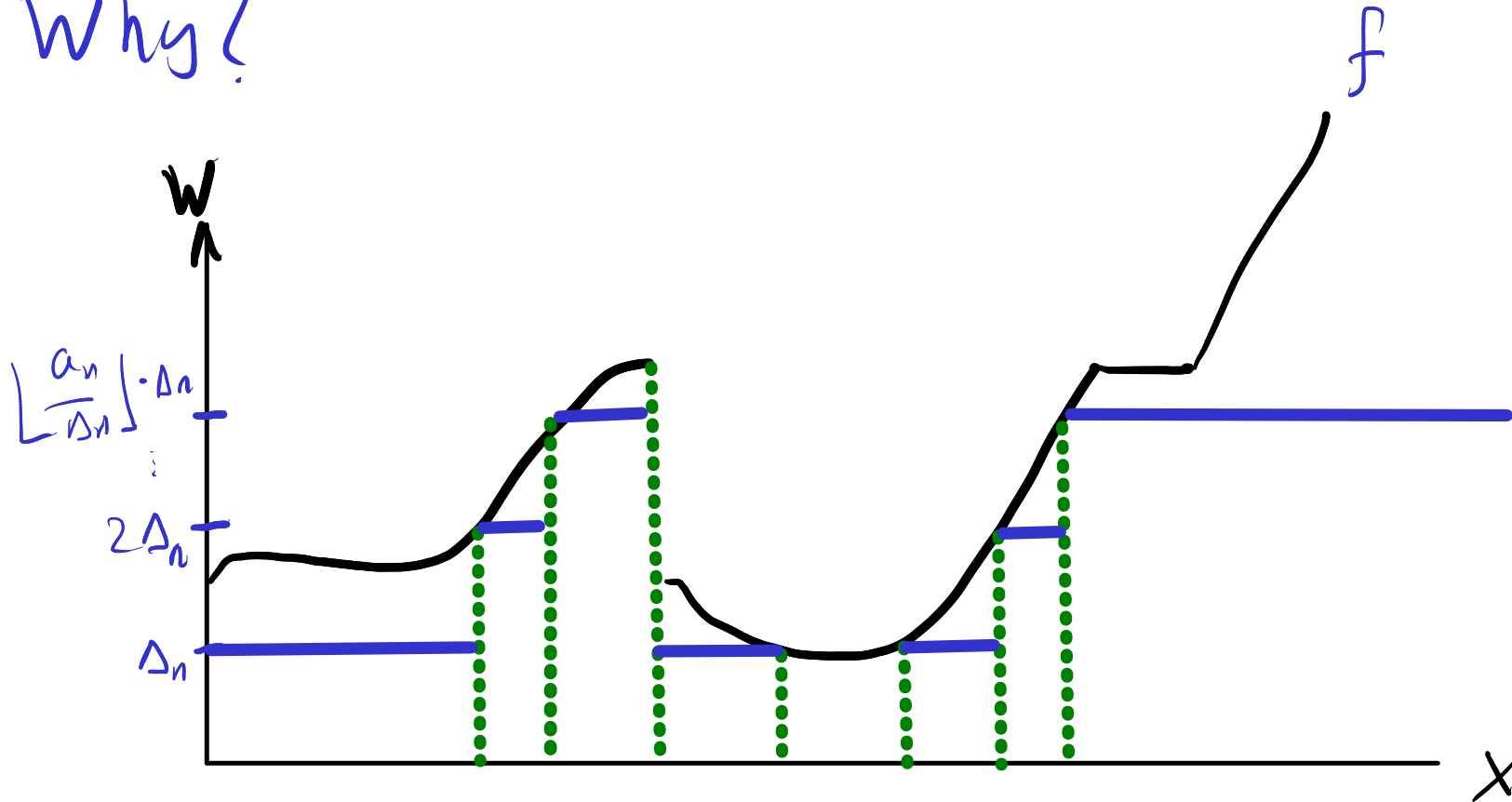
Why?



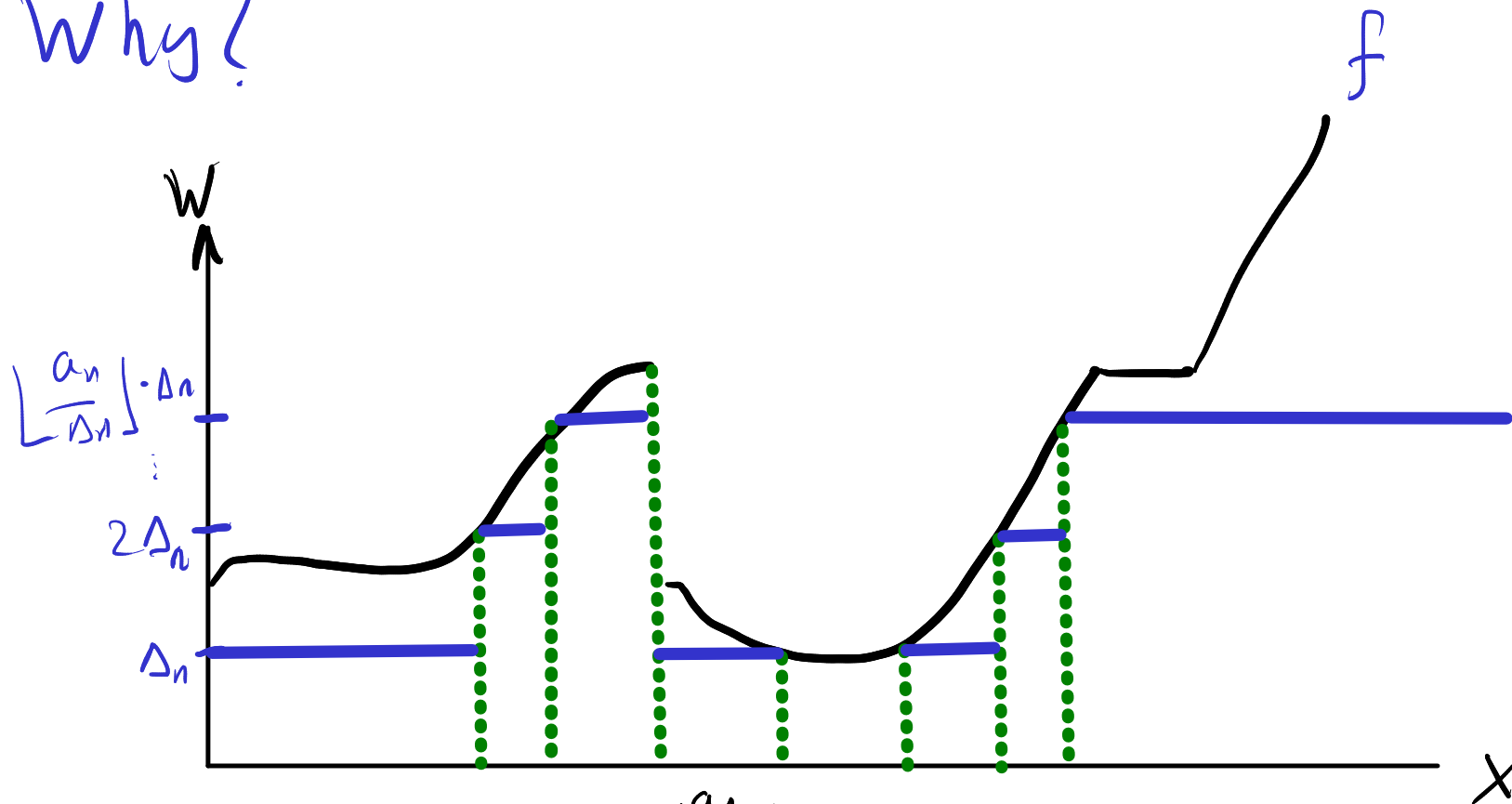
Why?



Why?



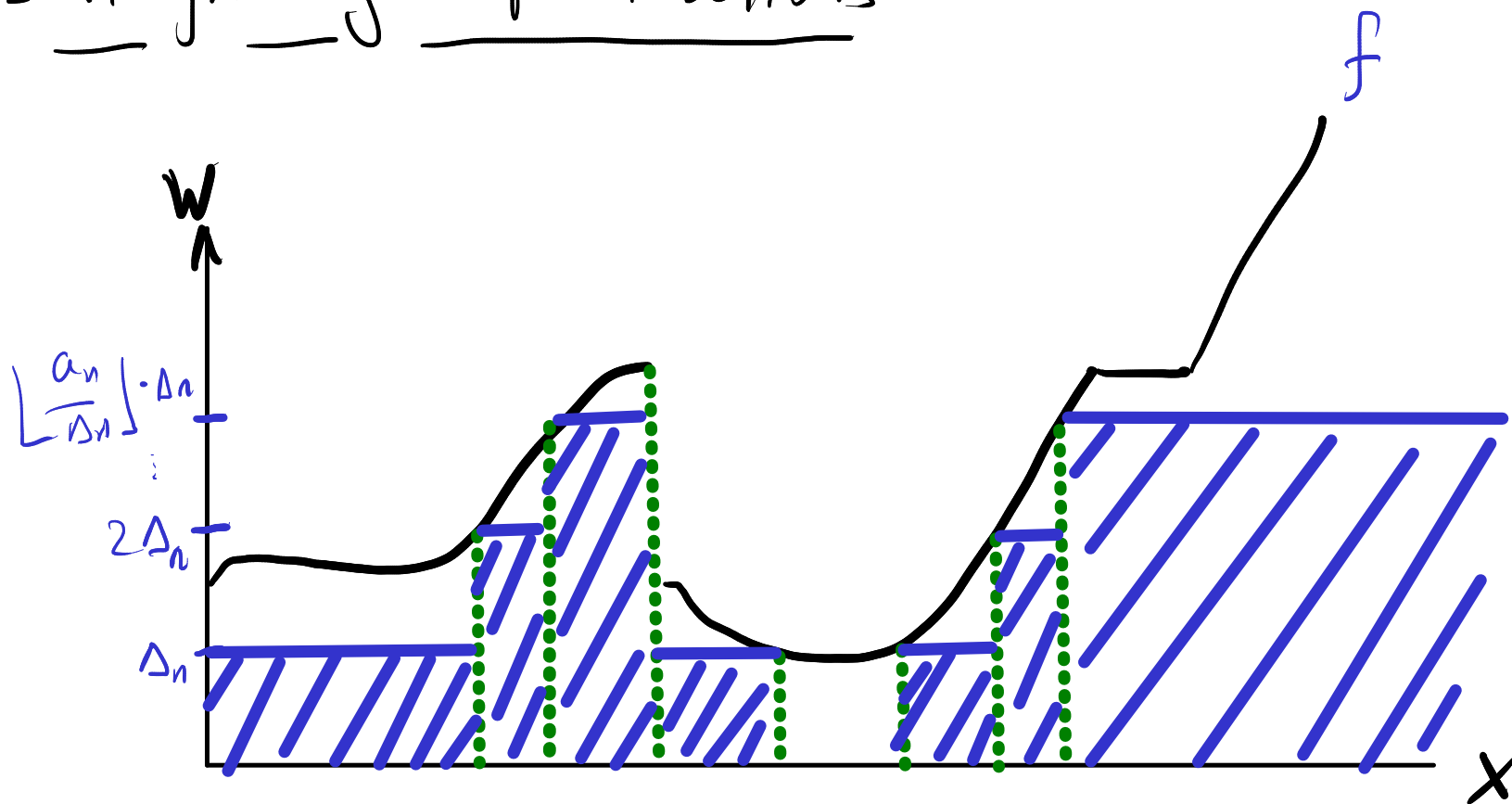
Why?



$$\left[ \text{Simple Approx}_{\Delta, a} f \right]_0 = \sum_{i=1}^{\lfloor \frac{a_n}{\Delta_n} \rfloor} i \cdot \Delta_n \left[ i \cdot \Delta_n \leq f < (i+1) \Delta_n \right] + \lfloor \frac{a_n}{\Delta_n} \rfloor \Delta_n \left[ f \geq \lfloor \frac{a_n}{\Delta_n} \rfloor \cdot \Delta_n \right] \in \text{Simple}$$



# Integrating Simple Functions



$$\int : \mathcal{G}X \times \text{Simple Code} \rightarrow \mathbb{W}$$

$$\int \mu(n, \vec{A}, \vec{r}) := \sum_{I \subseteq \{1, \dots, n\}} \left( \sum_{i \in I} r_i \right) \cdot \mu \left( \bigcap_{i \in I} A_i \setminus \bigcup_{i \notin I} A_i \right)$$

# Integration

Proper higher-order operation

$$\int : G \times W^X \longrightarrow W$$

$$\int \mu f := \sup \{ \int \mu \varphi \mid \varphi \in \text{Simple}, \varphi \leq f \}$$

we also write

$$\int \mu(dx) t$$

for  $\int \mu(x, t)$

$$= \lim_{n \rightarrow \infty} \int \mu(\text{Simple Approx}_{\Delta_n, \vec{a}_n} f)_n$$

measurable by type

$$\text{for } \frac{a_n}{\Delta_n} \rightarrow 0, \text{ e.g. } \Delta_n = \frac{1}{2^n}, a_n = n.$$

resolution

# The unrestricted Giry Strong Monad

Dirac:

$$\delta: X \rightarrow GX$$

$$x \mapsto \lambda A. \begin{cases} x \in A: 1 \\ x \notin A: 0 \end{cases}$$

Unlike the unrestricted Giry on Meas.

but: non-commutative

Kleisli extension / Kock integral:

$$\oint: GX \times GX^X \rightarrow GP$$

$$\oint \mu f := \lambda A. \int \mu(d\alpha) f(\alpha; A)$$

(Fubini fails,  
just line in  
Meas)

# Fubini-Tonelli fails

$$\# \in \mathbb{GIR} \quad \# E := \begin{cases} E \text{ finite:} & |E| \\ \text{o.w.} & ; \quad \infty \end{cases}$$

$$\lambda \in \mathbb{GIR}$$

Lebesgue

$$k: \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{W} \cong \mathbb{G}\mathbb{1}$$

$$\int \#(dr) \int \lambda(dx) k(x,y) = \int \# \underline{0} = \underline{0} \approx 0$$

$$y: \mathbb{R} + \{\infty\} \mapsto \lambda\{y\} \cdot 1 + \lambda\{y\}^c \cdot 0 = 0 \quad \#$$

$$k(x,y) := [x=y]$$

$$\int \lambda(dx) \int \#(dr) k(x,y) = \int \lambda(dx) \delta_{\infty} \approx \infty$$

$$x: \mathbb{R} + \{\infty\} \mapsto \# \{x\} + 0 = 1$$

# Randomisable measures monad

$$D \rightsquigarrow G$$

$$L_{DX} := \left\{ \lambda_\alpha \mid \alpha: \mathbb{R} \rightarrow X \right\}$$

$\lambda_A \int_{\text{Dom } \alpha} \lambda(\text{Dom } \alpha)$   
 Lebesgue measure

$$R_{DX} := \left\{ \lambda_\pi \cdot \lambda_{d\pi} \mid \alpha: \mathbb{R} \times \mathbb{R} \rightarrow X \right\}$$

$$\delta: X \rightarrow DX \quad \oint: D\Gamma^*(DX)^P \rightarrow DX \quad \text{lift along } D \rightsquigarrow G.$$

$D$  validates our measure axioms including Fubini-Tonelli  
 $\mu \in DX, \nu \in DY \vdash$

$$\oint \mu(dx) \oint \nu(dy) \delta_{(x,y)} = \oint \nu(dy) \oint \mu(dx) \delta_{(x,y)} =: \mu \otimes \nu$$

Thm: For sbs  $S$ ,  $PS$ ,  $D_{\leq 1}S$ ,  $D_{< \infty}S \in Sbs$   
and agree with their counterparts on Meas.

$$\mathcal{D}S_S = \{ \mu \mid \mu \text{ s-finite} \}$$

See [Staton'16]

$$\mathcal{R}_{DS} = \{ k: \mathbb{R} \rightarrow GD \mid k \text{ s-finite kernel} \}$$

Open: Is there a counterpart to  $D$  in Meas?

More modestly, is  $DS \in Sbs$ ?

(Hypothesis: **no**)

# Distribution Submonoids

A measure space

$$\Omega = (\Omega, \mu)$$

is a qbs  $\Omega$  with  
 $\mu \in \mathcal{D}_X$ .

Similarly: finite measure space  
- (sub) probability space.

$$\mathcal{P}_X := \{ \mu \in \mathcal{D}_X \mid \mu_X = 1 \}$$



$$\mathcal{P}_{\leq 1} X := \{ \mu \in \mathcal{D}_X \mid \mu_X \leq 1 \}$$



$$\mathcal{P}_{< \infty} X := \{ \mu \in \mathcal{D}_X \mid \mu_X < \infty \}$$



$\mathcal{D}_X$

# Full model

type: Obs     $w := [0, \infty]$      $\mathcal{B}_X \cong \mathcal{B}^X$

$\mathcal{D}X := ( \{ \lambda_\alpha \mid \alpha: \mathbb{R} \rightarrow X \}, \{ \lambda_{r, \lambda} \mid \alpha: \mathbb{R} \times \mathbb{R} \rightarrow X \} )$

$\mathcal{P}X := \{ \mu \in \mathcal{D}X \mid C_\mu[X] = 1 \}$

$C_\mu[E] := \mu E$      $\delta_x := E \mapsto \begin{cases} x \in E: 1 \\ x \notin E: 0 \end{cases}$

$\oint \mu k := \lambda E. \int \mu(\lambda) k(\lambda; E)$



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please ask questions!

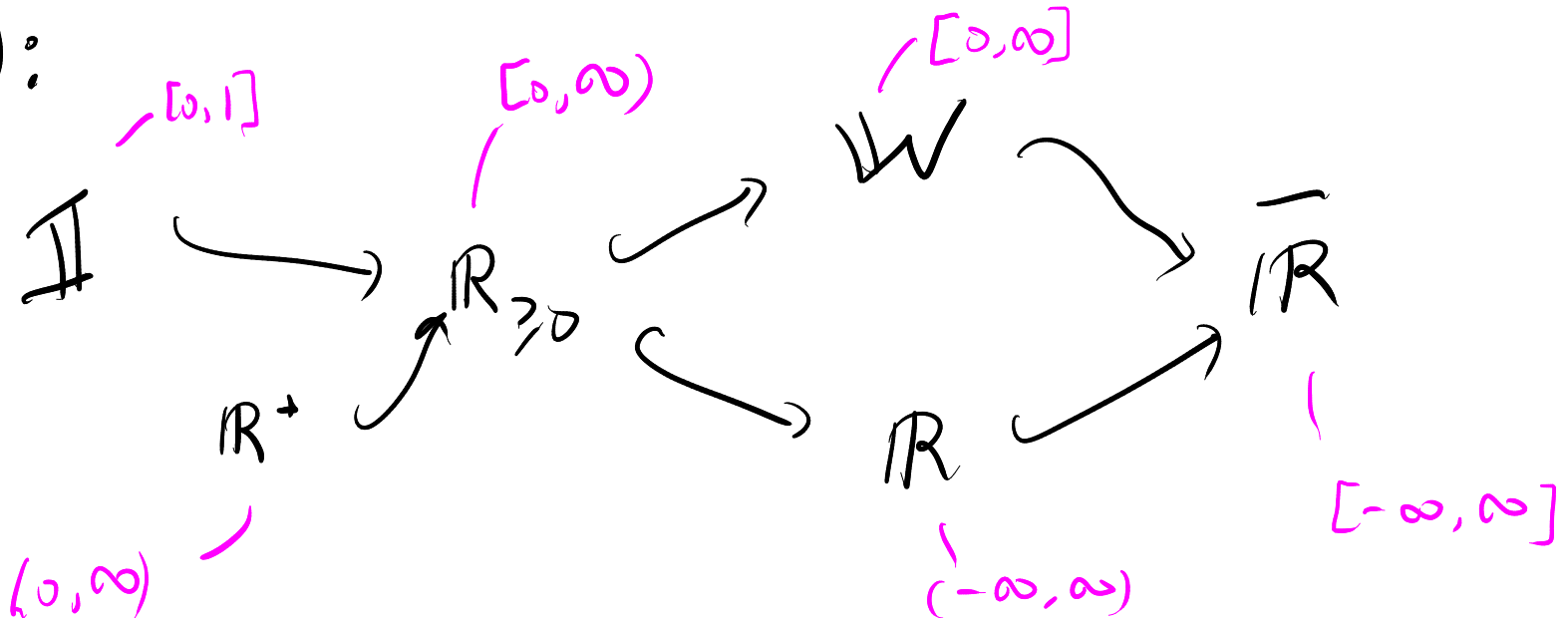


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Random variable  $\xi: \Omega \rightarrow \mathbb{H} \hookrightarrow \overline{\mathbb{R}}$

$\mathbb{H}$ :



- $\mathbb{H}^\Omega$  is a space
- $W^\Omega$  measurable  $\sigma$ -Semi-module for  $W$ :  $\sum_{n=0}^{\infty} \alpha_n \xi_n := \lambda W, \sum_{n=0}^{\infty} \alpha_n \cdot \xi_n$
- $\mathbb{R}^\Omega$  measurable vector space:  $\alpha \xi + \zeta := \lambda W, \alpha \cdot \xi W + \zeta W$

$$Pr: \mathcal{P}\Omega \times \mathcal{B}_\Omega \rightarrow \mathbb{W}$$

$$Pr_\lambda A := \text{eval}(\lambda, A) = \lambda A$$

Probability Space  $\Omega = (\Omega, \lambda_\Omega)$

$P: \mathcal{P}\Omega \vdash$  "  $Px$  holds  $\lambda(x)$ -almost surely "

for some  $Q \hookrightarrow \Omega$ ,  $P \supseteq Q$ ,  $[-\in Q] \cdot \lambda = \lambda$

Example  $(\xi, \zeta \in \mathbb{H}^\Omega)$

$\xi = \zeta$  a.s., when  $Pr_{\omega \sim \lambda} [\xi \omega \neq \zeta \omega] = 0$

# Integrating Random Variables (as discretely)

$$(-)_+, (-)_- : \mathbb{R}^{\Omega} \longrightarrow \mathbb{W}^{\Omega}$$

in Obs!

$$\xi_+ := \max(\xi, 0) \quad \xi_- := \max(-\xi, 0)$$

So:  $\xi = \xi_+ - \xi_-$

$$\int : \mathcal{P}\Omega \times \mathbb{W}^{\Omega} \longrightarrow \mathbb{W}$$

$\int$  respects  
a.s. equality:

$$\int \lambda \xi := \int \lambda \xi_+ - \int \lambda \xi_-$$

$$\xi = \zeta \text{ (a.s.)}$$

$$\Rightarrow \int \lambda \xi = \int \lambda \zeta$$

# Example

$$\lambda: \mathcal{P}\Omega \vdash \text{AS Converge}(\overline{\mathbb{R}})^{\Omega} : \mathcal{B}(\overline{\mathbb{R}}^{\mathbb{N} \times \Omega})$$
$$:= \left\{ \vec{z} \in \overline{\mathbb{R}}^{\mathbb{N} \times \Omega} \mid \Pr_{\omega \sim \lambda} [\lim_{n \rightarrow \infty} z_n \omega \neq \perp] \right\}$$

So:

$$f_{\text{lim}}^{\text{as}}: \overline{\mathbb{R}}^{\mathbb{N} \times \Omega} \longrightarrow \overline{\mathbb{R}}^{\Omega} \quad \text{Dom } f_{\text{lim}}^{\text{as}} := \text{AS Converge}(\overline{\mathbb{R}})^{\Omega}$$

$$f_{\text{lim}}^{\text{as}} \vec{z} := \lambda \omega. \limsup_{n \rightarrow \infty} f_n \omega$$

↳  $f_{\text{lim}}^{\text{as}}$  respects a.s. equality.

Thm (monotone convergence):

let  $\vec{\xi} \in W^{\mathbb{N} \times \Omega}$   $\lambda$ -a.s. monotone.

$$\vec{\xi} = \lim_{n \rightarrow \infty} \vec{\xi}_n \quad (\text{a.s.})$$



$$\int \lambda \vec{\xi} = \lim_{n \rightarrow \infty} \int \lambda \vec{\xi}_n$$

# Lebesgue space

( $\Omega$  Prob. Space,  
 $P \in [1, \infty)$ )

$$P: [1, \infty), \lambda: P_{\Omega} \quad L_{(\Omega, \lambda)}^P: \mathcal{B}(\mathbb{R}^{\Omega})$$

$$:= \left\{ \xi \in \mathbb{R}^{\Omega} \mid \int |\xi|^P < \infty \right\} \hookrightarrow \mathbb{R}^{\Omega}$$

$$\text{Ensemble } L_{\Omega} := \prod_{\substack{\lambda \in P_{\Omega} \\ P \in [1, \infty)}} L_{(\Omega, \lambda)}^P$$

$$L \quad P \leq q \Rightarrow L_{\Omega}^P \supseteq L_{\Omega}^q$$

$L^p$  semi norms

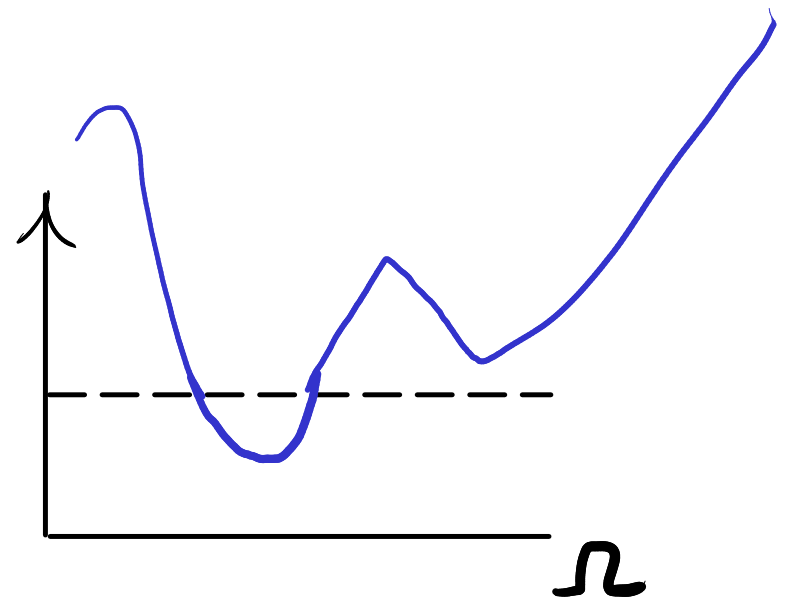
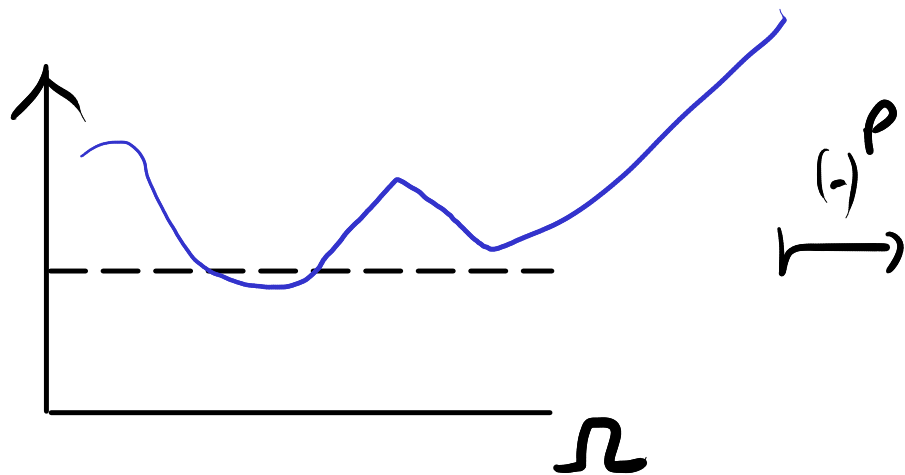
$$\| \cdot \| : \prod_{P, \lambda} L^p_{(\Omega, \lambda)} \rightarrow \mathbb{R}_{\geq 0} \quad \| \xi \|_p := \sqrt[p]{\int \lambda |\xi|^p}$$

$L^2$  inner product

$$\langle \cdot, \cdot \rangle : \prod_{P, \lambda} L^p_{(\Omega, \lambda)} \times L^p_{(\Omega, \lambda)} \rightarrow \mathbb{R}$$

$$\langle \xi, \eta \rangle_{P, \lambda} := \int \lambda \xi \eta$$





## Statistics

### Expectation

$$\mathbb{E} : \prod_{\lambda} L^1 \rightarrow \mathbb{R}$$

$$\mathbb{E}_{\lambda} \xi := \int_{\lambda} \xi$$

### Covariance and Correlation

$$\text{Cov}, \text{Corr} : \prod_{\lambda} L^2 \rightarrow \mathbb{R}$$

$$\text{Cov}(\xi, \zeta) := \langle \xi - \mathbb{E} \xi, \zeta - \mathbb{E} \zeta \rangle$$

$$\text{Corr}(\xi, \zeta) := \frac{\langle \xi, \zeta \rangle}{\|\xi\|_2 \cdot \|\zeta\|_2} = \cos(\text{angle}(\xi, \zeta))$$

# Sequential limits

$p \in [1, \infty)$ ,  $\mathcal{X}: \mathcal{P}(\mathcal{X})$  Cauchy  $L^p_{(\Omega, \mathcal{X})} \mathcal{B}(L^p_{(\mathcal{R}, \mathcal{H})})^{\mathbb{N}}$

$$= \left\{ \vec{z} \mid \forall \varepsilon \in \mathbb{Q}^+ \exists N \in \mathbb{N} \forall m, n \geq N \right. \\ \left. \left\| \sum_{k=n}^m z_k \right\|_p < \varepsilon \right\}$$

Thm:  $L^p_{\Omega}$  is Cauchy-complete

lim: Cauchy  $L^p \rightarrow L^p$  (convergence in mean)

Why?

1. Every Cauchy sequence has an a.s. converging subseq.
2. We can find it measurably

# Example

Then (dominated convergence)

For  $\vec{\Sigma}_n, \vec{\Sigma} \in L^1$  s.t.  $\vec{\Sigma}_n \leq \vec{\Sigma}$  a.s.:

1.  $\lim^{as} \vec{\Sigma} \in L^1$

2.  $\lim^1 \vec{\Sigma} = \lim^{as} \vec{\Sigma}$

3.  $\lim_{n \rightarrow \infty} \int \vec{\Sigma}_n = \int \lim_{n \rightarrow \infty} \vec{\Sigma}_n$

# Separability

Def:  $L^p$  separable: has countable dense subset

Fact: Separability is property of  $\mathbb{R}^2$ :

TFAE:

-  $\exists p \geq 1. L^p$  separable

-  $\forall p \geq 1. L^p$  separable

Measurably separability in  $I \hookrightarrow P\Omega \times [1, \infty)$

$$\vec{\beta} : \prod_{(\lambda, \rho) \in I} L^p_{(\Omega, \lambda)} \quad \text{s.t.}$$

$$\{ \vec{\beta}_n^{\lambda, \rho} \mid n \in \mathbb{N} \} \text{ dense in } L^p_{(\Omega, \lambda)}$$

Prop. - Every sbs  $S$  measurably separable in  $P_S \times [1, \infty)$

-  $I \hookrightarrow P\Omega \times \{2\}$  measurably separable

$$\Rightarrow \exists \vec{\beta} \in \prod_{\lambda \in I} L^2_{(\Omega, \lambda)} \text{ orthonormal system}$$

$$\begin{aligned} \langle \beta_n, \beta_m \rangle &= 0 \\ \|\beta_n\|_2 &= 1 \\ (\beta_n) &\text{ dense} \end{aligned}$$

## Example

Let  $S \hookrightarrow L^2$  closed vector subspace.

Orthogonal decomposition / linear in fact.

$$\langle P, P^\perp \rangle : L^2 \rightarrow S \times S^\perp$$

When  $S$  is separable with orthonormal system  $\beta$

We have a measurable version of

$$\langle P, P^\perp \rangle : L^2 \rightarrow S \times S^\perp$$

$$P \xi := \sum_{n=0}^{\infty} \langle \xi, \beta_n \rangle \beta_n$$

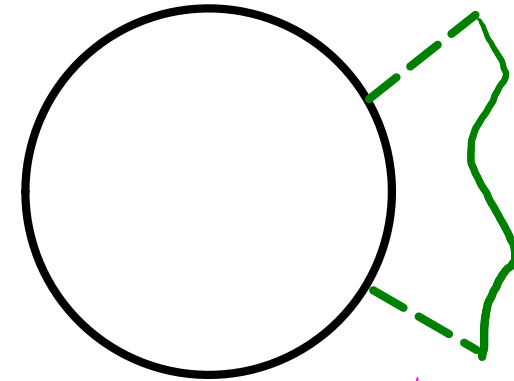
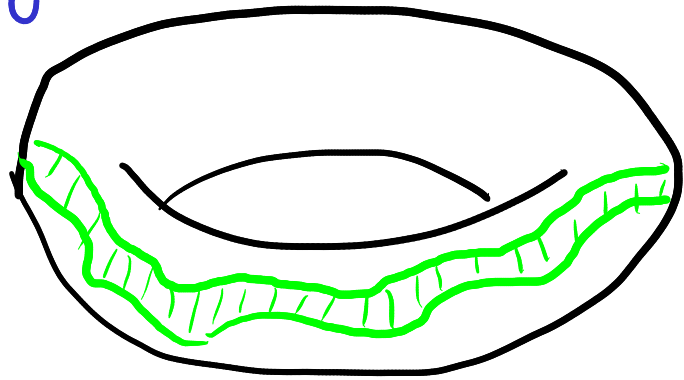
$$P^\perp := \text{Id} - P.$$

# Kolmogorov's Conditional Expectation

ground truth space

$\mathcal{H}$

Sample space



$\Sigma$   
Statistic  
of interest

Conditional  
expectation  
 $E[\Sigma | H = -]$   
Observed  
statistic

$\mathbb{R}$

# Kolmogorov's Conditional Expectation

A conditional expectation

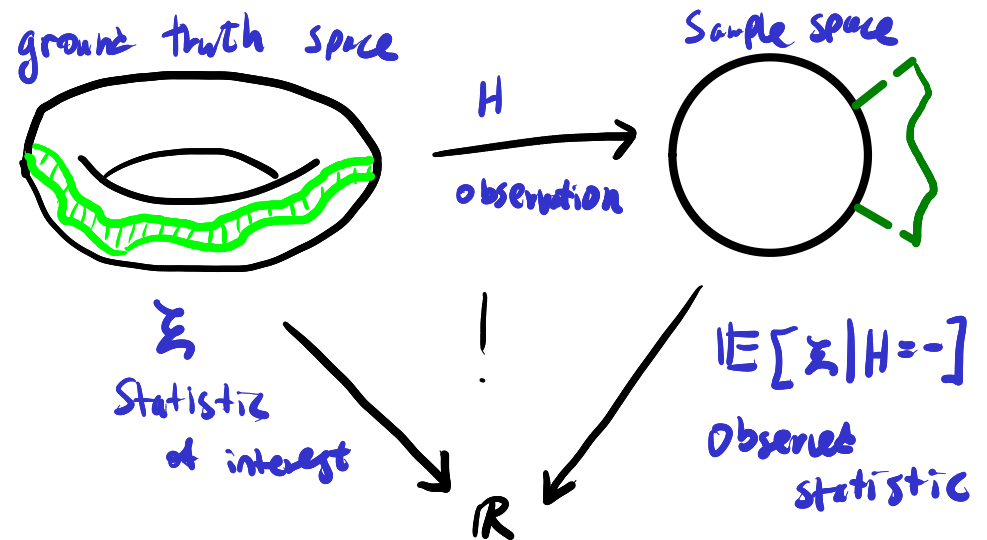
of  $Z \in \mathcal{L}'_\Omega$  wrt

$H: \Omega \rightarrow \mathbb{H}$  is

$Z \in \mathcal{L}'_{\mathbb{H}}$  s.t. for all  $A \in \mathcal{B}_{\mathbb{H}}$ :

$$\int_A \mu Z = \int_{H^{-1}[A]} \lambda Z$$

where  $\mu := \lambda_H$



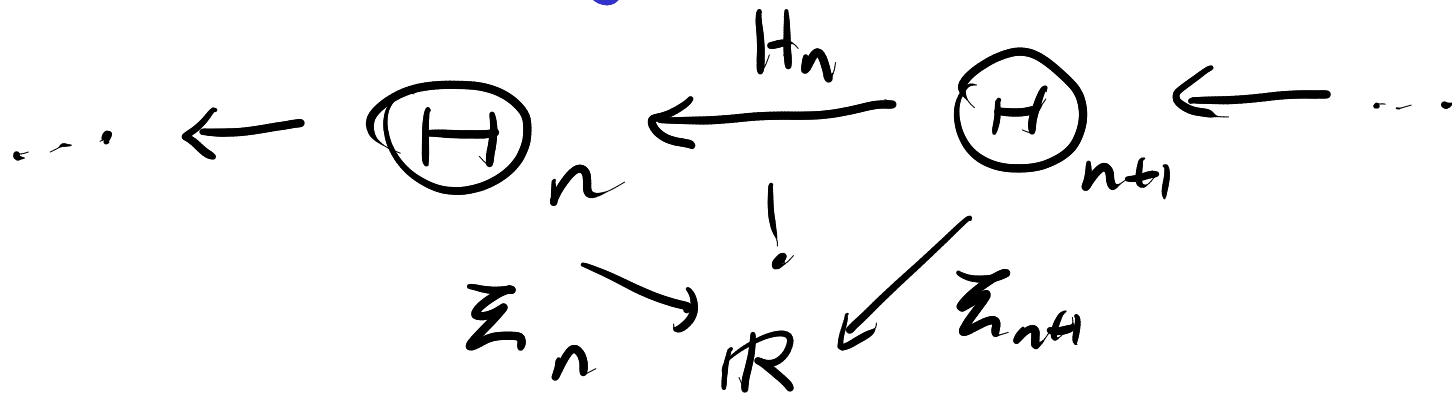


# Conditional expectations

1. unique a.s.

2. fundamental to modern Probability, eg.:

a martingale



s.t.  $\xi_n = \mathbb{E}[\xi_{n+1} | H_n = \cdot]$

Thm (Existence)

-  $\exists \mathbb{E}[-|\mathcal{H}=-]: \mathcal{L}'_{(\Omega, \lambda)} \rightarrow \mathcal{L}'_{(\mathbb{H}, \mu)}$

- When  $(\Omega, \lambda)$  is separable

$$\mathbb{E}[-|\mathcal{H}=-]: \mathcal{L}'_{(\Omega, \lambda)} \rightarrow \mathcal{L}'_{(\mathbb{H}, \mu)}$$

- When  $\mathbb{H}$  is  $\mathcal{I}$ -measurably separable

$$\mathbb{E}[-|\mathcal{H}=-]: \prod_{\substack{\mathbb{H} \in \mathbb{H} \\ \lambda \in \mathcal{H}_*^{\mathcal{I}}[\mathbb{I}]}} \mathcal{L}'_{(\Omega, \lambda)} \rightarrow \mathcal{L}'_{(\mathbb{H}, \mu)}$$

## Plan:

- 1) Type-driven probability: discrete case (Mon + Tue)
- 2) Borel sets & measurable spaces (Wed)
- 3) Quasi Borel spaces (Wed) Simple type structure (Thu)
- 4) Dependent type structure & standard Borel spaces (Thu)
- 5) Integration & random variables (Fri)



Course  
web  
page

# Discrete model

type : set       $W := [0, \infty]$        $\mathcal{B}_X := \mathcal{P}X$

$\mathcal{D}X := \{ \mu : X \rightarrow W \mid \text{supp } \mu \text{ countable} \}$

$\mathcal{P}X := \{ \mu \in \mathcal{D}X \mid \sum_{\mu} C_{\mu}[X] = 1 \}$

$C_{\mu}[E] := \sum_{x \in E} \mu x$        $\delta_x := \lambda x' . \begin{cases} x = x' : 0 \\ x \neq x' : 1 \end{cases}$

$\oint \mu k := \lambda x . \sum_{r \in \Gamma} \mu r . k(r; x)$

# Full model

$$\text{type} : \text{Obs} \quad w := [0, \infty] \quad \mathcal{B}_X \cong \mathcal{B}^X$$

$$DX := \left( \left\{ \lambda_\alpha \mid \alpha : \mathbb{R} \rightarrow X \right\}, \left\{ \lambda_{r, \lambda} \mid \alpha : \mathbb{R} \times \mathbb{R} \rightarrow X \right\} \right)$$

$$P_X := \left\{ \mu \in DX \mid \int_{\mu} C_e[X] = 1 \right\}$$

$$C_e[E] := \int_{\mu} E \quad \delta_x := E \mapsto \begin{cases} x \in E : 1 \\ x \notin E : 0 \end{cases}$$

$$\oint \mu k := \int E. \int \mu(\lambda) k(\lambda; E)$$